

Delivering High Availability Routed Networks

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Effects of Network Outage

■ Immediate Impact

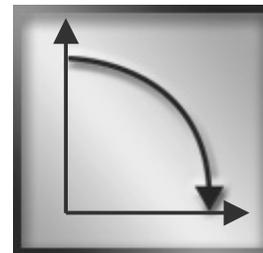
- Loss of Revenue
- Repair Costs
- SLA penalties
- Dissatisfied customers
- Project delays
- Management distraction

■ Long Term Impact

- Damage to corporate brand
- Customer churn, market share
- Competition
- Lawsuits
- Lack of internal confidence



Financial



Market Share

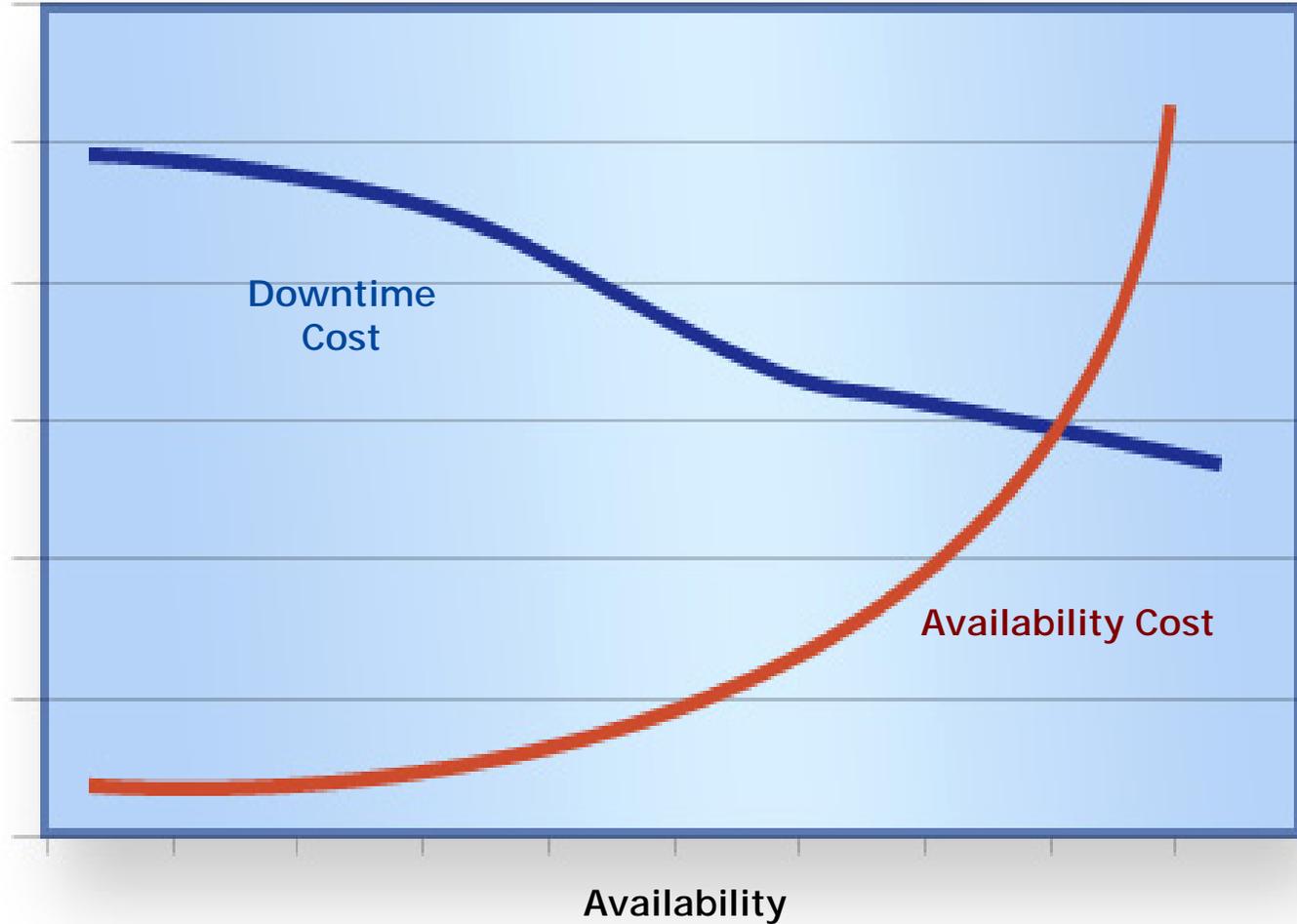


Brand Damage

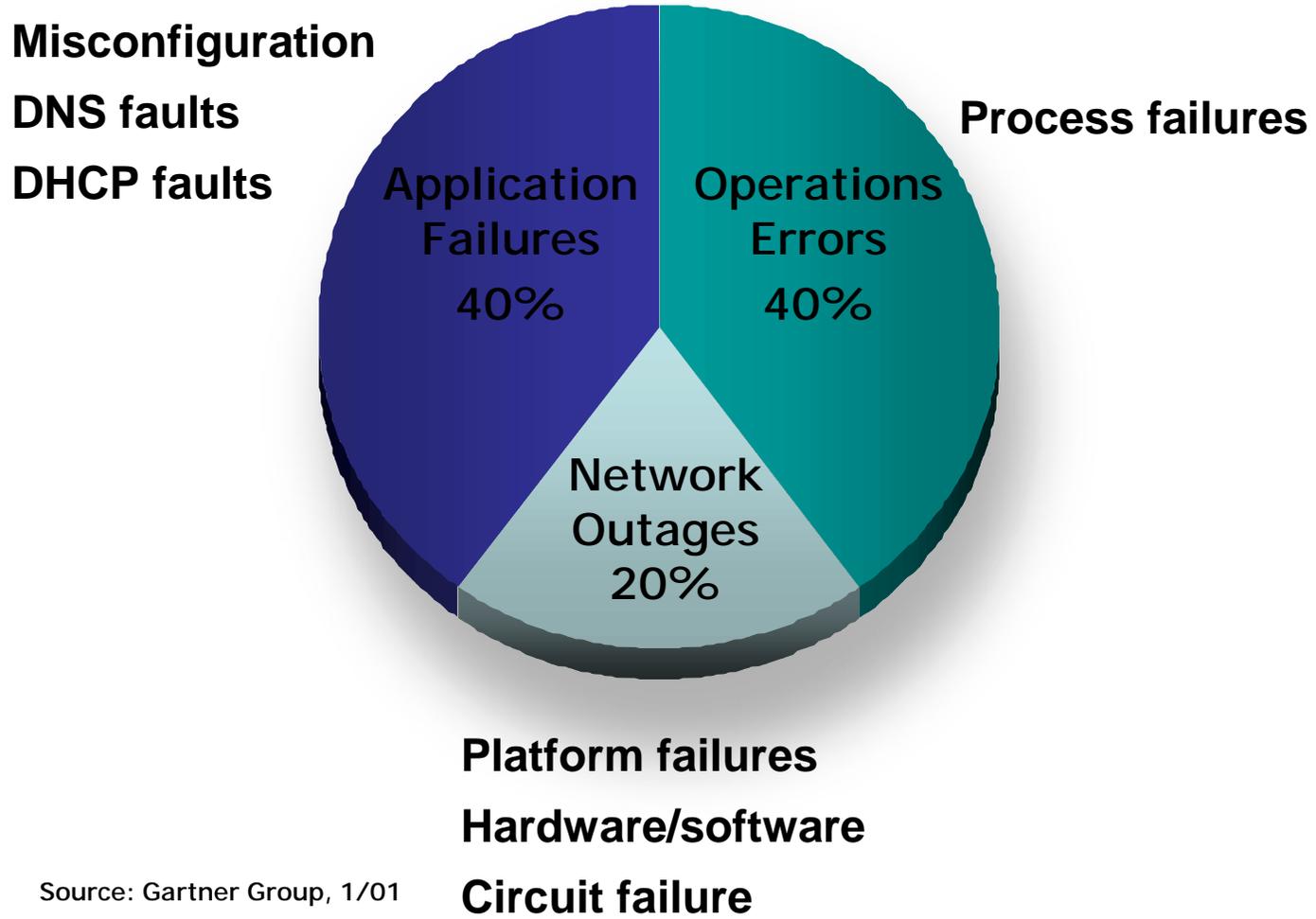
Business Case for High Availability



Cost



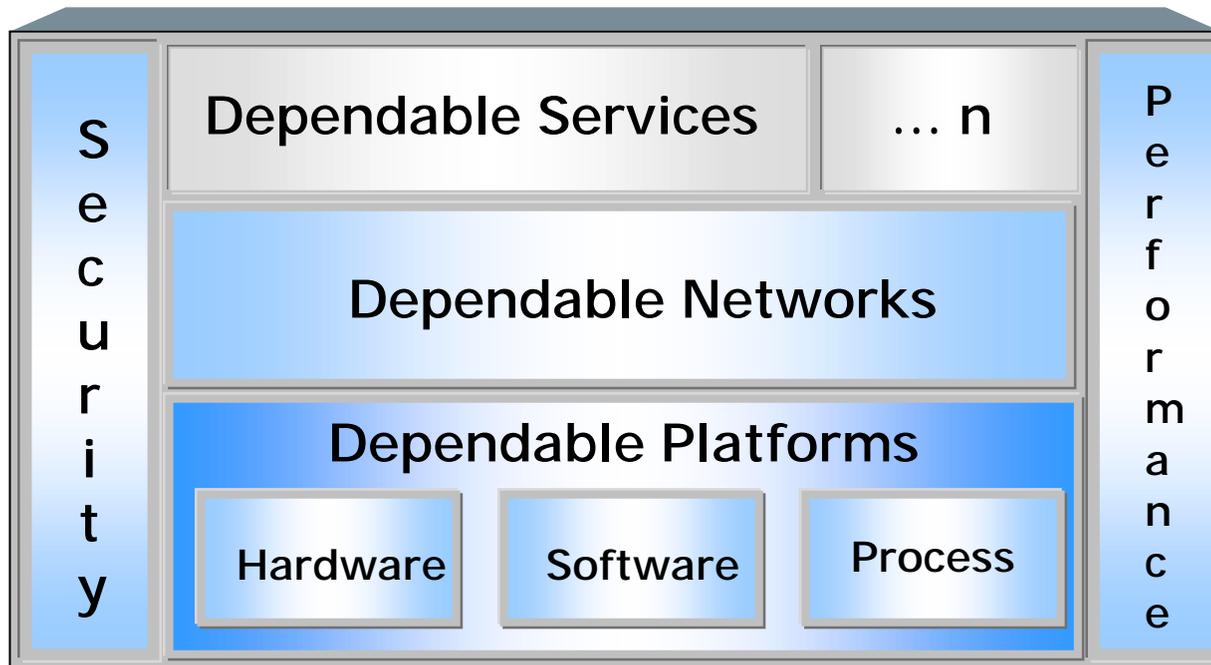
Threats to Dependability



Source: Gartner Group, 1/01

HA Solution Architecture

IP Carrier-Class Availability Is a Culture, Not a Single Feature, Protocol or Product



