

BGP Techniques for Internet Service Providers

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Presentation Slides

Will be available on

ftp://ftp-eng.cisco.com

/pfs/seminars/APRICOT2007-BGP-Techniques.pdf

And on the APRICOT 2007 website

Feel free to ask questions any time

BGP Techniques for Internet Service Providers

- BGP Basics
- Scaling BGP
- Using Communities
- Deploying BGP in an ISP network



BGP Basics

What is this BGP thing?

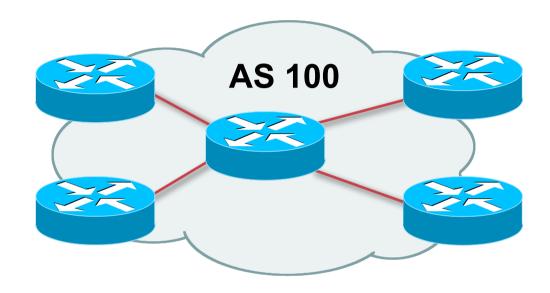
Border Gateway Protocol

Routing Protocol used to exchange routing information between networks

exterior gateway protocol

- Described in RFC4271
 - RFC4276 gives an implementation report on BGP-4 RFC4277 describes operational experiences using BGP-4
- The Autonomous System is BGP's fundamental operating unit It is used to uniquely identify networks with common routing policy

Autonomous System (AS)



- Collection of networks with same routing policy
- Single routing protocol
- Usually under single ownership, trust and administrative control
- Identified by a unique number

Autonomous System Number (ASN)

An ASN is a 16 bit number

1-64511 are assigned by the RIRs

64512-65534 are for private use and should never appear on the Internet

0 and 65535 are reserved

32 bit ASNs are here now!

www.ietf.org/internet-drafts/draft-ietf-idr-as4bytes-13.txt

www.ietf.org/internet-drafts/draft-michaelson-4byte-as-representation-02.txt

www.ietf.org/internet-drafts/draft-rekhter-as4octet-ext-community-01.txt

www.apnic.net/docs/policy/proposals/prop-032-v002.html

With AS 23456 reserved for the transition

Autonomous System Number (ASN)

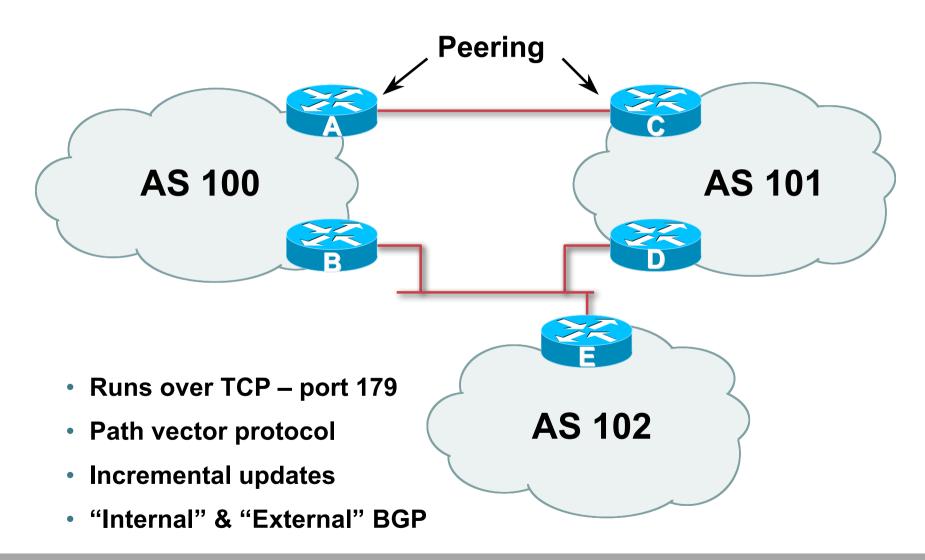
- ASNs are distributed by the Regional Internet Registries
- Also available from upstream ISPs who are members of one of the RIRs

Current ASN allocations up to 43007 have been made to the RIRs

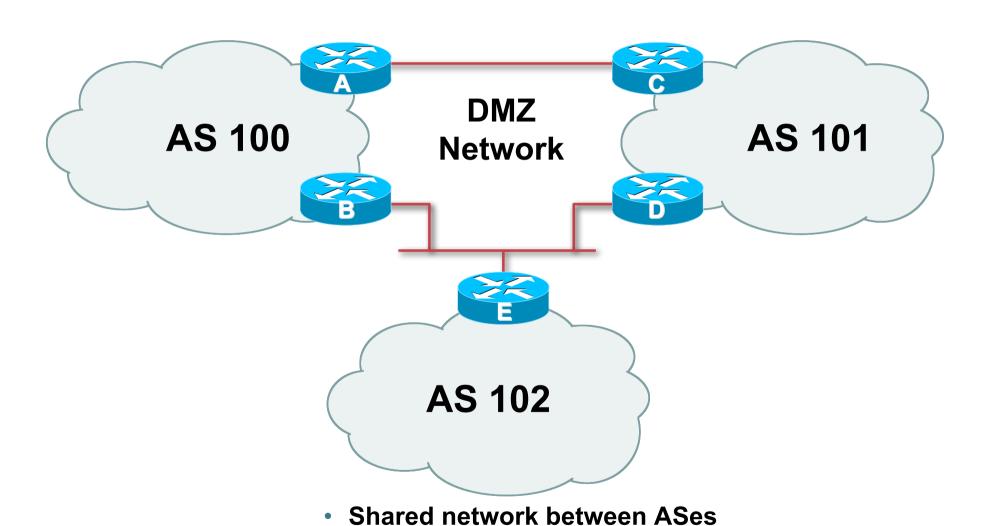
Of these, around 24500 are visible on the Internet

See www.iana.org/assignments/as-numbers

BGP Basics



Demarcation Zone (DMZ)



BGP General Operation

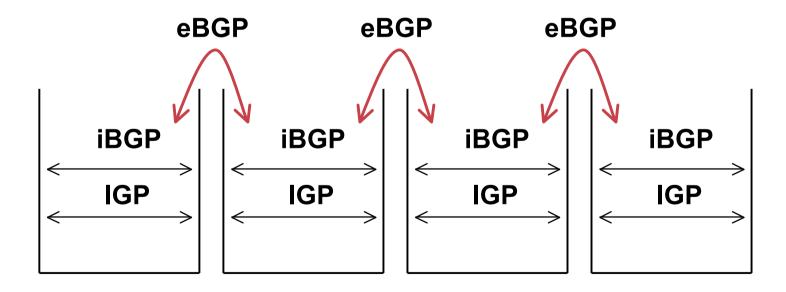
- Learns multiple paths via internal and external BGP speakers
- Picks the best path and installs in the forwarding table
- Best path is sent to external BGP neighbours
- Policies applied by influencing the best path selection

eBGP & iBGP

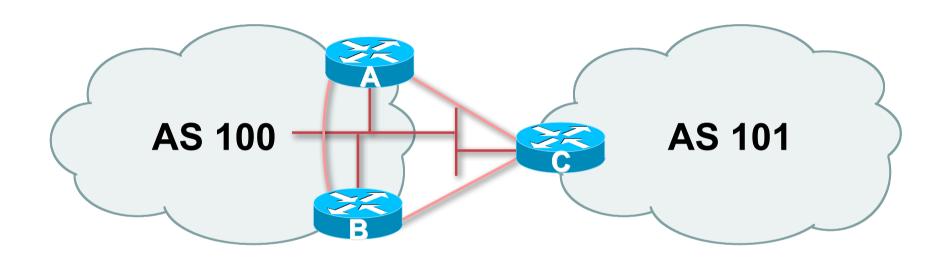
- BGP used internally (iBGP) and externally (eBGP)
- iBGP used to carry some/all Internet prefixes across ISP backbone ISP's customer prefixes
- eBGP used to
 exchange prefixes with other ASes
 implement routing policy

BGP/IGP model used in ISP networks

Model representation



External BGP Peering (eBGP)

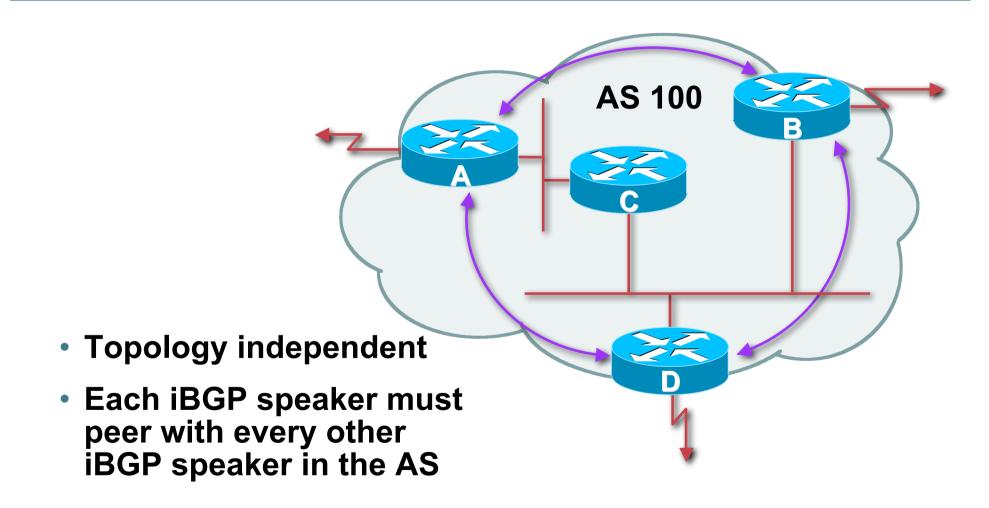


- Between BGP speakers in different AS
- Should be directly connected
- Never run an IGP between eBGP peers

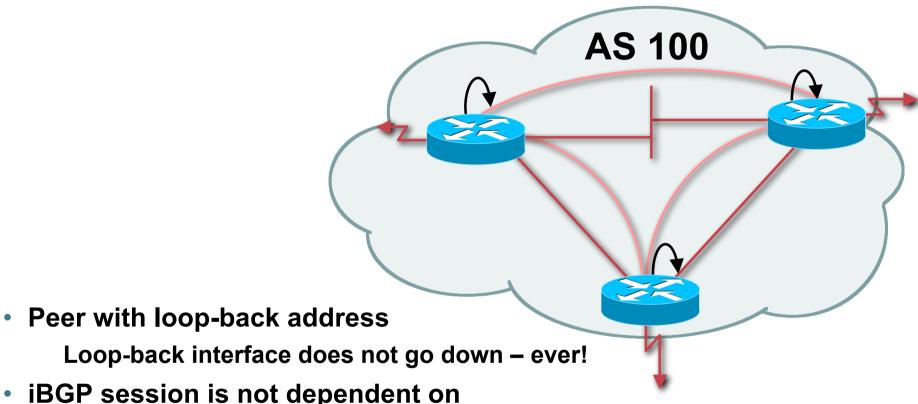
Internal BGP (iBGP)

- BGP peer within the same AS
- Not required to be directly connected
 IGP takes care of inter-BGP speaker connectivity
- iBGP speakers need to be fully meshed they originate connected networks they do not pass on prefixes learned from other iBGP speakers

Internal BGP Peering (iBGP)



Peering to loopback addresses



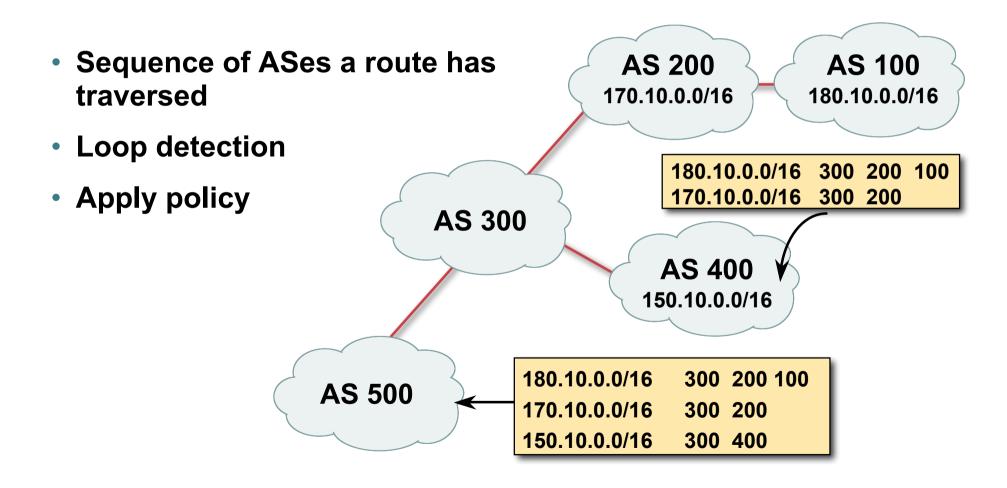
 iBGP session is not dependent on State of a single interface Physical topology