Reliable Hardware

- **Hard Fault Tolerance**

- Environmental sensors
- Component redundancy
- Redundant boot devices

- **Soft Fault Tolerance**

- Extensive internal diagnostics
- CRC-protected internal data paths
- ECC SDRAM

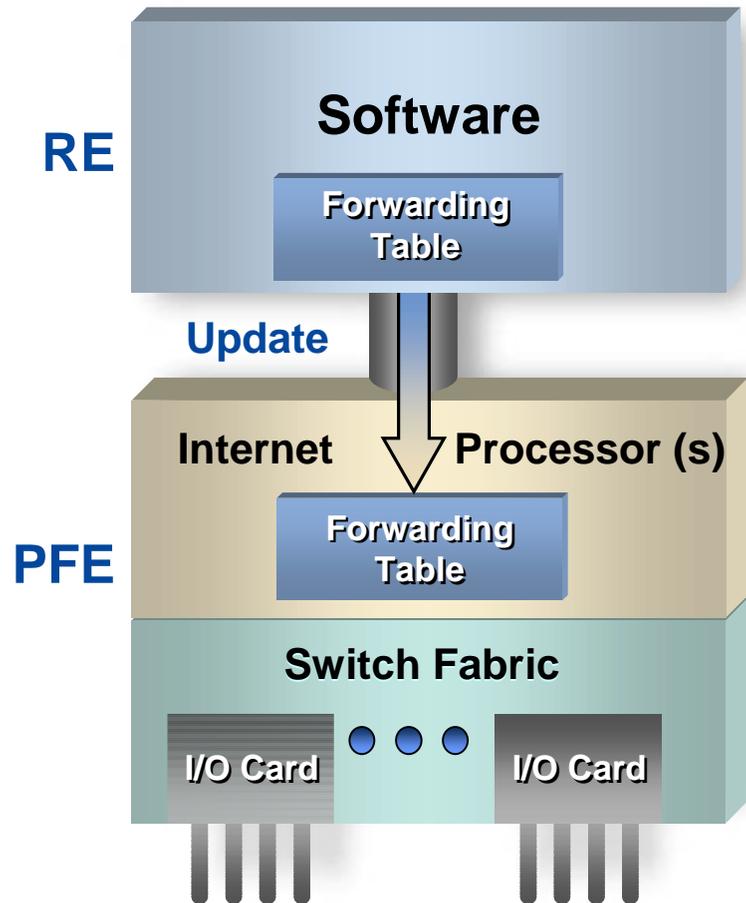
- **MTTR Reduction**

- Hot swappable components
- Field replaceable components

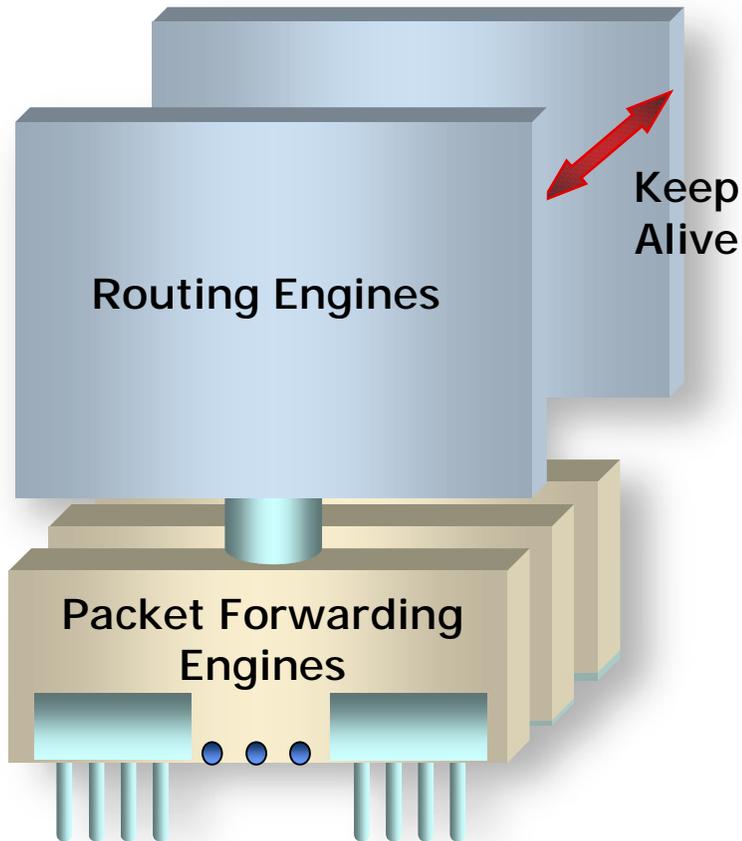


A Logical Platform View

- Hardware modularity is fundamental
- Clean separation of routing and packet forwarding functions
- Different vendors have different names, but for example:
 - **Routing Engine (RE)**
 - Routing protocol and management functions
 - **Packet Forwarding Engine (PFE)**
 - Packet forwarding and processing
- Multiples of each module allow redundancy and failover

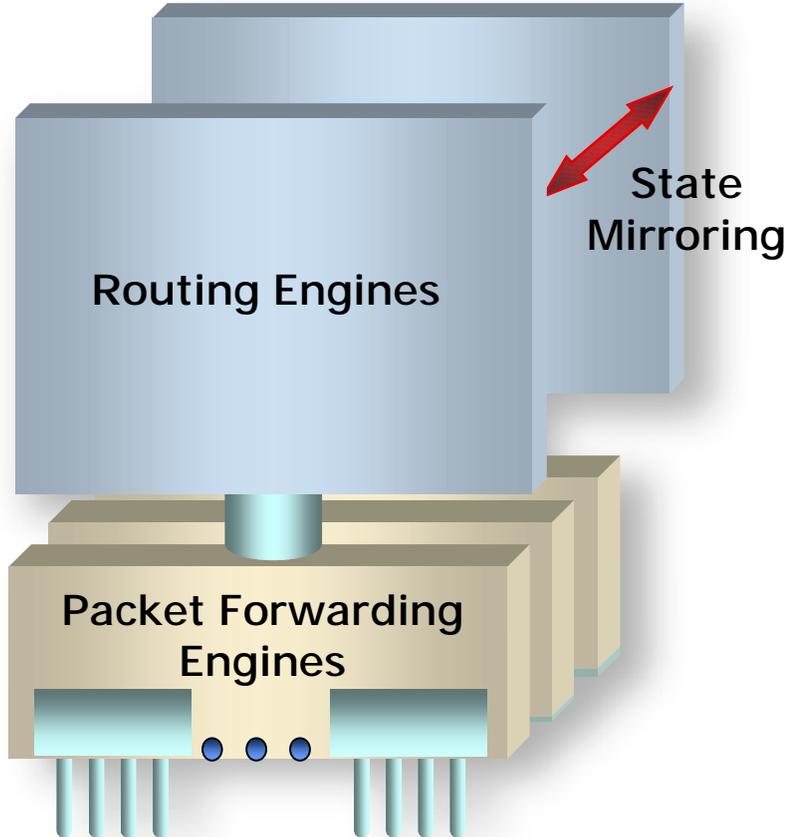


Simple RE Failover



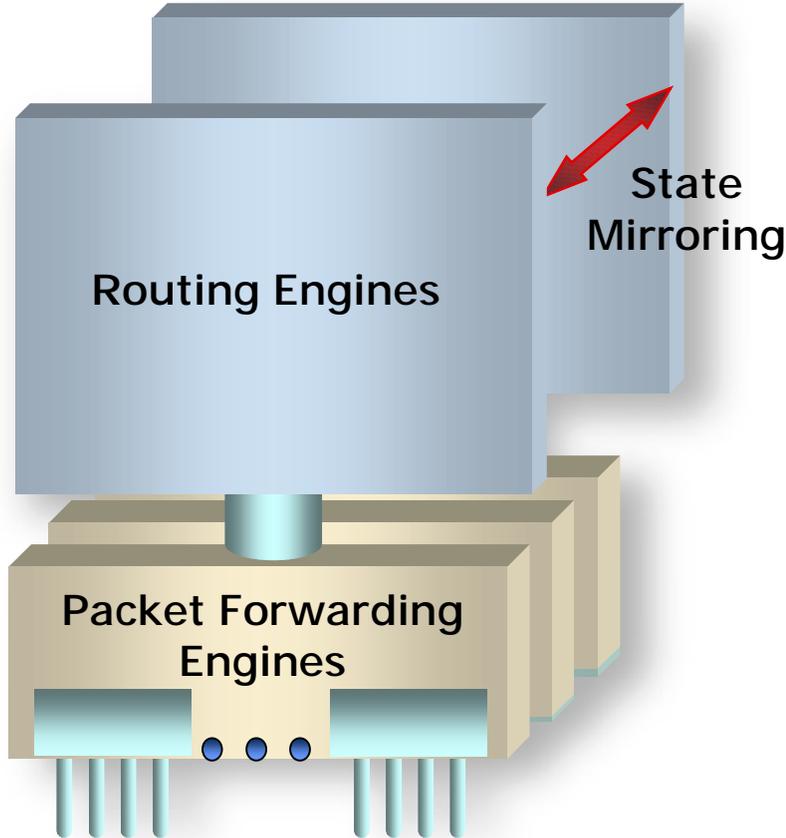
- Protects against Single Node Hardware Failure
- Redundant Routing Engines run keepalive process
- Automatic failover to secondary
- Configuration synchronized between RE's
- Configurable timer
- Routing Process restarts
- Requires PFE reset

Stateful Protocol Mirroring



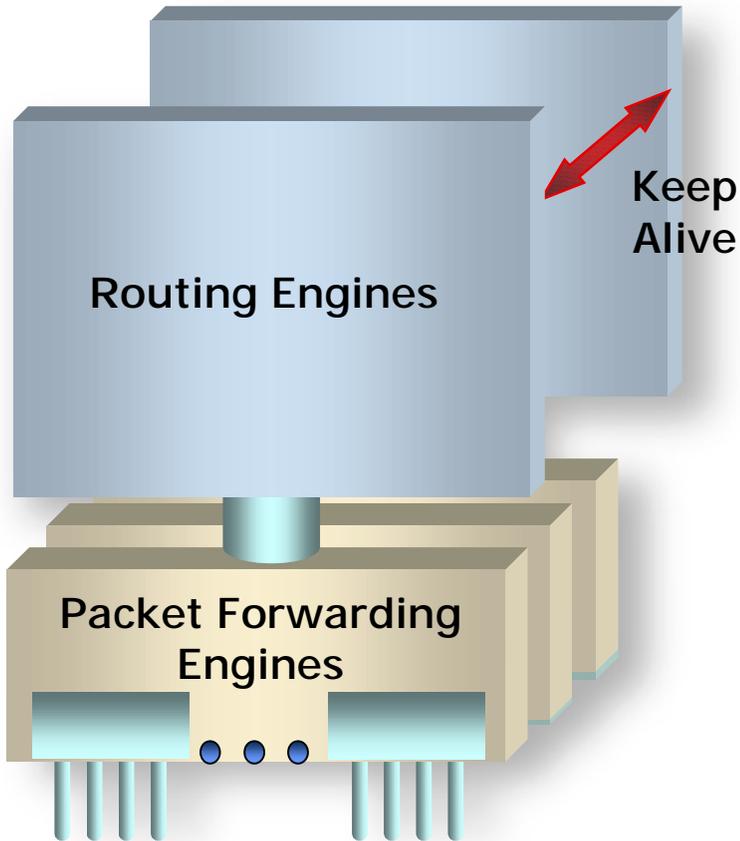
- Protects against Single Node Hardware Failure
- Redundant Routing Engines Mirror each others state
- BGP & TCP
- Theoretically ISIS & OSPF
- Automatic failover to secondary
- Advocated by some vendors, claiming Carrier-Class IP

Stateful Protocol Mirroring



- **Great Idea!**
- **Difficult to do without replicating errors as well as “good” state**
- **Potential for “bug mirroring”**
- **Much more challenging in a rich service environment than an IP-only core**