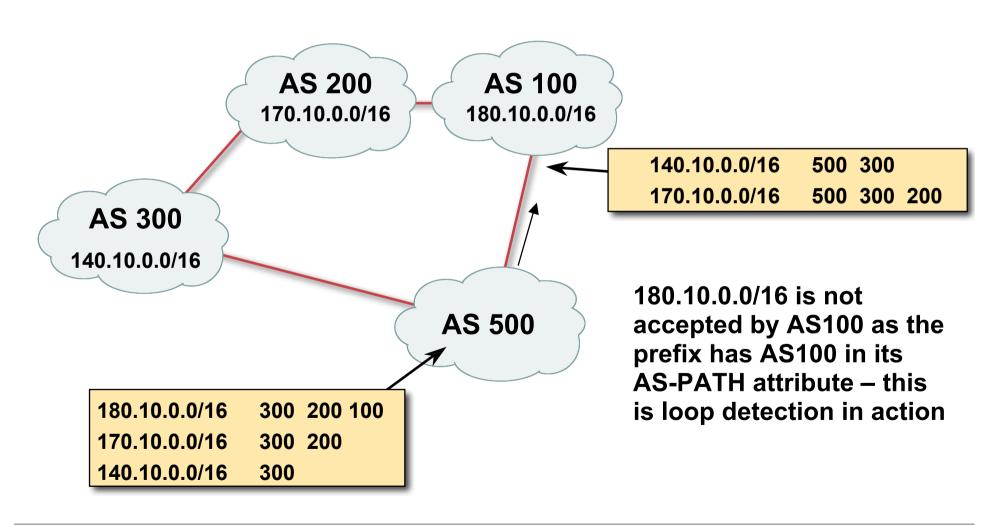
BGP Attributes

Information about BGP

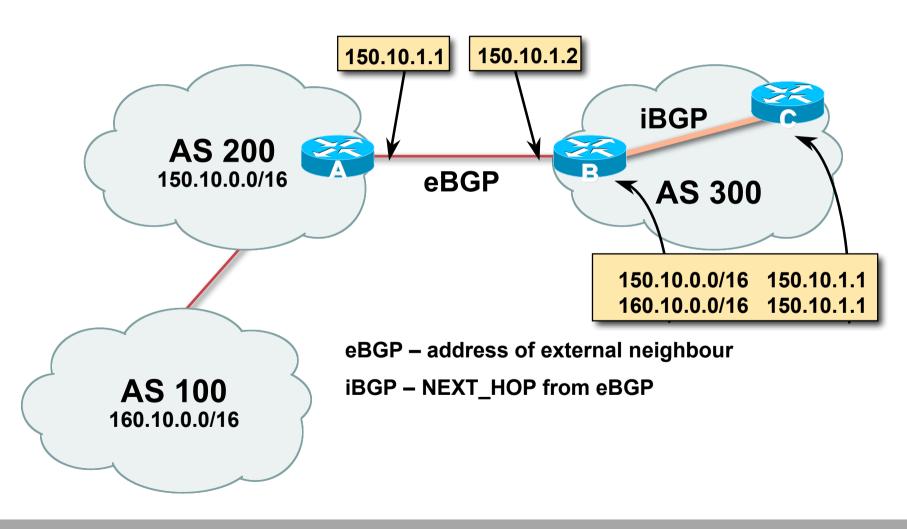
AS-Path



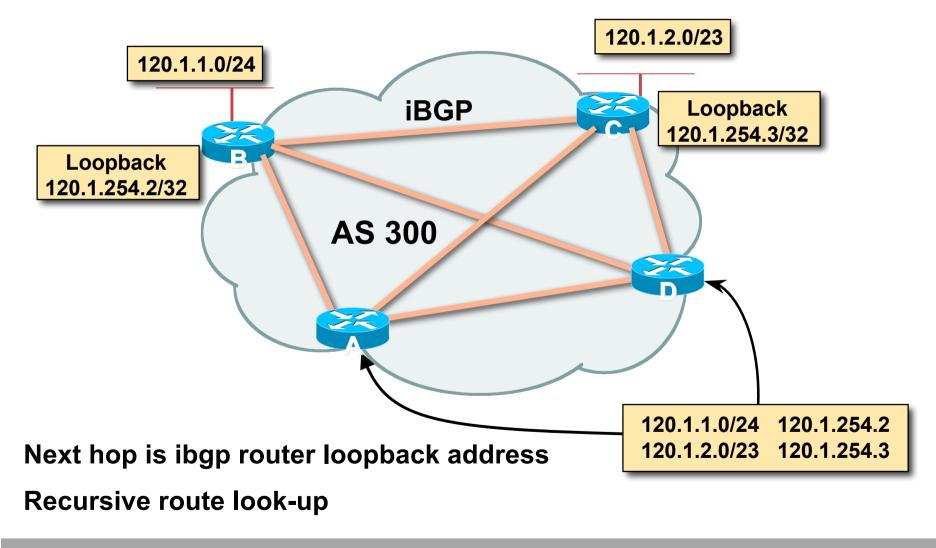
AS-Path loop detection



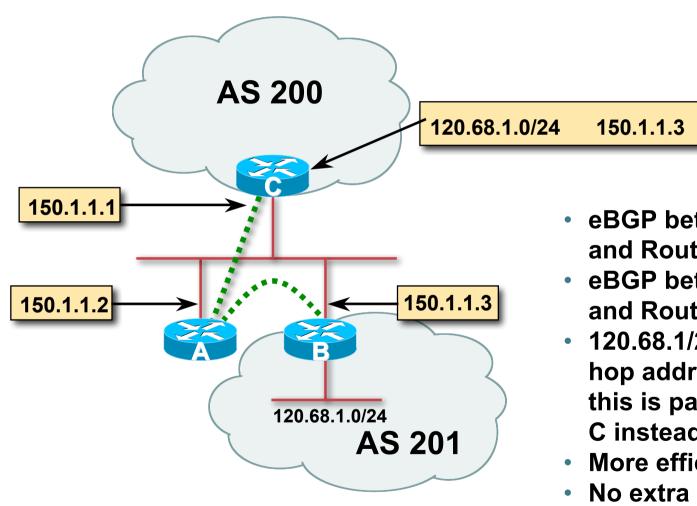
Next Hop



iBGP Next Hop



Third Party Next Hop



- eBGP between Router A and Router C
- eBGP between Router A and Router B
- 120.68.1/24 prefix has next hop address of 150.1.1.3 this is passed on to Router C instead of 150.1.1.2
- More efficient
- No extra config needed

Next Hop (Summary)

- IGP should carry route to next hops
- Recursive route look-up
- Unlinks BGP from actual physical topology
- Allows IGP to make intelligent forwarding decision

Origin

- Conveys the origin of the prefix
- Historical attribute

Used in transition from EGP to BGP

- Influences best path selection
- Three values: IGP, EGP, incomplete

IGP – generated by BGP network statement

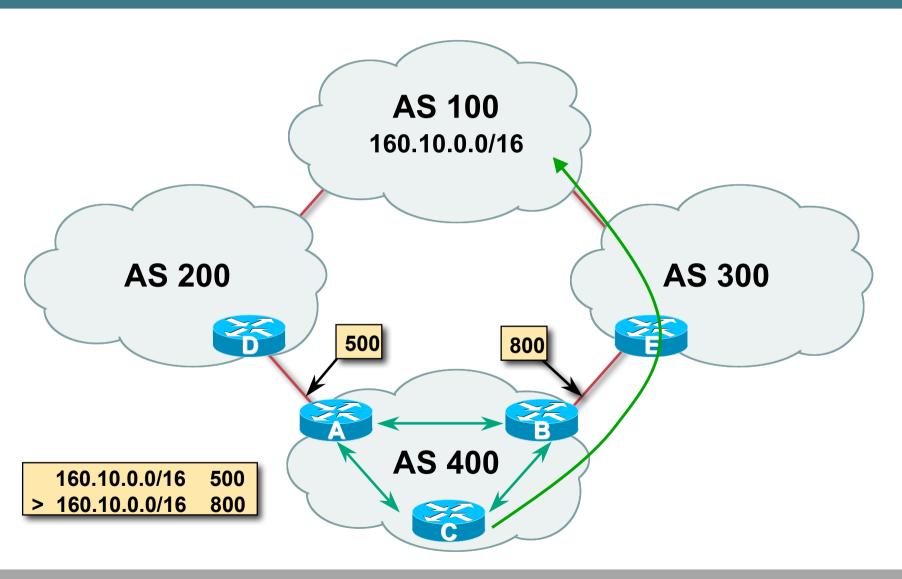
EGP – generated by **EGP**

incomplete – redistributed from another routing protocol

Aggregator

- Conveys the IP address of the router or BGP speaker generating the aggregate route
- Useful for debugging purposes
- Does not influence best path selection

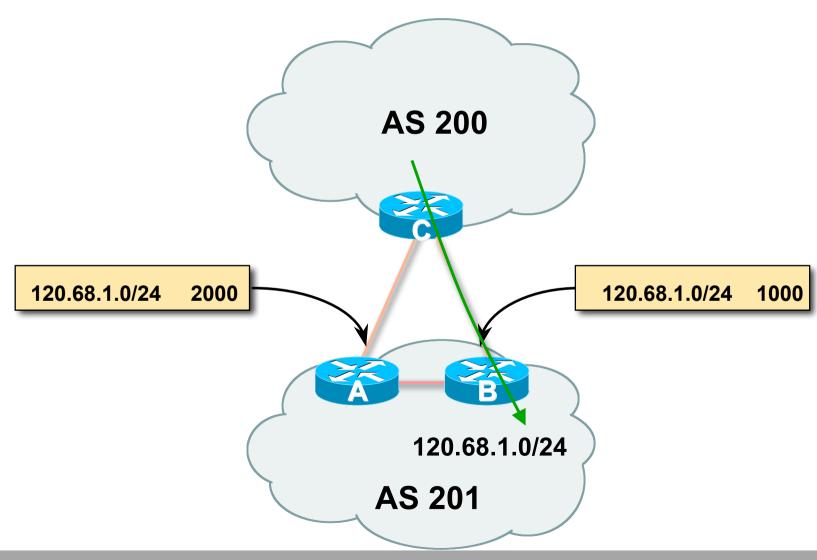
Local Preference



Local Preference

- Local to an AS non-transitive
 Default local preference is 100 (IOS)
- Used to influence BGP path selection determines best path for outbound traffic
- Path with highest local preference wins

Multi-Exit Discriminator (MED)



Multi-Exit Discriminator

- Inter-AS non-transitive & optional attribute
- Used to convey the relative preference of entry points determines best path for *inbound* traffic
- Comparable if paths are from same AS
 bgp always-compare-med allows comparisons of MEDs from different ASes
- Path with lowest MED wins
- Absence of MED attribute implies MED value of zero (RFC4271)

Multi-Exit Discriminator "metric confusion"

MED is non-transitive and optional attribute

Some implementations send learned MEDs to iBGP peers by default, others do not

Some implementations send MEDs to eBGP peers by default, others do not

 Default metric value varies according to vendor implementation

Original BGP spec made no recommendation

Some implementations said no metric was equivalent to 2^{32} -1 (the highest possible) or 2^{32} -2

Other implementations said no metric was equivalent to 0

Potential for "metric confusion"

Community

Communities are described in RFC1997

Transitive and Optional Attribute

32 bit integer

Represented as two 16 bit integers (RFC1998)

Common format is < local-ASN>:xx

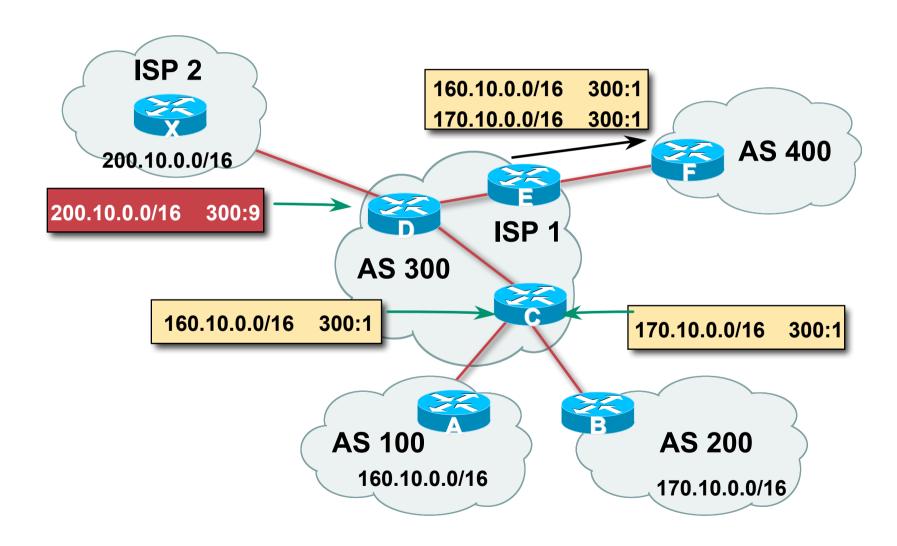
0:0 to 0:65535 and 65535:0 to 65535:65535 are reserved

Used to group destinations

Each destination could be member of multiple communities

Very useful in applying policies within and between ASes

Community



Well-Known Communities

Several well known communities

www.iana.org/assignments/bgp-well-known-communities

no-export

65535:65281

do not advertise to any eBGP peers

no-advertise

65535:65282

do not advertise to any BGP peer

no-export-subconfed

65535:65283

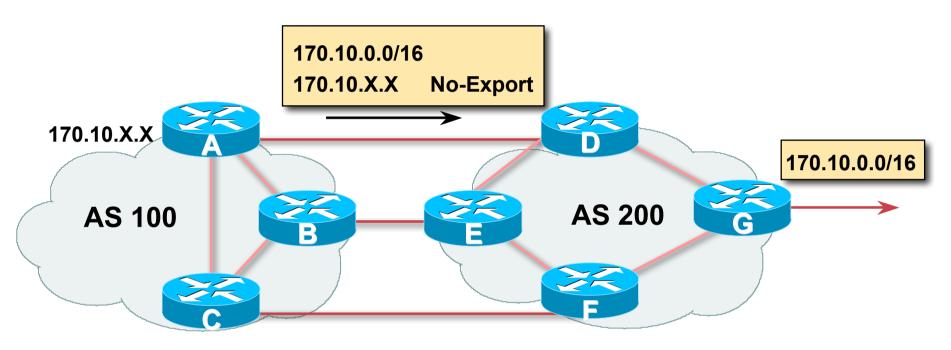
do not advertise outside local AS (only used with confederations)

no-peer

65535:65284

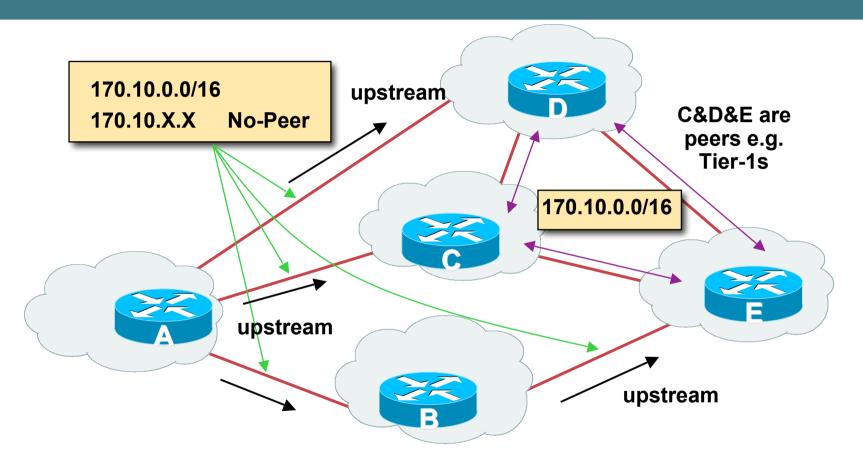
do not advertise to bi-lateral peers (RFC3765)

No-Export Community



- AS100 announces aggregate and subprefixes
 aim is to improve loadsharing by leaking subprefixes
- Subprefixes marked with no-export community
- Router G in AS200 does not announce prefixes with no-export community set

No-Peer Community



 Sub-prefixes marked with no-peer community are not sent to bilateral peers

They are only sent to upstream providers

Community Implementation details

Community is an optional attribute

Some implementations send communities to iBGP peers by default, some do not

Some implementations send communities to eBGP peers by default, some do not

Being careless can lead to community "confusion"

ISPs need consistent community policy within their own networks

And they need to inform peers, upstreams and customers about their community expectations



BGP Path Selection Algorithm

Why Is This the Best Path?