Graceful RE Switchover



- Protects against **Single Node Hardware Failure**
- **Primary (REP) and Secondary (RES) utilize keepalive process**
 - Automatic failover to RES
 - Synchronized Configuration
- **REP and RES share:**
 - Forwarding info + PFE config
- **REP failure does not reset PFE**
 - No forwarding interruption
 - Only Management sessions lost
 - Alarms, SNMP traps on failover

Reliable Software

■ Hard Fault Tolerance

- Redundant REs
- Different software versions

■ Soft Fault Tolerance

- Separate control and forwarding
- Modular processes can be restarted independently
- Processes protected in own memory space
- Individual process watchdogs

■ MTTR Reduction

- Incremental software upgrades
- Modularity to speed up testing

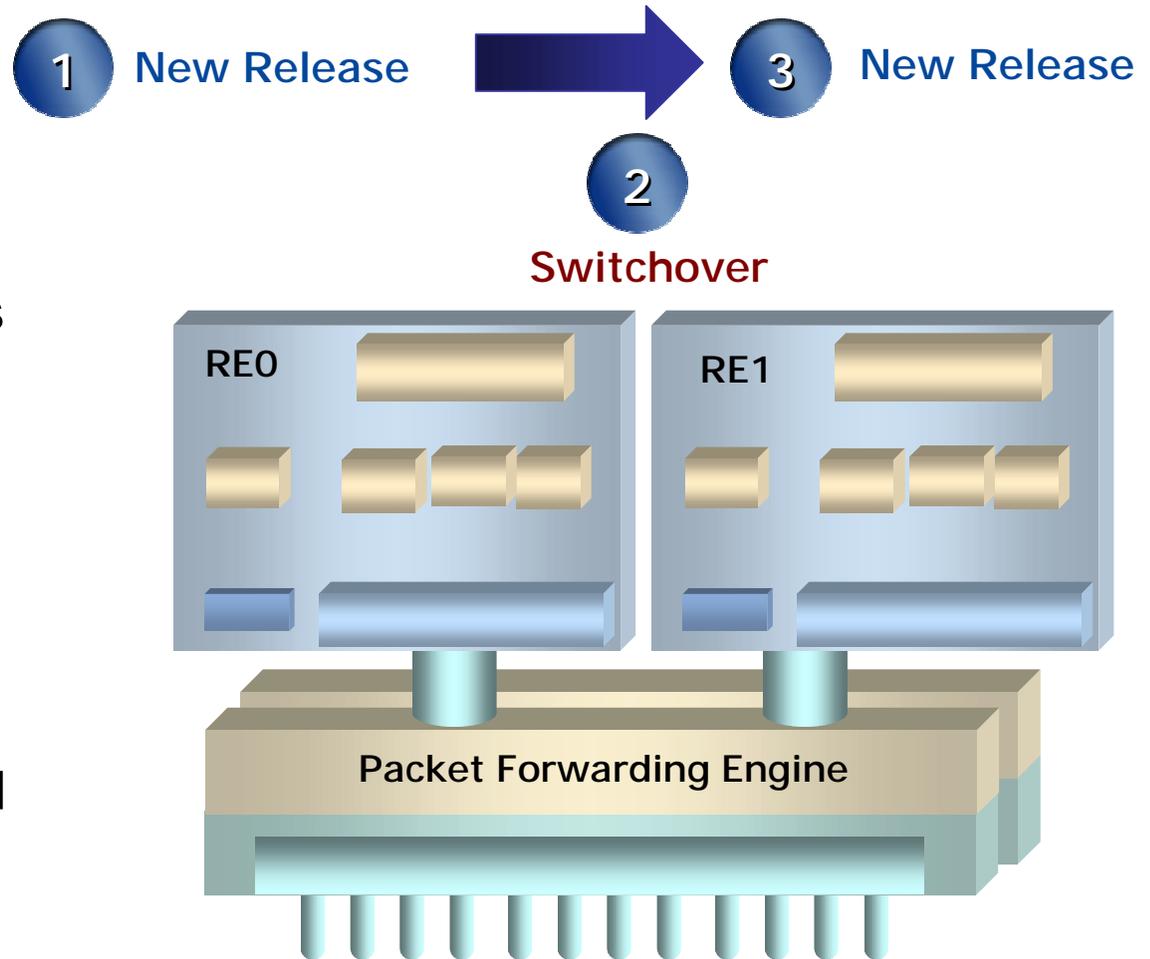


Software Reliability Principles

- **Loose coupling of modular components**
 - A single failing component will not crash the box
 - Localizes complexity
 - Creates conceptual boundaries to contain problems
 - Clean interfaces between system components (well-defined, efficient APIs)
- **Memory protection**
 - Processes cannot scribble on each others' memory
- **Adding complexity will not improve reliability**
 - If base software is not expandable, maintainable, reliable, then adding additional layers won't help
 - "Make it as simple as possible, but no simpler."
--Albert Einstein

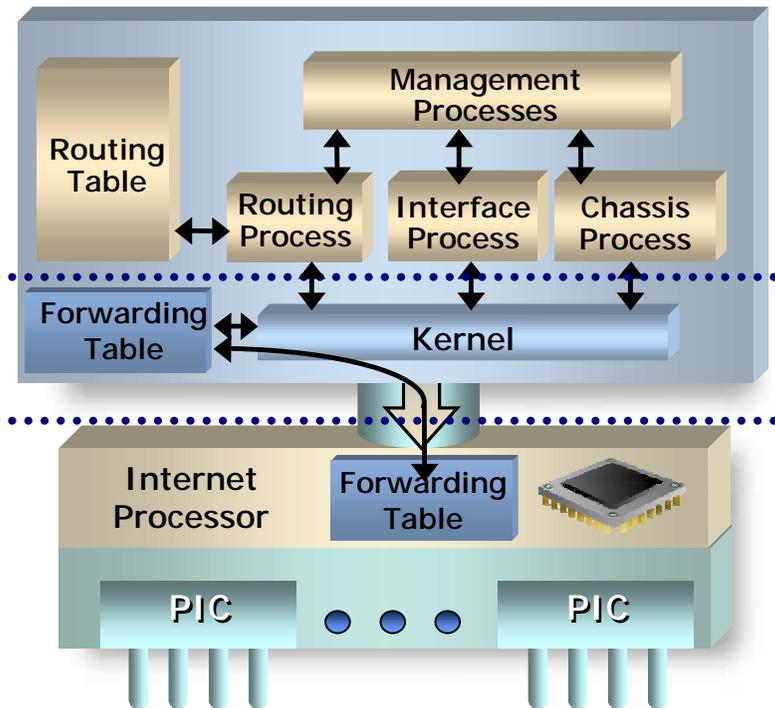
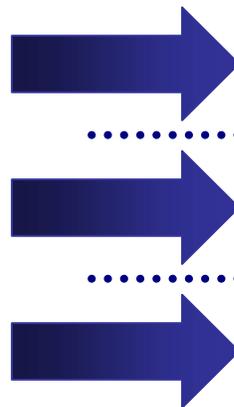
In-Service Software Upgrades

- **Leverages**
 - Graceful RE Switchover
 - Graceful Restart Protocol Extensions
- **Preserves forwarding**
 - In any RE failure
- **Delivers**
 - In-service software upgrades
- **Might also be enabled by stateful mirroring**

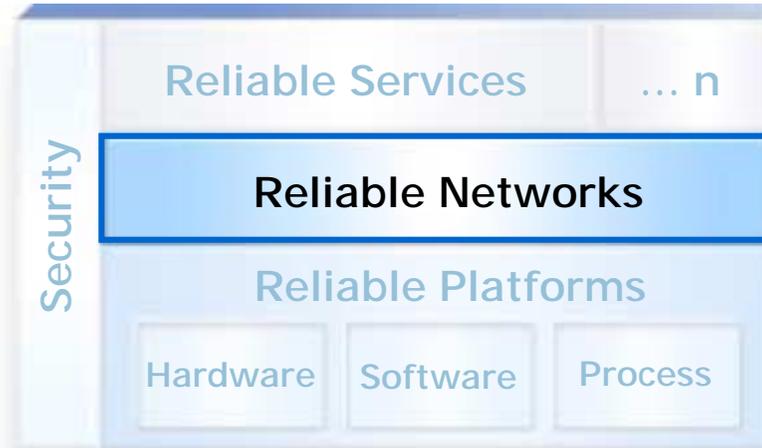


In-Service Software Upgrades

- When Software is modular:
- (JUNOS, for example)
 - “jinstall” is a complete software distribution
- “jroute”
 - Routing protocols
- “jkernel”
 - Operating system
- “jpfe”
 - PFE software



Reliable Networks

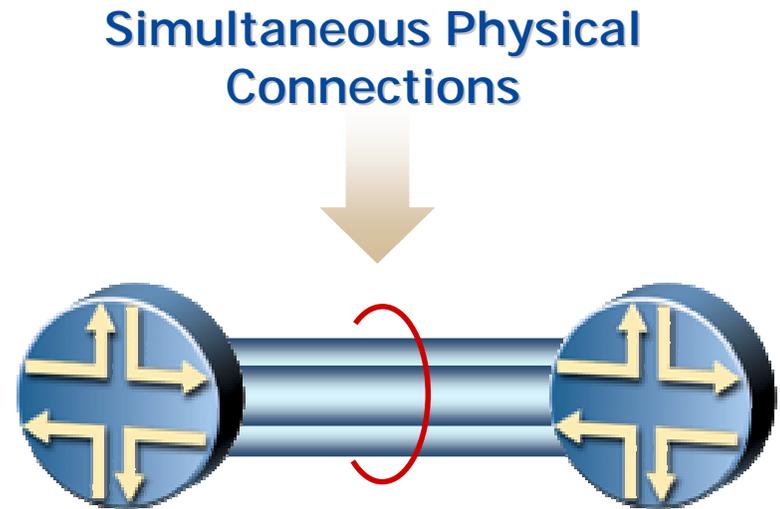


Protection and Recovery from failures

- ◆ MPLS
 - ◆ Fast reroute
 - ◆ Secondary LSPs
- ◆ VRRP
- ◆ Convergence improvements
- ◆ Graceful Restart
- ◆ Link Redundancy
- ◆ Multi-Homing
- ◆ SONET APS/SDH MSP

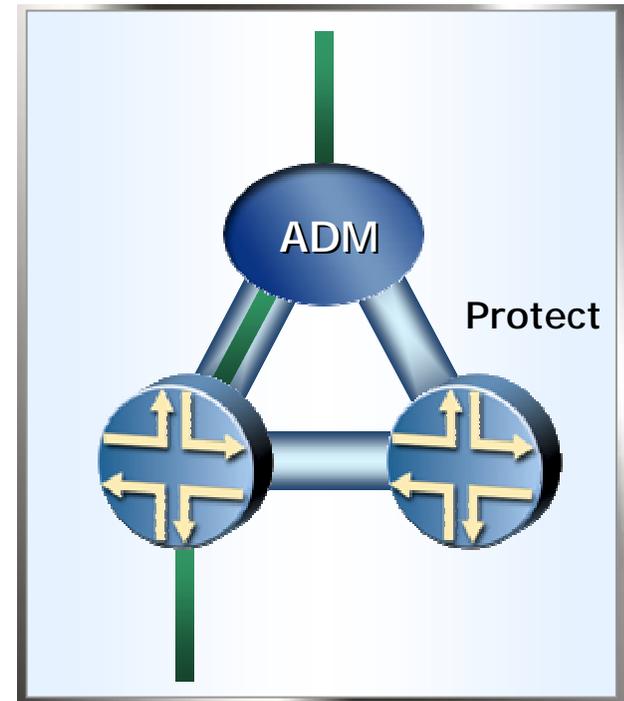
Link Redundancy

- **Reliable Links**
 - Link failure does not affect forwarding
 - Load redistributed among other members
- **Parallel Link Technologies**
 - MLPPP – T1/E1 Link aggregation
 - Multi-Link Frame Relay
 - 802.3ad – Ethernet aggregation
 - SONET/SDH aggregation