BGP Path Selection Algorithm for IOS Part One

- Do not consider path if no route to next hop
- Do not consider iBGP path if not synchronised (Cisco IOS)
- Highest weight (local to router)
- Highest local preference (global within AS)
- Prefer locally originated route
- Shortest AS path

BGP Path Selection Algorithm for IOS Part Two

- Lowest origin code
 IGP < EGP < incomplete
- Lowest Multi-Exit Discriminator (MED)

If bgp deterministic-med, order the paths before comparing If bgp always-compare-med, then compare for all paths otherwise MED only considered if paths are from the same AS (default)

BGP Path Selection Algorithm for IOS Part Three

- Prefer eBGP path over iBGP path
- Path with lowest IGP metric to next-hop
- Lowest router-id (originator-id for reflected routes)
- Shortest Cluster-List

Client must be aware of Route Reflector attributes!

Lowest neighbour IP address

BGP Path Selection Algorithm

In multi-vendor environments:

Make sure the path selection processes are understood for each brand of equipment

Each vendor has slightly different implementations, extra steps, extra features, etc

Watch out for possible MED confusion



Applying Policy with BGP

Control!

Applying Policy in BGP: Why?

Policies are applied to:

Influence BGP Path Selection by setting BGP attributes

Determine which prefixes are announced or blocked

Determine which AS-paths are preferred, permitted, or denied

Determine route groupings and their effects

Decisions are generally based on prefix, AS-path and community

Applying Policy with BGP: Tools

 Most implementations have tools to apply policies to BGP:

Prefix manipulation/filtering

AS-PATH manipulation/filtering

Community Attribute setting and matching

 Implementations also have policy language which can do various match/set constructs on the attributes of chosen BGP routes



Extending BGP

- Documented in RFC2842
- Capabilities parameters passed in BGP open message
- Unknown or unsupported capabilities will result in NOTIFICATION message
- Codes:

0 to 63 are assigned by IANA by IETF consensus64 to 127 are assigned by IANA "first come first served"128 to 255 are vendor specific

Current capabilities are:

0	Reserved	[RFC3392]
1	Multiprotocol Extensions for BGP-4	[RFC4760]
2	Route Refresh Capability for BGP-4	[RFC2918]
3	Cooperative Route Filtering Capability	[ID]
4	Multiple routes to a destination capability	[RFC3107]
64	Graceful Restart Capability	[RFC4724]
65	Support for 4 octet ASNs	[ID]
66	Deprecated 2003-03-06	
67	Support for Dynamic Capability	[ID]

See www.iana.org/assignments/capability-codes

Multiprotocol extensions

This is a whole different world, allowing BGP to support more than IPv4 unicast routes

Examples include: v4 multicast, IPv6, v6 multicast, VPNs

Another tutorial (or many!)

- Route refresh is a well known scaling technique covered shortly
- The other capabilities are still in development or not widely implemented or deployed yet

BGP for Internet Service Providers

- BGP Basics
- Scaling BGP
- Using Communities
- Deploying BGP in an ISP network



BGP Scaling Techniques

BGP Scaling Techniques

- How does a service provider:
 - Scale the iBGP mesh beyond a few peers?
 - Implement new policy without causing flaps and route churning?
 - Keep the network stable, scalable, as well as simple?

BGP Scaling Techniques

- Route Refresh
- Route Reflectors
- Confederations



Dynamic Reconfiguration

Route Refresh

Route Refresh

BGP peer reset required after every policy change

Because the router does not store prefixes which are rejected by policy

Hard BGP peer reset:

Terminates BGP peering & Consumes CPU
Severely disrupts connectivity for all networks

Soft BGP peer reset (or Route Refresh):

BGP peering remains active

Impacts only those prefixes affected by policy change

Route Refresh Capability

- Facilitates non-disruptive policy changes
- For most implementations, no configuration is needed

Automatically negotiated at peer establishment

- No additional memory is used
- Requires peering routers to support "route refresh capability" – RFC2918

Dynamic Reconfiguration

- Use Route Refresh capability if supported find out from the BGP neighbour status display Non-disruptive, "Good For the Internet"
- If not supported, see if implementation has a workaround
- Only hard-reset a BGP peering as a last resort

Consider the impact to be equivalent to a router reboot



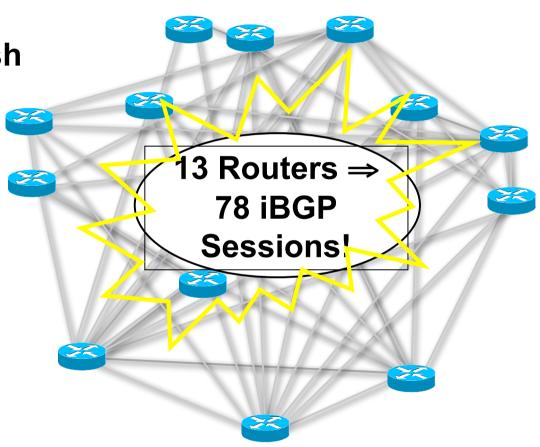
Route Reflectors

Scaling the iBGP mesh

Scaling iBGP mesh

Avoid ½n(n-1) iBGP mesh

n=1000 ⇒ nearly half a million ibgp sessions!

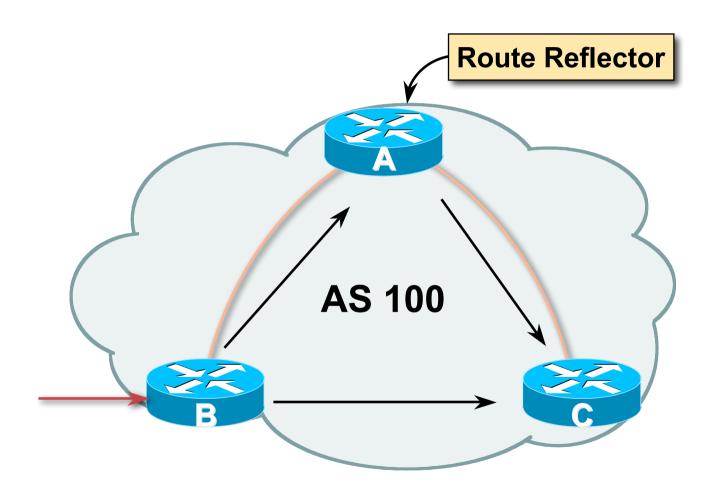


Two solutions

Route reflector – simpler to deploy and run

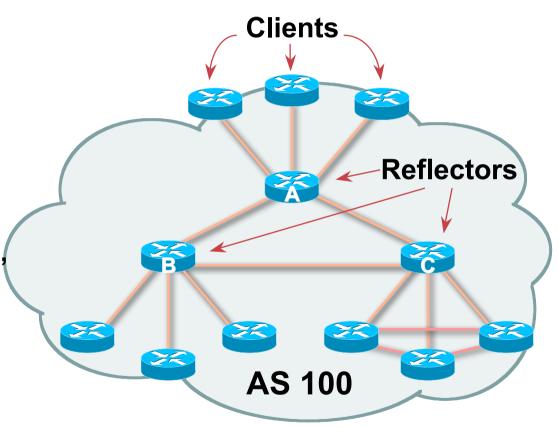
Confederation – more complex, has corner case advantages

Route Reflector: Principle



Route Reflector

- Reflector receives path from clients and non-clients
- Selects best path
- If best path is from client, reflect to other clients and non-clients
- If best path is from non-client, reflect to clients only
- Non-meshed clients
- Described in RFC4456



Route Reflector Topology

- Divide the backbone into multiple clusters
- At least one route reflector and few clients per cluster
- Route reflectors are fully meshed
- Clients in a cluster could be fully meshed
- Single IGP to carry next hop and local routes

Route Reflectors: Loop Avoidance

Originator_ID attribute

Carries the RID of the originator of the route in the local AS (created by the RR)

Cluster_list attribute

The local cluster-id is added when the update is sent by the RR

Best to set cluster-id is from router-id (address of loopback)

(Some ISPs use their own cluster-id assignment strategy – but needs to be well documented!)

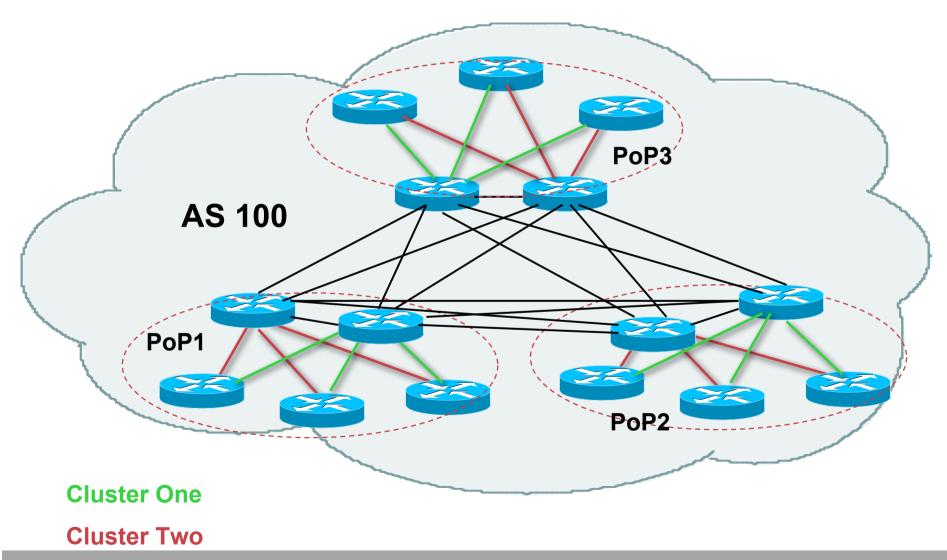
Route Reflectors: Redundancy

- Multiple RRs can be configured in the same cluster
 - not advised!

All RRs in the cluster must have the same cluster-id (otherwise it is a different cluster)

- A router may be a client of RRs in different clusters
 - Common today in ISP networks to overlay two clusters redundancy achieved that way
 - → Each client has two RRs = redundancy

Route Reflectors: Redundancy



Route Reflector: Benefits

- Solves iBGP mesh problem
- Packet forwarding is not affected
- Normal BGP speakers co-exist
- Multiple reflectors for redundancy
- Easy migration
- Multiple levels of route reflectors

Route Reflectors: Migration

Where to place the route reflectors?

Always follow the physical topology!

This will guarantee that the packet forwarding won't be affected

Typical ISP network:

PoP has two core routers

Core routers are RR for the PoP

Two overlaid clusters

Route Reflectors: Migration

Typical ISP network:

Core routers have fully meshed iBGP
Create further hierarchy if core mesh too big
Split backbone into regions

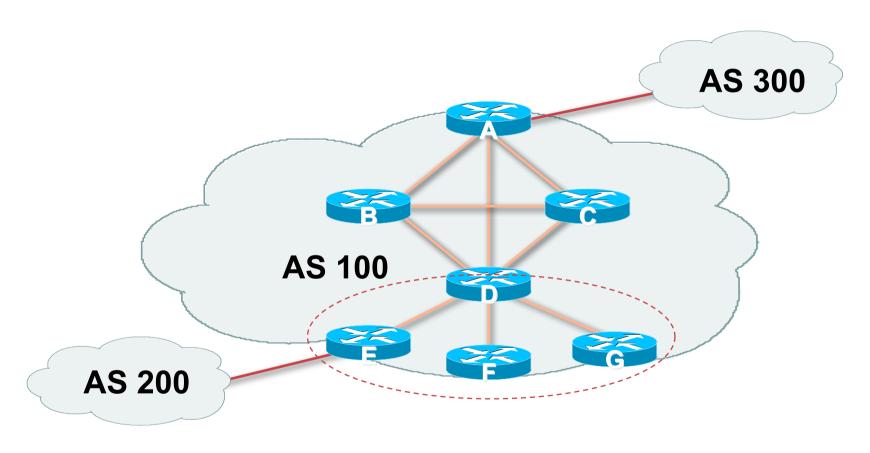
Configure one cluster pair at a time

Eliminate redundant iBGP sessions

Place maximum one RR per cluster

Easy migration, multiple levels

Route Reflector: Migration



 Migrate small parts of the network, one part at a time



BGP Confederations

Confederations

Divide the AS into sub-AS

eBGP between sub-AS, but some iBGP information is kept

Preserve NEXT_HOP across the sub-AS (IGP carries this information)

Preserve LOCAL_PREF and MED

- Usually a single IGP
- Described in RFC3065

Confederations (Cont.)

 Visible to outside world as single AS – "Confederation Identifier"

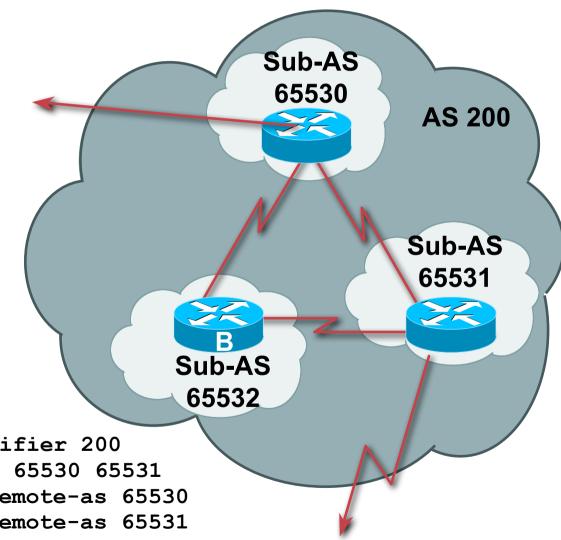
Each sub-AS uses a number from the private AS range (64512-65534)

iBGP speakers in each sub-AS are fully meshed

The total number of neighbours is reduced by limiting the full mesh requirement to only the peers in the sub-AS

Can also use Route-Reflector within sub-AS

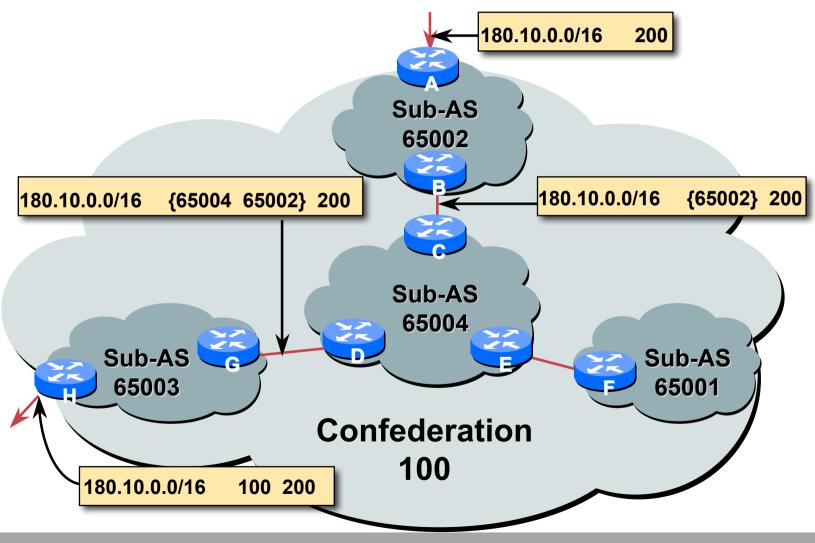
Confederations



Configuration (rtr B):

bgp confederation identifier 200 bgp confederation peers 65530 65531 neighbor 141.153.12.1 remote-as 65530 neighbor 141.153.17.2 remote-as 65531

Confederations: AS-Sequence



Route Propagation Decisions

Same as with "normal" BGP:

From peer in same sub-AS → only to external peers

From external peers → to all neighbors

"External peers" refers to

Peers outside the confederation

Peers in a different sub-AS

Preserve LOCAL_PREF, MED and NEXT_HOP

RRs or Confederations

	Internet Connectivity	Multi-Level Hierarchy	Policy Control	Scalability	Migration Complexity
Confederations	Anywhere in the Network	Yes	Yes	Medium	Medium to High
Route Reflectors	Anywhere in the Network	Yes	Yes	Very High	Very Low

Most new service provider networks now deploy Route Reflectors from Day One

More points about confederations

 Can ease "absorbing" other ISPs into you ISP – e.g., if one ISP buys another

Or can use AS masquerading feature available in some implementations to do a similar thing

 Can use route-reflectors with confederation sub-AS to reduce the sub-AS iBGP mesh

BGP Scaling Techniques

Route Refresh

Use should be mandatory

Route Reflectors/Confederations
 The only way to scale iBGP mesh



Route Flap Damping

Network Stability for the 1990s

Network Instability for the 21st Century!

Route Flap Damping

- For many years, Route Flap Damping was a strongly recommended practice
- Now it is strongly discouraged as it causes far greater network instability than it cures
- But first, the theory...

Route Flap Damping

Route flap

Going up and down of path or change in attribute

BGP WITHDRAW followed by UPDATE = 1 flap

eBGP neighbour going down/up is NOT a flap

Ripples through the entire Internet

Wastes CPU

Damping aims to reduce scope of route flap propagation

Route Flap Damping (continued)

Requirements

Fast convergence for normal route changes

History predicts future behaviour

Suppress oscillating routes

Advertise stable routes

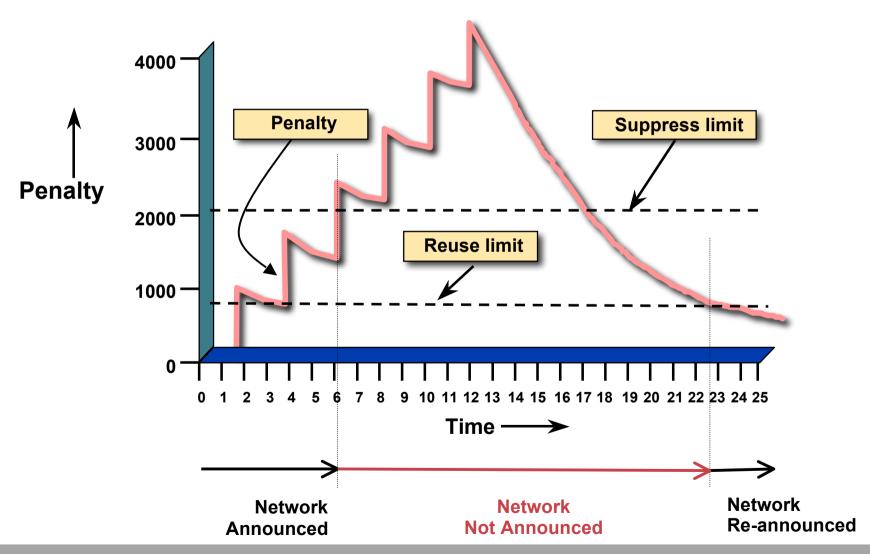
Implementation described in RFC 2439

Operation

Add penalty (1000) for each flap
 Change in attribute gets penalty of 500

- Exponentially decay penalty half life determines decay rate
- Penalty above suppress-limit do not advertise route to BGP peers
- Penalty decayed below reuse-limit
 re-advertise route to BGP peers
 penalty reset to zero when it is half of reuse-limit

Operation



Operation

- Only applied to inbound announcements from eBGP peers
- Alternate paths still usable
- Controllable by at least:

Half-life

reuse-limit

suppress-limit

maximum suppress time

Configuration

 Implementations allow various policy control with flap damping

Fixed damping, same rate applied to all prefixes

Variable damping, different rates applied to different ranges of prefixes and prefix lengths

Route Flap Damping History

- First implementations on the Internet by 1995
- Vendor defaults too severe

RIPE Routing Working Group recommendations in ripe-178, ripe-210, and ripe-229

http://www.ripe.net/ripe/docs

But many ISPs simply switched on the vendors' default values without thinking

Serious Problems:

 "Route Flap Damping Exacerbates Internet Routing Convergence"

Zhuoqing Morley Mao, Ramesh Govindan, George Varghese & Randy H. Katz, August 2002

- "What is the sound of one route flapping?"
 Tim Griffin, June 2002
- Various work on routing convergence by Craig Labovitz and Abha Ahuja a few years ago
- "Happy Packets"

Closely related work by Randy Bush et al

Problem 1:

One path flaps:

BGP speakers pick next best path, announce to all peers, flap counter incremented

Those peers see change in best path, flap counter incremented

After a few hops, peers see multiple changes simply caused by a single flap → prefix is suppressed

Problem 2:

- Different BGP implementations have different transit time for prefixes
 - Some hold onto prefix for some time before advertising Others advertise immediately
- Race to the finish line causes appearance of flapping, caused by a simple announcement or path change → prefix is suppressed

Solution:

- Do NOT use Route Flap Damping whatever you do!
- RFD will unnecessarily impair access

to your network and

to the Internet

 More information contained in RIPE Routing Working Group recommendations:

www.ripe.net/ripe/docs/ripe-378.[pdf,html,txt]

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Service Provider use of Communities

Some examples of how ISPs make life easier for themselves

BGP Communities

- Another ISP "scaling technique"
- Prefixes are grouped into different "classes" or communities within the ISP network
- Each community means a different thing, has a different result in the ISP network

BGP Communities

 Communities are generally set at the edge of the ISP network

Customer edge: customer prefixes belong to different communities depending on the services they have purchased

Internet edge: transit provider prefixes belong to difference communities, depending on the loadsharing or traffic engineering requirements of the local ISP, or what the demands from its BGP customers might be

Two simple examples follow to explain the concept

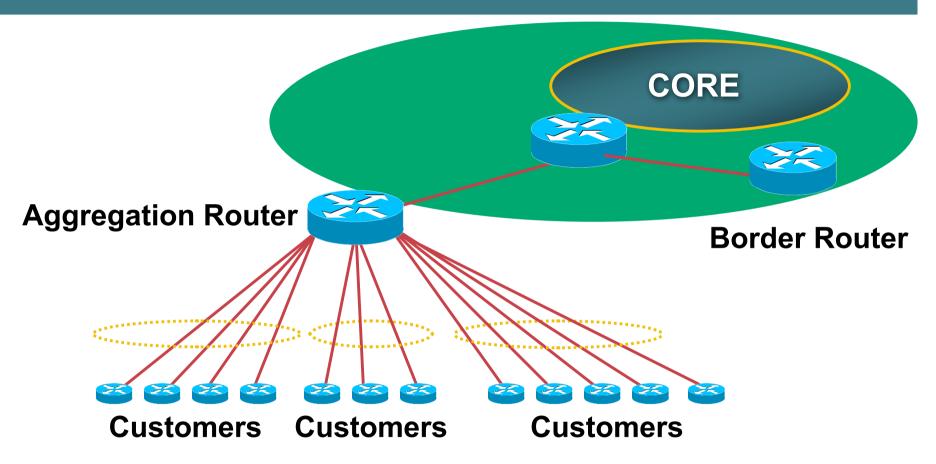
- This demonstrates how communities might be used at the customer edge of an ISP network
- ISP has three connections to the Internet:
 - IXP connection, for local peers
 - Private peering with a competing ISP in the region
 - Transit provider, who provides visibility to the entire Internet
- Customers have the option of purchasing combinations of the above connections

Community assignments:

IXP connection: community 100:2100

Private peer: community 100:2200

- Customer who buys local connectivity (via IXP) is put in community 100:2100
- Customer who buys peer connectivity is put in community 100:2200
- Customer who wants both IXP and peer connectivity is put in 100:2100 and 100:2200
- Customer who wants "the Internet" has no community set
 We are going to announce his prefix everywhere



Communities set at the aggregation router where the prefix is injected into the ISP's iBGP

- No need to alter filters at the network border when adding a new customer
- New customer simply is added to the appropriate community

Border filters already in place take care of announcements

⇒ Ease of operation!

Community Example – Internet Edge

- This demonstrates how communities might be used at the peering edge of an ISP network
- ISP has four types of BGP peers:

Customer

IXP peer

Private peer

Transit provider

- The prefixes received from each can be classified using communities
- Customers can opt to receive any or all of the above

Community Example – Internet Edge

Community assignments:

Customer prefix: community 100:3000

IXP prefix: community 100:3100

Private peer prefix: community 100:3200

- BGP customer who buys local connectivity gets 100:3000
- BGP customer who buys local and IXP connectivity receives community 100:3000 and 100:3100
- BGP customer who buys full peer connectivity receives community 100:3000, 100:3100, and 100:3200
- Customer who wants "the Internet" gets everything

Gets default route originated by aggregation router

Or pays money to get all 190k prefixes

Community Example – Internet Edge

No need to create customised filters when adding customers

Border router already sets communities

Installation engineers pick the appropriate community set when establishing the customer BGP session

⇒ Ease of operation!

Community Example – Summary

- Two examples of customer edge and internet edge can be combined to form a simple community solution for ISP prefix policy control
- More experienced operators tend to have more sophisticated options available

Advice is to start with the easy examples given, and then proceed onwards as experience is gained

Some ISP Examples

- ISPs also create communities to give customers bigger routing policy control
- Public policy is usually listed in the IRR

Following examples are all in the IRR

Examples build on the configuration concepts from the introductory example

Consider creating communities to give policy control to customers

Reduces technical support burden

Reduces the amount of router reconfiguration, and the chance of mistakes

Some ISP Examples: Sprintlink





http://www.sprintlink.net/policy/bgp.html











WHAT YOU CAN CONTROL

AS-PATH PREPENDS

Sprint allows customers to use AS-path prepending to adjust route preference on the network. Such prepending will be received and passed on properly without notifiving Sprint of your change in announcments.

Additionally, Sprint will prepend AS1239 to eBGP sessions with certain autonomous systems depending on a received community. Currently, the following ASes are supported: 1668, 209, 2914, 3300, 3356, 3549, 3561, 4635, 701, 7018, 702 and 8220.