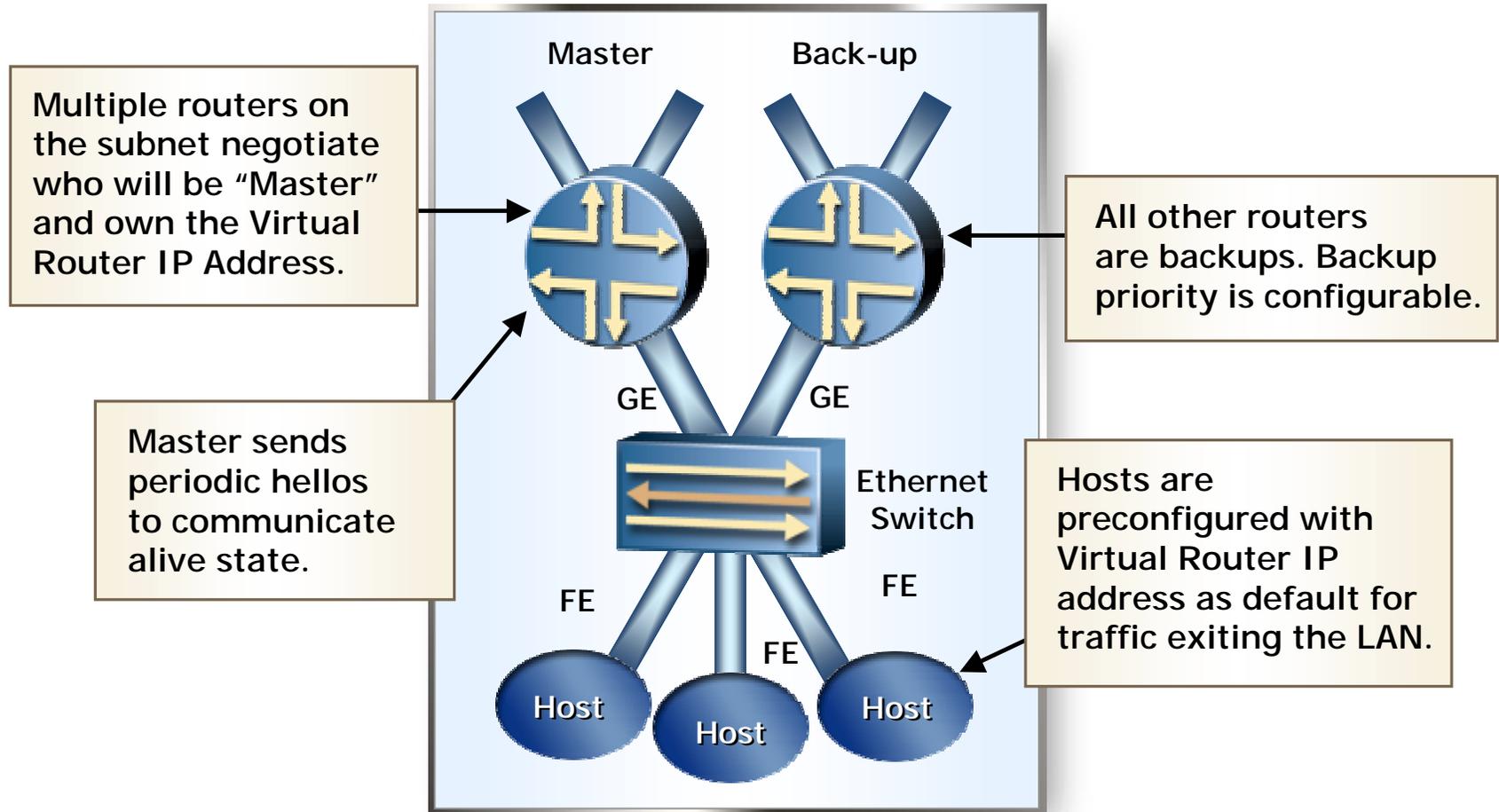
SONET/SDH Protection Switching

- **SONET APS & SDH MSP**
 - Redundant routers share uplink
- **Rapid circuit failure recovery**
 - Used on router-to-ADM links
 - Layer 3 protocol convergence longer
- **Interoperable with standard ADM**
- **Working & protect circuits**
 - May reside on different routers
 - May reside on same router