String	Resulting AS Path to ASXXX			
65000:XXX	Do not advertise to ASXXX			
65001:XXX	1239 (default)			
65002:XXX	1239 1239			
65003:XXX	1239 1239 1239			
65004:XXX	1239 1239 1239 1239			
String	Resulting AS Path to ASXXX in Asia			
65070:XXX	Do not advertise to ASXXX			
65071:XXX	1239 (default)			
65072:XXX	1239 1239			
65073:XXX	1239 1239 1239			
65074:XXX	1239 1239 1239 1239			
String F	Resulting AS Path to ASXXX in Europe			
65050:XXX	Do not advertise to ASXXX			
65051:XXX	1239 (default)			
65052:XXX	1239 1239			
65053:XXX	1239 1239 1239			
65054:XXX	1239 1239 1239 1239			
04	Resulting AS Path to ASXXX in North			
String	America			
65010:XXX	Do not advertise to ASXXX			
65011:XXX	1239 (default)			
65012:XXX	1239 1239			
65013:XXX	1239 1239 1239			
65014:XXX	1239 1239 1239 1239			
String Re	esulting AS Path to all supported ASes			
65000:0	Do not advertise			
65001:0	1239 (default)			
65002:0	1239 1239			

4000 4000 4000

More info at

www.sprintlink.net/policy/bgp.html

Some ISP Examples AAPT

- Australian ISP
- Run their own Routing Registry

Whois.connect.com.au

 Offer 6 different communities to customers to aid with their traffic engineering

Some ISP Examples AAPT

```
AS2764
aut-num:
              ASN-CONNECT-NET
as-name:
descr:
             AAPT Limited
admin-c:
              CNO2-AP
tech-c:
              CNO2-AP
              Community support definitions
remarks:
remarks:
              Community Definition
remarks:
remarks:
              2764:2 Don't announce outside local POP
remarks:
remarks:
              2764:4 Lower local preference by 15
              2764:5 Lower local preference by 5
remarks:
remarks:
              2764:6 Announce to customers and all peers
                           (incl int'l peers), but not transit
remarks:
              2764:7 Announce to customers only
remarks:
              2764:14 Announce to AANX
notify:
              routing@connect.com.au
mnt-by:
              CONNECT-AU
changed:
              nobody@connect.com.au 20050225
              CCAIR
source:
```

More at http://info.connect.com.au/docs/routing/general/multi-faq.shtml#q13

Some ISP Examples MCI Europe

- MCI's European operation
- Permits customers to send communities which determine
 - local preferences within MCI's network
 - Reachability of the prefix
 - How the prefix is announced outside of MCI's network

Some ISP Examples MCI Europe

```
aut-num: AS702
descr: MCI EMEA - Commercial IP service provider in Europe
remarks: MCI uses the following communities with its customers:
        702:80
                  Set Local Pref 80 within AS702
        702:120
                  Set Local Pref 120 within AS702
        702:20
                 Announce only to MCI AS'es and MCI customers
        702:30
                  Keep within Europe, don't announce to other MCI AS's
        702:1
                  Prepend AS702 once at edges of MCI to Peers
        702:2
                  Prepend AS702 twice at edges of MCI to Peers
        702:3
                  Prepend AS702 thrice at edges of MCI to Peers
        Advanced communities for customers
        702:7020 Do not announce to AS702 peers with a scope of
                  National but advertise to Global Peers, European
                  Peers and MCI customers.
         702:7001 Prepend AS702 once at edges of MCI to AS702
                  peers with a scope of National.
        702:7002 Prepend AS702 twice at edges of MCI to AS702
                  peers with a scope of National.
(more)
```

Some ISP Examples MCI Europe

(more) 702:7003 Prepend AS702 thrice at edges of MCI to AS702 peers with a scope of National. 702:8020 Do not announce to AS702 peers with a scope of European but advertise to Global Peers, National Peers and MCI customers. 702:8001 Prepend AS702 once at edges of MCI to AS702 peers with a scope of European. 702:8002 Prepend AS702 twice at edges of MCI to AS702 peers with a scope of European. 702:8003 Prepend AS702 thrice at edges of MCI to AS702 peers with a scope of European. Additional details of the MCI communities are located at: http://global.mci.com/uk/customer/bgp/ mnt-by: WCOM-EMEA-RICE-MNT changed: rice@lists.mci.com 20040523 source: RIPE

Some ISP Examples BT Ignite

One of the most comprehensive community lists around

Seems to be based on definitions originally used in Tiscali's network

whois -h whois.ripe.net AS5400 reveals all

 Extensive community definitions allow sophisticated traffic engineering by customers

Some ISP Examples BT Ignite

aut-num:	AS5400	
descr:	BT Ignite European Backbone	
remarks:		
remarks:	Community to	Community to
remarks:	Not announce To peer:	AS prepend 5400
remarks:		
remarks:	5400:1000 All peers & Transits	5400:2000
remarks:		
remarks:	5400:1500 All Transits	5400:2500
remarks:	5400:1501 Sprint Transit (AS1239)	5400:2501
remarks:	5400:1502 SAVVIS Transit (AS3561)	5400:2502
remarks:	5400:1503 Level 3 Transit (AS3356)	5400:2503
remarks:	5400:1504 AT&T Transit (AS7018)	5400:2504
remarks:	5400:1505 UUnet Transit (AS701)	5400:2505
remarks:		
remarks:	5400:1001 Nexica (AS24592)	5400:2001
remarks:	5400:1002 Fujitsu (AS3324)	5400:2002
remarks:	5400:1004 C&W EU (1273)	5400:2004
<snip></snip>		
notify:	notify@eu.bt.net	any
mnt-by:	CTP-MNT	
source:	RIPE many n	nore!

Some ISP Examples Carrier1

- European ISP
- Another very comprehensive list of community definitions

whois -h whois.ripe.net AS8918 reveals all

Some ISP Examples Carrier1

```
aut-num:
          AS8918
          Carrier1 Autonomous System
descr:
<snip>
remarks:
          Community Definition
remarks:
          8918:2000 Do not announce to C1 customers
remarks:
remarks:
          8918:2010 Do not announce to C1 peers, peers+ and transit
          8918:2015 Do not announce to C1 transit providers
remarks:
remarks:
        8918:2020 Do not announce to Teleglobe (AS 6453)
remarks:
remarks: 8918:2035 Do not announce to UUNet (AS 702)
remarks:
          8918:2050 Do not announce to T-Systems (AS 3320)
remarks:
          8918:2070 Do not announce to AMS-IX peers
remarks:
          8918:2080 Do not announce to NL-IX peers
remarks:
<snip>
          inoc@carrier1.net
notify:
mnt-by:
          CARRIER1-MNT
source:
          RIPE
                                                 And many
                                                many more!
```

Some ISP Examples Level 3

- Highly detailed AS object held on the RIPE Routing Registry
- Also a very comprehensive list of community definitions

whois -h whois.ripe.net AS3356 reveals all

Some ISP Examples Level 3

```
AS3356
aut-num:
             Level 3 Communications
descr:
<snip>
remarks:
remarks:
             customer traffic engineering communities - Suppression
remarks.
             64960:XXX - announce to AS XXX if 65000:0
remarks:
remarks:
             65000:0 - announce to customers but not to peers
remarks:
             65000:XXX - do not announce at peerings to AS XXX
remarks:
remarks:
             customer traffic engineering communities - Prepending
remarks:
             65001:0 - prepend once to all peers
remarks:
remarks:
             65001:XXX - prepend once at peerings to AS XXX
<snip>
remarks:
             3356:70 - set local preference to 70
remarks:
             3356:80 - set local preference to 80
             3356:90 - set local preference to 90
remarks:
remarks:
             3356:9999 - blackhole (discard) traffic
<snip>
mnt-by:
             LEVEL3-MNT
                                                   And many
source:
             RIPE
                                                  many more!
```

BGP for Internet Service Providers

- BGP Basics
- Scaling BGP
- Using Communities
- Deploying BGP in an ISP network



Deploying BGP in an ISP Network

Okay, so we've learned all about BGP now; how do we use it on our network??

Deploying BGP

- The role of IGPs and iBGP
- Aggregation
- Receiving Prefixes
- Configuration Tips



The role of IGP and iBGP

Ships in the night?

Or

Good foundations?

BGP versus OSPF/ISIS

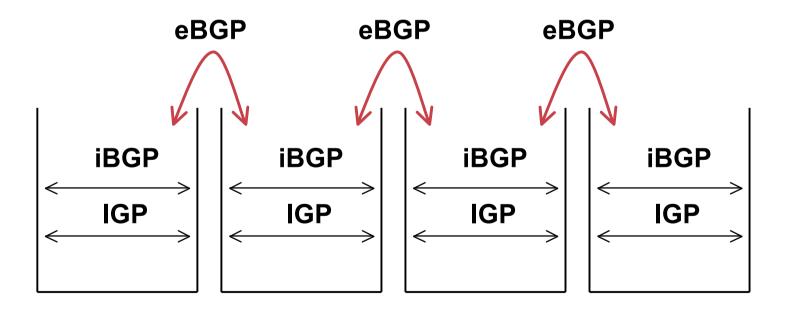
- Internal Routing Protocols (IGPs)
 - examples are ISIS and OSPF
 - used for carrying infrastructure addresses
 - **NOT** used for carrying Internet prefixes or customer prefixes
 - design goal is to minimise number of prefixes in IGP to aid scalability and rapid convergence

BGP versus OSPF/ISIS

- BGP used internally (iBGP) and externally (eBGP)
- iBGP used to carry some/all Internet prefixes across backbone customer prefixes
- eBGP used to
 exchange prefixes with other ASes
 implement routing policy

BGP/IGP model used in ISP networks

Model representation



BGP versus OSPF/ISIS

• DO NOT:

distribute BGP prefixes into an IGP distribute IGP routes into BGP use an IGP to carry customer prefixes

YOUR NETWORK WILL NOT SCALE

Injecting prefixes into iBGP

- Use iBGP to carry customer prefixes don't ever use IGP
- Point static route to customer interface
- Enter network into BGP process

Ensure that implementation options are used so that the prefix always remains in iBGP, regardless of state of interface

i.e. avoid iBGP flaps caused by interface flaps



Aggregation

Quality or Quantity?

Aggregation

- Aggregation means announcing the address block received from the RIR to the other ASes connected to your network
- Subprefixes of this aggregate may be:
 - **Used internally in the ISP network**
 - Announced to other ASes to aid with multihoming
- Unfortunately too many people are still thinking about class Cs, resulting in a proliferation of /24s in the Internet routing table

Aggregation

- Address block should be announced to the Internet as an aggregate
- Subprefixes of address block should NOT be announced to Internet unless special circumstances (more later)
- Aggregate should be generated internally

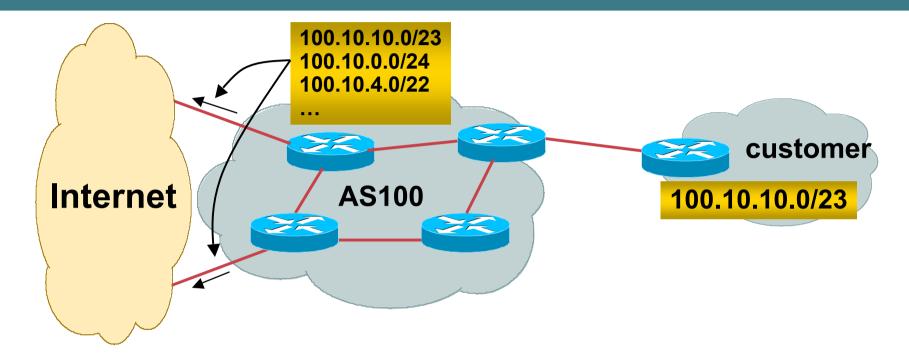
Not on the network borders!

Announcing an Aggregate

- ISPs who don't and won't aggregate are held in poor regard by community
- Registries publish their minimum allocation size Anything from a /20 to a /22 depending on RIR
- No real reason to see anything longer than a /22 prefix in the Internet

BUT there are currently >113000 /24s!

Aggregation – Example



- Customer has /23 network assigned from AS100's /19 address block
- AS100 announces customers' individual networks to the Internet

Aggregation – Bad Example

Customer link goes down

Their /23 network becomes unreachable

/23 is withdrawn from AS100's iBGP

 Their ISP doesn't aggregate its /19 network block

/23 network withdrawal announced to peers

starts rippling through the Internet

added load on all Internet backbone routers as network is removed from routing table

Customer link returns

Their /23 network is now visible to their ISP

Their /23 network is re-advertised to peers

Starts rippling through Internet

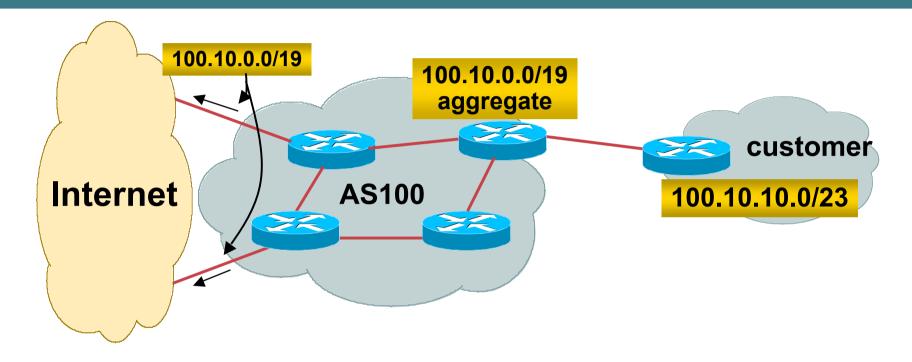
Load on Internet backbone routers as network is reinserted into routing table

Some ISP's suppress the flaps

Internet may take 10-20 min or longer to be visible

Where is the Quality of Service???

Aggregation – Example



- Customer has /23 network assigned from AS100's /19 address block
- AS100 announced /19 aggregate to the Internet

Aggregation – Good Example

- Customer link goes down their /23 network becomes unreachable /23 is withdrawn from AS100's **iBGP**
- announced no BGP hold down problems no BGP propagation delays no damping by other ISPs

/19 aggregate is still being

- **Customer link returns**
 - Their /23 network is visible again

The /23 is re-injected into AS100's iBGP

- The whole Internet becomes visible immediately
- **Customer has Quality of** Service perception

Aggregation – Summary

Good example is what everyone should do!

Adds to Internet stability

Reduces size of routing table

Reduces routing churn

Improves Internet QoS for everyone

Bad example is what too many still do!

Why? Lack of knowledge?

Laziness?

The Internet Today (February 2007)

Current Internet Routing Table Statistics

BGP Routing Table Entries	213110
Prefixes after maximum aggregation	114382
Unique prefixes in Internet	103747
Prefixes smaller than registry alloc	110121
/24s announced	113112
only 5746 /24s are from 192.0.0.0/8	
ASes in use	24456

"The New Swamp"

Swamp space is name used for areas of poor aggregation

The original swamp was 192.0.0.0/8 from the former class C block

Name given just after the deployment of CIDR

The new swamp is creeping across all parts of the Internet Not just RIR space, but "legacy" space too

"The New Swamp" RIR Space – February 1999

RIR blocks contribute 49393 prefixes or 88% of the Internet Routing Table

Block	Networks	Block	Networks	Block	Networks	Block	Networks
24/8	165	77/8	0	118/8	0	203/8	3622
41/8	0	78/8	0	119/8	0	204/8	3792
58/8	0	79/8	0	120/8	0	205/8	2584
59/8	0	80/8	0	121/8	0	206/8	3127
60/8	0	81/8	0	122/8	0	207/8	2723
61/8	3	82/8	0	123/8	0	208/8	2817
62/8	87	83/8	0	124/8	0	209/8	2574
63/8	20	84/8	0	125/8	0	210/8	617
64/8	0	<i>85/8</i>	0	126/8	0	211/8	0
65/8	0	86/8	0	189/8	0	212/8	717
66/8	0	87/8	0	190/8	0	213/8	1
67/8	0	88/8	0	192/8	6275	216/8	943
68/8	0	89/8	0	193/8	2390	217/8	0
69/8	0	90/8	0	194/8	2932	218/8	0
70/8	0	91/8	0	195/8	1338	219/8	0
71/8	0	96/8	0	196/8	513	220/8	0
72/8	0	97/8	0	198/8	4034	221/8	0
73/8	0	98/8	0	199/8	3495	222/8	0
74/8	0	99/8	0	200/8	1348		
75/8	0	116/8	0	201/8	0		
76/8	0	117/8	0	202/8	2276		Cioco Public

"The New Swamp" RIR Space – February 2007

RIR blocks contribute 192490 prefixes or 90% of the Internet Routing Table

Block	Networks	Block	Networks	Block	Networks	Block	Networks
24/8	2930	77/8	1214	118/8	3	203/8	10459
41/8	288	78/8	8	119/8	3	204/8	5569
58/8	1097	79/8	2	120/8	3	205/8	2892
59/8	1152	80/8	2053	121/8	426	206/8	3857
60/8	604	81/8	1695	122/8	698	207/8	4331
61/8	2589	82/8	1564	123/8	534	208/8	4258
62/8	2193	83/8	1172	124/8	1340	209/8	5540
63/8	2967	84/8	1269	125/8	1554	210/8	4759
64/8	5501	85/8	1891	126/8	41	211/8	2733
65/8	3917	86/8	800	189/8	169	212/8	2900
66/8	6575	87/8	1157	190/8	1077	213/8	3052
67/8	2015	88/8	847	192/8	6927	216/8	6930
68/8	2770	89/8	1970	193/8	5704	217/8	2615
69/8	3702	90/8	105	194/8	4652	218/8	1561
70/8	1693	91/8	577	195/8	4279	219/8	1197
71/8	1188	96/8	8	196/8	1600	220/8	1988
72/8	2878	97/8	1	198/8	4748	221/8	894
73/8	273	98/8	3	199/8	4184	222/8	1241
74/8	1483	99/8	0	200/8	7482		
75/8	483	116/8	3 3	201/8	2927		
76/8	194	117/8	3	202/8	10529		Cioco Dublio

"The New Swamp" Summary

RIR space shows creeping deaggregation

It seems that an RIR /8 block averages around 5000 prefixes once fully allocated

So their existing 80 /8s will eventually cause 400000 prefix announcements

Food for thought:

Remaining 54 unallocated /8s and the 80 RIR /8s combined will cause:

600000 prefixes with 5000 prefixes per /8 density

804000 prefixes with 6000 prefixes per /8 density

Plus 12% due to "non RIR space deaggregation"

→ Routing Table size of 900480 prefixes

"The New Swamp" Summary

- Rest of address space is showing similar deaggregation too
- What are the reasons?

Main justification is traffic engineering

Real reasons are:

Lack of knowledge

Laziness

Deliberate & knowing actions

BGP Report (bgp.potaroo.net)

- 199336 total announcements in October 2006
- 129795 prefixes

After aggregating including full AS PATH info

i.e. including each ASN's traffic engineering

35% saving possible

109034 prefixes

After aggregating by Origin AS

i.e. ignoring each ASN's traffic engineering

10% saving possible

The excuses

- Traffic engineering causes 10% of the Internet Routing table
- Deliberate deaggregation causes 35% of the Internet Routing table

Efforts to improve aggregation

The CIDR Report

Initiated and operated for many years by Tony Bates

Now combined with Geoff Huston's routing analysis

www.cidr-report.org

Results e-mailed on a weekly basis to most operations lists around the world

Lists the top 30 service providers who could do better at aggregating

Efforts to improve aggregation The CIDR Report

- Also computes the size of the routing table assuming ISPs performed optimal aggregation
- Website allows searches and computations of aggregation to be made on a per AS basis

Flexible and powerful tool to aid ISPs

Intended to show how greater efficiency in terms of BGP table size can be obtained without loss of routing and policy information

Shows what forms of origin AS aggregation could be performed and the potential benefit of such actions to the total table size

Very effectively challenges the traffic engineering excuse











Status Summary

Table History

Date	Prefixes	CIDR Aggregated
02-05-0	5 157356	108023
03-05-0	5 157392	108044
04-05-0	5 157505	108133
05-05-0	5 157530	108201
06-05-0	5 157716	108341
07-05-0	5 157747	108272
08-05-0	5 157845	108355
09-05-0	5 157874	108388



Plot: BGP Table Size

AS Summary

19498 Number of ASes in routing system

7996 Number of ASes announcing only one prefix

1467 Largest number of prefixes announced by an AS

AS7018: ATT-INTERNET4 - AT&T WorldNet

Services

90497280 Largest address span announced by an AS (/32s)

AS721: DLA-ASNBLOCK-AS - DoD Network

Information Center



Plot: AS count

Plot: Average announcements per origin AS

Report: ASes ordered by originating address span Report: ASes ordered by transit address span

Report: Autonomous System number-to-name mapping (from Registry WHOIS data)













Aggregation Summary

The algorithm used in this report proposes aggregation only when there is a precise match using AS path so as to preserve traffic transit policies. Aggregation is also proposed across non-advertised address space ('holes').

09Ma	y05				
ASnum	NetsNow	NetsAggr	NetGain	% Gain	Description

Table	157925	108381	49544	31.4%	All ASes
AS4323	1098	223	875	79.7%	TWTC - Time Warner Telecom
AS18566	805	8	797	99.0%	COVAD - Covad Communications
AS4134	893	220	673	75.4%	CHINANET-BACKBONE No.31, Jin-rong Street
AS721	1117	564	553	49.5%	DLA-ASNBLOCK-AS - DoD Network Information Center
AS7018	1467	939	528	36.0%	ATT-INTERNET4 - AT&T WorldNet Services
AS27364	539	22	517	95.9%	ACS-INTERNET - Armstrong Cable Services
AS22773	483	23	460	95.2%	CCINET-2 - Cox Communications Inc.
AS6197	900	506	394	43.8%	BATI-ATL - BellSouth Network Solutions, Inc
AS3602	509	146	363	71.3%	SPRINT-CA-AS - Sprint Canada Inc.
AS17676	431	78	353	81.9%	JPNIC-JP-ASN-BLOCK Japan Network Information Center
AS9929	350	46	304	86.9%	CNCNET-CN China Netcom Corp.
AS4766	574	279	295		KIXS-AS-KR Korea Telecom
AS6478	416	123	293	70.4%	ATT-INTERNET3 - AT&T WorldNet Services
AS6140	399	135	264	66.2%	IMPSAT-USA - ImpSat
AS14654	264	6	258		WAYPORT - Wayport
AS9583	735	483	252	34.3%	SIFY-AS-IN Sify Limited
AS9443	374	123	251	67.1%	INTERNETPRIMUS-AS-AP Primus Telecommunications
AS7545	493	247	246	49.9%	TPG-INTERNET-AP TPG Internet Pty Ltd
AS1239	886	644	242	27.3%	SPRINTLINK - Sprint
AS15270	272	37	235	86.4%	AS-PAETEC-NET - PaeTec.net -a division of PaeTecCommunications, Inc.
AS23126	254	23	231		KMCTELCOM-DIA - KMC Telecom, Inc.
AS4755	516	287	229		VSNL-AS Videsh Sanchar Nigam Ltd. Autonomous System
AS7725	415	186	229		CCH-AS7 - Comcast Cable Communications Holdings, Inc
AS6198	464	236	228		BATI-MIA - BellSouth Network Solutions, Inc
AS5668	488	264	224		AS-5668 - CenturyTel Internet Holdings, Inc.
AS2386	853	634	219		INS-AS - AT&T Data Communications Services
AS9498	296	79	217		BBIL-AP BHARTI BT INTERNET LTD.
AS11456	319	110	209		NUVOX - NuVox Communications, Inc.
AS6167	264	67	197		CELLCO-PART - Cellco Partnership
AS6517	319	128	191	59.9%	YIPESCOM - Yipes Communications, Inc.
Total	17193	6866	10327	AN 12	Ton 30 total











Top 20 Added Routes this week per Originating AS

Nttp://www.cidr-report.org/

Prefixes ASnum AS Description 154 AS7725 CCH-AS7 - Comcast Cable Communications Holdings, Inc. 108 AS4755 VSNL-AS Videsh Sanchar Nigam Ltd. Autonomous System 52 AS35911 BNO-1 - Telebec 36 AS13645 BROADBANDONE - BroadbandONE, Inc. 19 AS17488 HATHWAY-NET-AP Hathway IP Over Cable Internet 16 AS9576 SOOKMYUNG-AS SOOKMYUNG WOMEN'S UNIVERSITY 16 AS174 COGENT Cogent/PSI 15 AS18633 GIANTWEB - Giant Technologies Inc. 16 AS18042 KBT Koos Broadband Telecom 16 AS32613 IWEB-AS - Groupe iWeb Technologies inc. 15 AS19632 Metropolis Intercom 15 AS30340 AS-LLIX - Liberty Lake Internet Portal 13 AS19916 ASTRUM-0001 - OLM LLC 13 AS22047 VTR BANDA ANCHA S.A. 13 AS21882 PRIORITYNETWORKS - Priority Networks Inc. 42 AS9940 WOLCST-AS-AP World online AS, Cybersoft Technologies. 12 AS12715 JAZZNET Jazz Telecom S.A. 12 AS22927 Telefonica de Argentina 11 AS30533 CONNEXION-BY-BOEING-LTN - Connexion by Boeing 11 AS25454 TELEMEDIAAS Telemedia SA Autonomous System

Top 20 Withdrawn Routes this week per Originating AS

Deofivos	4 Coum	AS Description
		LH - Lighthouse Communications, Inc.
		WEBCENTRAL-AS WebCentral
-31	AS8921	I-CONNEXION ICX Autonomous System
-23	AS4513	Globix Corporation
-20	AS1239	SPRINTLINK - Sprint
		ACDNET-ASN1 - ACD.net
-17	AS29257	CBB-IE-AS Connexion by Boeing Ireland, Ltd.
		CHARTER-NET-HKY-NC - Charter Communications
-16	AS6167	CELLCO-PART - Cellco Partnership
		PKTELECOM-AS-AP Pakistan Telecom
-14	AS9152	MEGADAT Autonomous System
-14	AS16154	TELECOMS-AS Telecoms-Net Ltd.
-14	AS24219	NFI-AS-AP No Fuss Internet
-13	AS174	COGENT Cogent/PSI
-13	AS10125	DACCESS-AP DATA ACCESS INDIA LIMITED
-13	AS30857	TAURUS-AS Taurus Telecom PJSC
-12	AS17854	CABLELINE-AS-KR BANDOCABLELINE
-12	AS7049	S&M International S.A.
-12	AS4323	TWTC - Time Warner Telecom
-12	AS3561	SAWIS - Savvis







Report: Announced Route count per Originating AS Report: Withdrawn Route count per Originating AS

More Specifics

A list of route advertisements that appear to be more specific than the original Class-based prefix mask, or more specific than the registry allocation size.

Top 20 ASes advertising more specific prefixes

More Specifics	Total Prefixes	ASnum	AS Description
1103	1467	AS7018	ATT-INTERNET4 - AT&T WorldNet Services
1012	1180	AS174	COGENT Cogent/PSI
974	1098	AS4323	TWTC - Time Warner Telecom
880	900	AS6197	BATI-ATL - BellSouth Network Solutions, Inc
801	1117	AS721	DLA-ASNBLOCK-AS - DoD Network Information Center
798	805	AS18566	COVAD - Covad Communications
780	853	AS2386	INS-AS - AT&T Data Communications Services
742	893	AS4134	CHINANET-BACKBONE No.31, Jin-rong Street
730	735	AS9583	SIFY-AS-IN Sify Limited
621		AS1239	SPRINTLINK - Sprint
594	994	AS701	ALTERNET-AS - UUNET Technologies, Inc.
583	595	AS20115	CHARTER-NET-HKY-NC - Charter Communications
540	574	AS4766	KIXS-AS-KR Korea Telecom
533	539	AS27364	ACS-INTERNET - Armstrong Cable Services
500	516	AS4755	VSNL-AS Videsh Sanchar Nigam Ltd. Autonomous System
475	488	AS5668	AS-5668 - CenturyTel Internet Holdings, Inc.
470	483	AS22773	CCINET-2 - Cox Communications Inc.
456	493	AS7545	TPG-INTERNET-AP TPG Internet Pty Ltd
453		AS3602	SPRINT-CA-AS - Sprint Canada Inc.
452	464	AS6198	BATI-MIA - BellSouth Network Solutions, Inc

Report: ASes ordered by number of more specific prefixes

Report: More Specific prefix list (by AS)

Report: More Specific prefix list (ordered by prefix)









Rank AS Type Originate Addr Space (pfx) Transit Addr space (pfx) Description
24 AS1239 ORG+TRN Originate: 11982080 /8.49 Transit: 145498112 /4.88 SPRINTLINK - Sprint

Aggregation Suggestions

This report does not take into account conditions local to each origin AS in terms of policy or traffic engineering requirements, so this is an approximate guideline as to aggregation possibilities.

Rank AS	AS Name	Current	Wthdw	Aggte	Annce	Redotn	*
20 AS1239	SPRINTLINK - Sprint	886	307	65	644	242	27.31%

Prefix (AS Path) Aggregation Action	
12.9.182.0/23 4637 1239	
12.22.206.0/24 4637 1239	
24.56.144.0/21 4637 1239	
24.137.128.0/21 4637 1239	
24.221.0.0/17 4637 1239 + Announce - aggregate of 24.221.0.0/18 (4637 1239) and 24.221	1.64.0/18 (4637 1239)
24.221.0.0/18 4637 1239 - Withdrawn - aggregated with 24.221.64.0/18 (4637 1239)	
24.221.64.0/19 4637 1239 - Withdrawn - aggregated with 24.221.96.0/19 (4637 1239)	
24.221.96.0/19 4637 1239 - Withdrawn - aggregated with 24.221.64.0/19 (4637 1239)	
24.221.128.0/18 4637 1239 + Announce - aggregate of 24.221.128.0/19 (4637 1239) and 24.2	221.160.0/19 (4637 1239)
24.221.128.0/19 4637 1239 - Withdrawn - aggregated with 24.221.160.0/19 (4637 1239)	
24.221.160.0/19 4637 1239 - Withdrawn - aggregated with 24.221.128.0/19 (4637 1239)	
24.221.192.0/20 4637 1239	
24.221.220.0/22 4637 1239	
24.221.224.0/20 4637 1239 + Announce - aggregate of 24.221.224.0/21 (4637 1239) and 24.2	221.232.0/21 (4637 1239)
24.221.224.0/21 4637 1239 - Withdrawn - aggregated with 24.221.232.0/21 (4637 1239)	
24.221.232.0/22 4637 1239 - Withdrawn - aggregated with 24.221.236.0/22 (4637 1239)	
24.221.236.0/22 4637 1239 - Withdrawn - aggregated with 24.221.232.0/22 (4637 1239)	
24.221.242.0/23 4637 1239	
24.221.244.0/22 4637 1239	
24.221.248.0/21 4637 1239	
38.113.4.0/24 4637 1239	
63.90.4.0/24 4637 1239	
63.113.210.0/24 4637 1239	
63.122.77.0/24 4637 1239	
63.122.78.0/23 4637 1239	
63.134.0.0/17 4637 1239	
63.160.0.0/12 4637 1239	
63.178.251.0/24 4637 1239	
63.237.89.0/24 4637 1239	
64.6.224.0/19 4637 1239	
64.9.45.0/24 4637 1239	
64.9.86.0/24 4637 1239	
64.17.64.0/22 4637 1239	



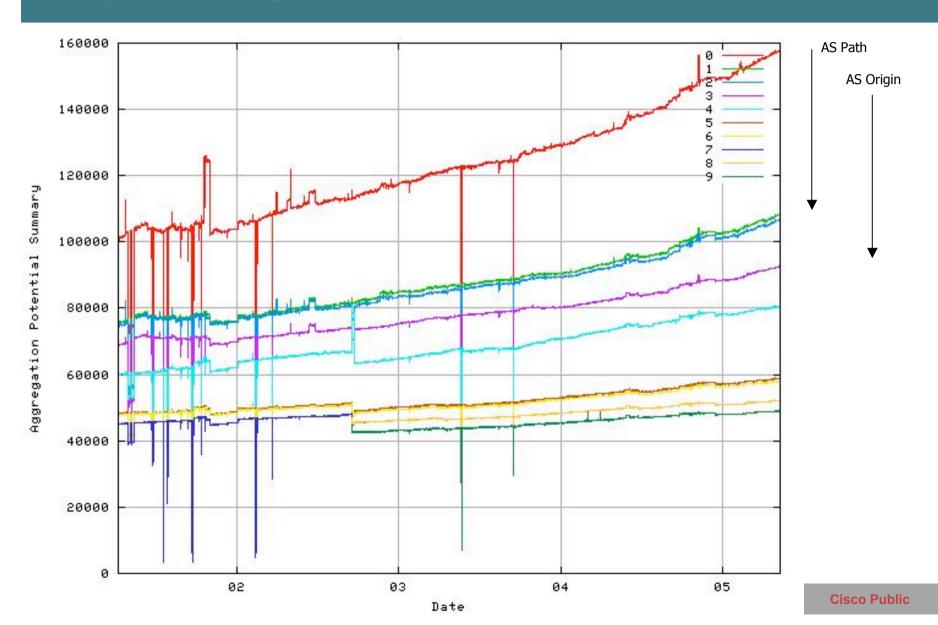




Rank AS AS Name Current Wthdw Aggte Annce Redctn % 49 AS701 ALTERNET-AS - UUNET Technologies, Inc. 994 208 68 854 140 14.08%

```
AS 701: ALTERNET-AS - UUNET Technologies, Inc.
  Prefix (AS Path)
                                       Aggregation Action
17.255.232.0/24
                     4637 701
24.32.66.0/24
                     4637 701
24.32.68.0/22
                     4637 701
                                           + Announce - aggregate of 24.32.68.0/23 (4637 701) and 24.32.70.0/23 (4637 701)
24.32.68.0/24
                     4637 701
                                           - Withdrawn - aggregated with 24.32.69.0/24 (4637 701)
24.32.69.0/24
                     4637 701
                                           - Withdrawn - aggregated with 24.32.68.0/24 (4637 701)
24.32.70.0/24
                                           - Withdrawn - aggregated with 24.32.71.0/24 (4637 701)
                     4637 701
24.32.71.0/24
                     4637 701
                                           - Withdrawn - aggregated with 24.32.70.0/24 (4637 701)
24.32.130.0/24
                     4637 701
24.32.144.0/22
                     4637 701
                                           + Announce - aggregate of 24.32.144.0/23 (4637 701) and 24.32.146.0/23 (4637 701)
24.32.144.0/23
                     4637 701
                                           - Withdrawn - aggregated with 24.32.146.0/23 (4637 701)
                     4637 701
                                           - Withdrawn - aggregated with 24.32.144.0/23 (4637 701)
24.32.146.0/23
24.32.163.0/24
                     4637 701
24.32.164.0/24
                     4637 701
24.206.172.0/24
                     4637 701
                     4637 701
24.216.0.0/16
                     4637 701
                                           - Withdrawn - matching aggregate 24.216.0.0/16 4637 701
24.216.82.0/24
24.216.94.0/23
                     4637 701
                                           - Withdrawn - matching aggregate 24.216.0.0/16 4637 701
24.216.174.0/24
                     4637 701
24.240.0.0/15
                     4637 701
55.191.7.0/24
                     4637 701
62.70.23.0/24
                     4637 701
63.0.0.0/9
                     4637 701
                                           + Announce - aggregate of 63.0.0.0/10 (4637 701) and 63.64.0.0/10 (4637 701)
63.0.0.0/12
                     4637 701
                                           - Withdrawn - aggregated with 63.16.0.0/12 (4637 701)
63.16.0.0/12
                     4637 701
                                           - Withdrawn - aggregated with 63.0.0.0/12 (4637 701)
63.32.0.0/12
                                           - Withdrawn - aggregated with 63.48.0.0/12 (4637 701)
                     4637 701
63.48.0.0/12
                     4637 701
                                           - Withdrawn - aggregated with 63.32.0.0/12 (4637 701)
63.64.0.0/12
                     4637 701
                                           - Withdrawn - aggregated with 63.80.0.0/12 (4637 701)
63.80.0.0/12
                     4637 701
                                           - Withdrawn - aggregated with 63.64.0.0/12 (4637 701)
63.96.0.0/12
                     4637 701
                                           - Withdrawn - aggregated with 63.112.0.0/12 (4637 701)
63.112.0.0/12
                     4637 701
                                           - Withdrawn - aggregated with 63.96.0.0/12 (4637 701)
63.134.153.0/24
                     4637 701
63.134.154.0/24
                     4637 701
63.134.161.0/24
                     4637 701
63.134.162.0/23
                     4637 701
                                           + Announce - aggregate of 63.134.162.0/24 (4637 701) and 63.134.163.0/24 (4637 701)
63.134.162.0/24
                     4637 701
                                           - Withdrawn - aggregated with 63.134.163.0/24 (4637 701)
63.134.163.0/24
                     4637 701
                                           - Withdrawn - aggregated with 63.134.162.0/24 (4637 701)
                     4637 701
63.134.164.0/24
63.134.168.0/23
                     4637 701
63.134.176.0/24
                     4637 701
                     4637 701
63.134.179.0/24
63.141.42.0/24
                     4637 701
```

Aggregation Potential (source: bgp.potaroo.net/as4637/)



Aggregation Summary

Aggregation on the Internet could be MUCH better

35% saving on Internet routing table size is quite feasible

Tools are available

Commands on the routers are not hard

CIDR-Report webpage

RIPE Routing WG aggregation recommendation

RIPE-399 — http://www.ripe.net/ripe/docs/ripe-399.html



Receiving Prefixes

Receiving Prefixes

 There are three scenarios for receiving prefixes from other ASNs

Customer talking BGP

Peer talking BGP

Upstream/Transit talking BGP

 Each has different filtering requirements and need to be considered separately

Receiving Prefixes: From Customers

- ISPs should only accept prefixes which have been assigned or allocated to their downstream customer
- If ISP has assigned address space to its customer, then the customer IS entitled to announce it back to his ISP
- If the ISP has NOT assigned address space to its customer, then:

Check in the four RIR databases to see if this address space really has been assigned to the customer

The tool: whois –h whois.apnic.net x.x.x.0/24

Receiving Prefixes: From Customers

 Example use of whois to check if customer is entitled to announce address space:

pfs-pc\$ whois -h whois.apnic.net 202.12.29.0

inetnum: 202.12.29.0 - 202.12.29.255

netname: APNIC-AP-AU-BNE

descr: APNIC Pty Ltd - Brisbane Offices + Servers

descr: Level 1, 33 Park Rd

descr: PO Box 2131, Milton

descr: Brisbane, QLD.

country: AU

admin-c: HM20-AP

tech-c: NO4-AP

mnt-by: APNIC-HM

changed: hm-changed@apnic.net 20030108

status: ASSIGNED PORTABLE

source: APNIC

Portable – means its an assignment

to the customer, the customer can

announce it to you

Receiving Prefixes: From Customers

 Example use of whois to check if customer is entitled to announce address space:

\$ whois -h whois.ripe.net 193.128.2.0

inetnum: 193.128.2.0 - 193.128.2.15

descr: Wood Mackenzie

country: GB

admin-c: DB635-RIPE tech-c: DB635-RIPE

status: ASSIGNED PA

mnt-by: AS1849-MNT

changed: davids@uk.uu.net 20020211

source: RIPE

route: 193.128.0.0/14

descr: PIPEX-BLOCK1

origin: AS1849

notify: routing@uk.uu.net

mnt-by: AS1849-MNT

changed: beny@uk.uu.net 20020321

source: RIPE

ASSIGNED PA - means that it is

to the ISP who assigned it

Provider Aggregatable address space

and can only be used for connecting

Receiving Prefixes: From Peers

 A peer is an ISP with whom you agree to exchange prefixes you originate into the Internet routing table

Prefixes you accept from a peer are only those they have indicated they will announce

Prefixes you announce to your peer are only those you have indicated you will announce

Receiving Prefixes: From Peers

Agreeing what each will announce to the other:

Exchange of e-mail documentation as part of the peering agreement, and then ongoing updates

OR

Use of the Internet Routing Registry and configuration tools such as the IRRToolSet

www.isc.org/sw/IRRToolSet/

Receiving Prefixes: From Upstream/Transit Provider

- Upstream/Transit Provider is an ISP who you pay to give you transit to the WHOLE Internet
- Receiving prefixes from them is not desirable unless really necessary

special circumstances – see later

Ask upstream/transit provider to either:

originate a default-route

OR

announce one prefix you can use as default

Receiving Prefixes: From Upstream/Transit Provider

 If necessary to receive prefixes from any provider, care is required

```
don't accept RFC1918 etc prefixes
```

ftp://ftp.rfc-editor.org/in-notes/rfc3330.txt

don't accept your own prefixes

don't accept default (unless you need it)

don't accept prefixes longer than /24

Check Rob Thomas' list of "bogons"

http://www.cymru.com/Documents/bogon-list.html

Receiving Prefixes

 Paying attention to prefixes received from customers, peers and transit providers assists with:

The integrity of the local network

The integrity of the Internet

Responsibility of all ISPs to be good Internet citizens



Preparing the network

Before we begin...

Preparing the Network

- We will deploy BGP across the network before we try and multihome
- BGP will be used therefore an ASN is required
- If multihoming to different ISPs, public ASN needed:

Either go to upstream ISP who is a registry member, or Apply to the RIR yourself for a one off assignment, or Ask an ISP who is a registry member, or

Join the RIR and get your own IP address allocation too (this option strongly recommended)!

Preparing the Network Initial Assumptions

- The network is not running any BGP at the moment single statically routed connection to upstream ISP
- The network is not running any IGP at all Static default and routes through the network to do "routing"

Preparing the Network Example One

- The network is not running any BGP at the moment single statically routed connection to upstream ISP
- The network is not running any IGP at all Static default and routes through the network to do "routing"

Preparing the Network First Step: IGP

- Decide on IGP: OSPF or ISIS ©
- Assign loopback interfaces and /32 addresses to each router which will run the IGP

Loopback is used for OSPF and BGP router id anchor Used for iBGP and route origination

Deploy IGP (e.g. OSPF)

IGP can be deployed with NO IMPACT on the existing static routing

e.g. OSPF distance might be 110, static distance is 1

Smallest distance wins

Preparing the Network IGP (cont)

 Be prudent deploying IGP – keep the Link State Database Lean!

Router loopbacks go in IGP

WAN point to point links go in IGP

(In fact, any link where IGP dynamic routing will be run should go into IGP)

Summarise on area/level boundaries (if possible) – i.e. think about your IGP address plan

Preparing the Network IGP (cont)

Routes which don't go into the IGP include:

Dynamic assignment pools (DSL/Cable/Dial)

Customer point to point link addressing

(using next-hop-self in iBGP ensures that these do NOT need to be in IGP)

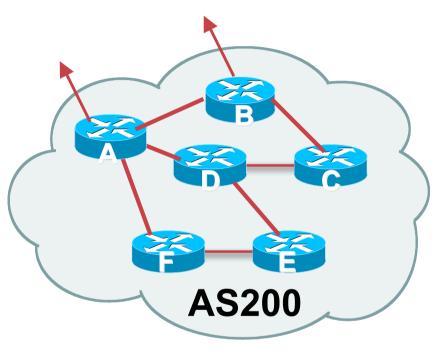
Static/Hosting LANs

Customer assigned address space

Anything else not listed in the previous slide

Preparing the Network Second Step: iBGP

- Second step is to configure the local network to use iBGP
- iBGP can run on
 all routers, or
 a subset of routers, or
 just on the upstream edge
- iBGP must run on all routers which are in the transit path between external connections



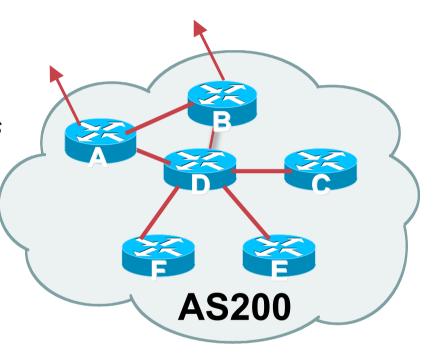
Preparing the Network Second Step: iBGP (Transit Path)

- iBGP must run on all routers which are in the transit path between external connections
- Routers C, E and F are not in the transit path

Static routes or IGP will suffice

Router D is in the transit path

Will need to be in iBGP mesh, otherwise routing loops will result



Preparing the Network Layers

- Typical SP networks have three layers:
 - Core the backbone, usually the transit path
 - Distribution the middle, PoP aggregation layer
 - Aggregation the edge, the devices connecting customers

Preparing the Network Aggregation Layer

iBGP is optional

Many ISPs run iBGP here, either partial routing (more common) or full routing (less common)

Full routing is not needed unless customers want full table

Partial routing is cheaper/easier, might usually consist of internal prefixes and, optionally, external prefixes to aid external load balancing

Communities and peer-groups make this administratively easy

Many aggregation devices can't run iBGP

Static routes from distribution devices for address pools IGP for best exit

Preparing the Network Distribution Layer

Usually runs iBGP

Partial or full routing (as with aggregation layer)

But does not have to run iBGP

IGP is then used to carry customer prefixes (does not scale)
IGP is used to determine nearest exit

 Networks which plan to grow large should deploy iBGP from day one

Migration at a later date is extra work

No extra overhead in deploying iBGP, indeed IGP benefits

Preparing the Network Core Layer

- Core of network is usually the transit path
- iBGP necessary between core devices

Full routes or partial routes:

Transit ISPs carry full routes in core

Edge ISPs carry partial routes only

Core layer includes AS border routers

Decide on:

Best iBGP policy

Will it be full routes everywhere, or partial, or some mix?

iBGP scaling technique

Community policy?

Route-reflectors?

Techniques such as peer groups and peer templates?

Then deploy iBGP:

Step 1: Introduce iBGP mesh on chosen routers make sure that iBGP distance is greater than IGP distance (it usually is)

Step 2: Install "customer" prefixes into iBGP

Check! Does the network still work?

Step 3: Carefully remove the static routing for the prefixes now in IGP and iBGP

Check! Does the network still work?

Step 4: Deployment of eBGP follows

Install "customer" prefixes into iBGP?

- Customer assigned address space
 Network statement/static route combination
 Use unique community to identify customer assignments
- Customer facing point-to-point links
 - Redistribute connected through filters which only permit point-topoint link addresses to enter iBGP
 - Use a unique community to identify point-to-point link addresses (these are only required for your monitoring system)
- Dynamic assignment pools & local LANs
 - Simple network statement will do this
 - Use unique community to identify these networks

Carefully remove static routes?

Work on one router at a time:

Check that static route for a particular destination is also learned by the iBGP

If so, remove it

If not, establish why and fix the problem

(Remember to look in the RIB, not the FIB!)

- Then the next router, until the whole PoP is done
- Then the next PoP, and so on until the network is now dependent on the IGP and iBGP you have deployed

Preparing the Network Completion

Previous steps are NOT flag day steps

Each can be carried out during different maintenance periods, for example:

Step One on Week One

Step Two on Week Two

Step Three on Week Three

And so on

And with proper planning will have NO customer visible impact at all

Preparing the Network Example Two

- The network is not running any BGP at the moment single statically routed connection to upstream ISP
- The network is running an IGP though All internal routing information is in the IGP By IGP, OSPF or ISIS is assumed

Preparing the Network IGP

 If not already done, assign loopback interfaces and /32 addresses to each router which is running the IGP

Loopback is used for OSPF and BGP router id anchor Used for iBGP and route origination

Ensure that the loopback /32s are appearing in the IGP

Preparing the Network iBGP

- Go through the iBGP decision process as in Example One
- Decide full or partial, and the extent of the iBGP reach in the network

Preparing the Network iBGP Implementation

Then deploy iBGP:

Step 1: Introduce iBGP mesh on chosen routers
make sure that iBGP distance is greater than IGP distance (it usually is)

Step 2: Install "customer" prefixes into iBGP

Check! Does the network still work?

Step 3: Reduce BGP distance to be less than the IGP (so that iBGP routes take priority)

Step 4: Carefully remove the "customer" prefixes from the IGP

Check! Does the network still work?

Step 5: Restore BGP distance to less than IGP

Step 6: Deployment of eBGP follows

Preparing the Network iBGP implementation

Install "customer" prefixes into iBGP?

Customer assigned address space

Network statement/static route combination

Use unique community to identify customer assignments

Customer facing point-to-point links

Redistribute connected through filters which only permit point-topoint link addresses to enter iBGP

Use a unique community to identify point-to-point link addresses (these are only required for your monitoring system)

Dynamic assignment pools & local LANs

Simple network statement will do this

Use unique community to identify these networks

Preparing the Network iBGP implementation

Carefully remove "customer" routes from IGP?

Work on one router at a time:

Check that IGP route for a particular destination is also learned by iBGP

If so, remove it from the IGP

If not, establish why and fix the problem

(Remember to look in the RIB, not the FIB!)

- Then the next router, until the whole PoP is done
- Then the next PoP, and so on until the network is now dependent on the iBGP you have deployed

Preparing the Network Completion

Previous steps are NOT flag day steps

Each can be carried out during different maintenance periods, for example:

Step One on Week One

Step Two on Week Two

Step Three on Week Three

And so on

And with proper planning will have NO customer visible impact at all

Preparing the Network Configuration Summary

- IGP essential networks are in IGP
- Customer networks are now in iBGP iBGP deployed over the backbone
 Full or Partial or Upstream Edge only
- BGP distance is greater than any IGP
- Now ready to deploy eBGP



Configuration Tips

Of templates, passwords, tricks, and more templates

iBGP and IGPs Reminder!

- Make sure loopback is configured on router iBGP between loopbacks, NOT real interfaces
- Make sure IGP carries loopback /32 address
- Consider the DMZ nets:
 - Use unnumbered interfaces?
 - Use next-hop-self on iBGP neighbours
 - Or carry the DMZ /30s in the iBGP
 - Basically keep the DMZ nets out of the IGP!

Next-hop-self

Used by many ISPs on edge routers

Preferable to carrying DMZ /30 addresses in the IGP

Reduces size of IGP to just core infrastructure

Alternative to using unnumbered interfaces

Helps scale network

BGP speaker announces external network using local address (loopback) as next-hop

Templates

Good practice to configure templates for everything

Vendor defaults tend not to be optimal or even very useful for ISPs

ISPs create their own defaults by using configuration templates

eBGP and iBGP examples follow

Also see Project Cymru's BGP templates

www.cymru.com/Documents

iBGP Template Example

- iBGP between loopbacks!
- Next-hop-self
 Keep DMZ and external point-to-point out of IGP
- Always send communities in iBGP
 Otherwise accidents will happen
- Hardwire BGP to version 4
 Yes, this is being paranoid!
- Use passwords on iBGP session
 Not being paranoid, VERY necessary

eBGP Template Example

BGP damping

Do NOT use it unless you understand the impact

Do NOT use the vendor defaults without thinking

- Remove private ASes from announcements Common omission today
- Use extensive filters, with "backup"

Use as-path filters to backup prefix filters

Keep policy language for implementing policy, rather than basic filtering

Use password agreed between you and peer on eBGP session

eBGP Template Example continued

Use maximum-prefix tracking

Router will warn you if there are sudden increases in BGP table size, bringing down eBGP if desired

Log changes of neighbour state

...and monitor those logs!

Make BGP admin distance higher than that of any IGP

Otherwise prefixes heard from outside your network could override your IGP!!

Limiting AS Path Length

 Some BGP implementations have problems with long AS_PATHS

Memory corruption

Memory fragmentation

 Even using AS_PATH prepends, it is not normal to see more than 20 ASes in a typical AS_PATH in the Internet today

The Internet is around 5 ASes deep on average

Largest AS_PATH is usually 16-20 ASNs

Limiting AS Path Length

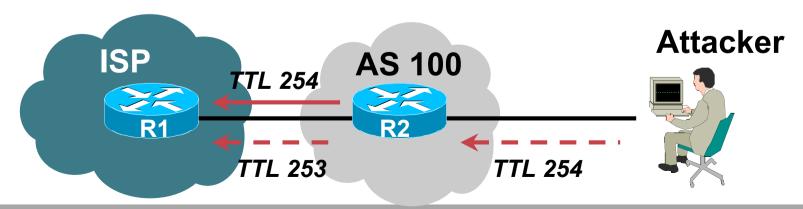
 Some announcements have ridiculous lengths of AS-paths:

This example is an error in one IPv6 implementation

 If your implementation supports it, consider limiting the maximum AS-path length you will accept

BGP TTL "hack"

- Implement RFC3682 on BGP peerings
 - **Neighbour sets TTL to 255**
 - Local router expects TTL of incoming BGP packets to be 254
 - No one apart from directly attached devices can send BGP packets which arrive with TTL of 254, so any possible attack by a remote miscreant is dropped due to TTL mismatch



BGP TTL "hack"

TTL Hack:

Both neighbours must agree to use the feature TTL check is much easier to perform than MD5 (Called BTSH – BGP TTL Security Hack)

Provides "security" for BGP sessions

In addition to packet filters of course

MD5 should still be used for messages which slip through the TTL hack

See www.nanog.org/mtg-0302/hack.html for more details

Passwords on BGP sessions

- Yes, I am mentioning passwords again
- Put password on the BGP session
 - It's a secret shared between you and your peer
 - If arriving packets don't have the correct MD5 hash, they are ignored
 - Helps defeat miscreants who wish to attack BGP sessions
- Powerful preventative tool, especially when combined with filters and the TTL "hack"

Using Communities

Use communities to:

Scale iBGP management

Ease iBGP management

Come up with a strategy for different classes of customers

Which prefixes stay inside network

Which prefixes are announced by eBGP

...etc...

Using Communities: Strategy

BGP customers

Offer max 3 types of feeds (easier than custom configuration per peer)

Use communities

Static customers

Use communities

Differentiate between different types of prefixes
 Makes eBGP filtering easy

Using Communities: BGP Customer Aggregation Guidelines

Define at least three groups of peers:

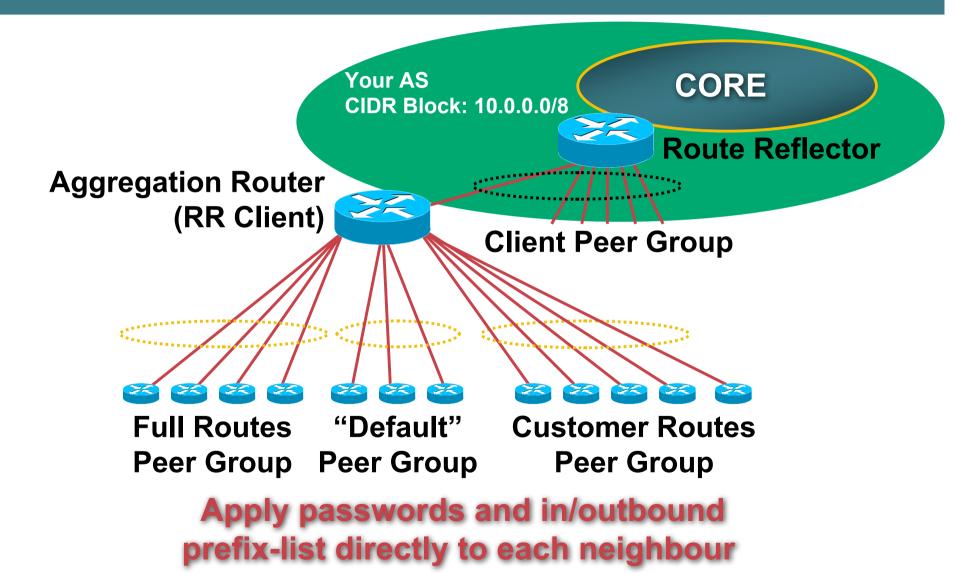
```
cust-default—send default route only cust-cust—send customer routes only cust-full —send full Internet routes
```

Identify routes via communities e.g.

```
100:4100=customers; 100:4500=peers
```

- Apply passwords per neighbour
- Apply inbound & outbound prefix filters per neighbour

BGP Customer Aggregation



Using Communities: Static Customer Aggregation Guidelines

Identify routes via communities, e.g.

100:4000 = my address blocks

100:4100 = "specials" from my blocks

100:4200 = customers from my blocks

100:4300 = customers outside my blocks

Helps with aggregation, iBGP, filtering

 Set correct community as networks are installed in BGP on aggregation routers

Using Communities: Sample core configuration

eBGP peers and upstreams

Send communities 100:4000, 100:4100 and 100:4300, receive everything

iBGP full routes

Send everything (only to network core)

iBGP partial routes

Send communities 100:4000, 100:4100, 100:4200, 100:4300 and 100:4500 (to edge routers, peering routers, IXP routers)

Summary

- Use configuration templates
- Standardise the configuration
- Be aware of standard "tricks" to avoid compromise of the BGP session
- Anything to make your life easier, network less prone to errors, network more likely to scale
- It's all about scaling if your network won't scale, then it won't be successful



BGP Techniques for Internet Service Providers

Philip Smith <pfs@cisco.com>
APRICOT 2007
Bali, Indonesia
27th February 2007