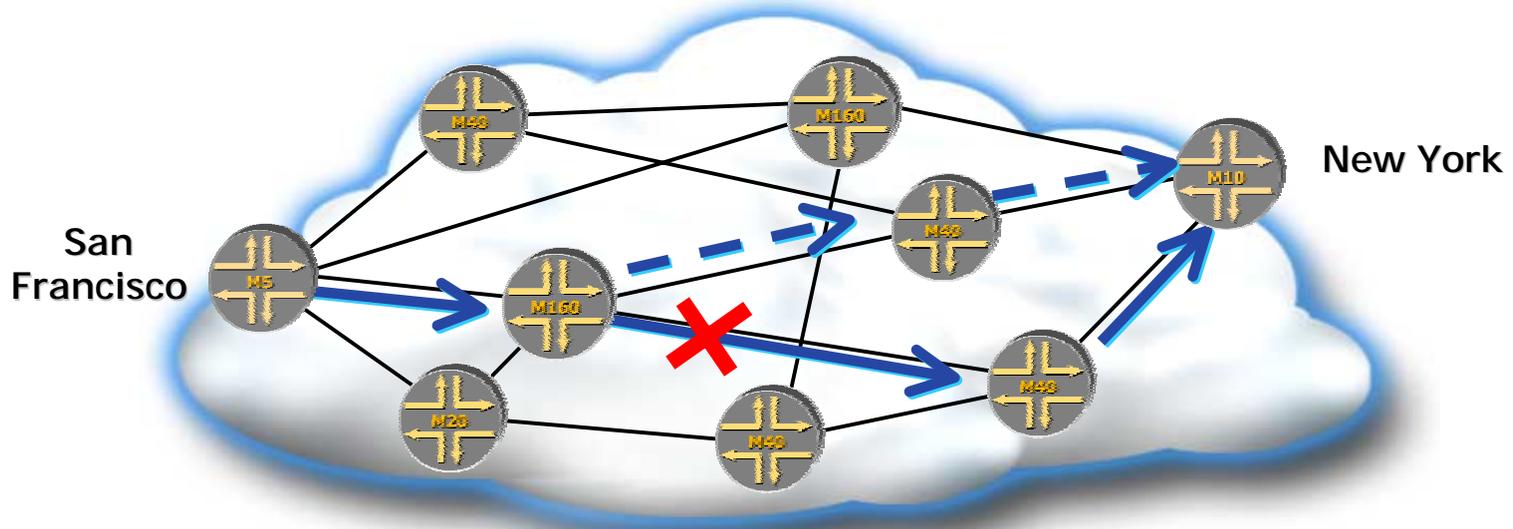


Virtual Router Redundancy Protocol

- Redundant default gateways–VRRP (RFC 2338)



IP Dynamic Routing



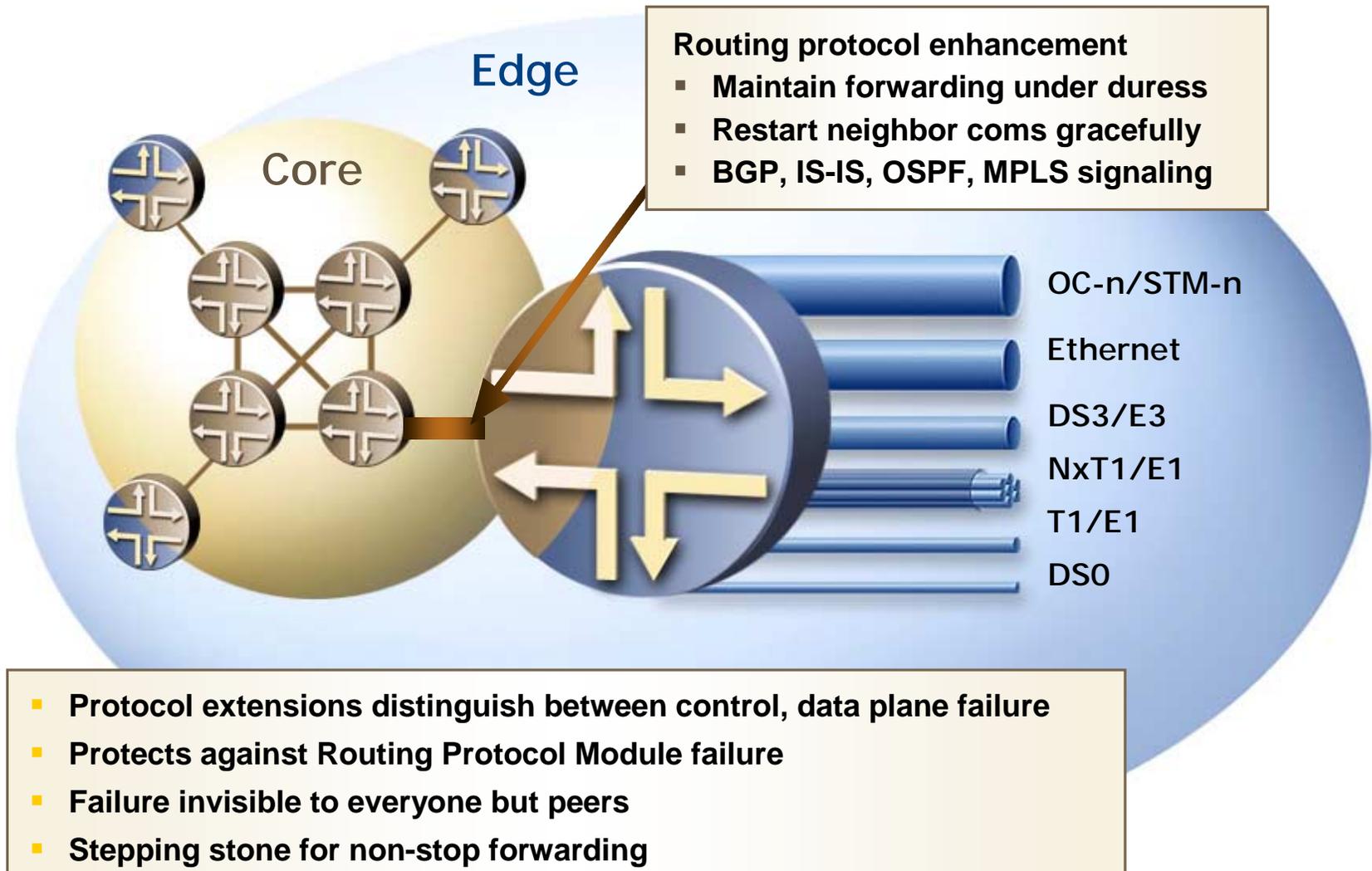
- **OSPF or IS-IS computes path**
- **If link or node fails, New path is computed**
- **Response times: Typically a few seconds**
- **Completion time: Typically a few minutes, but very dependant on topology**

Faster Router Convergence

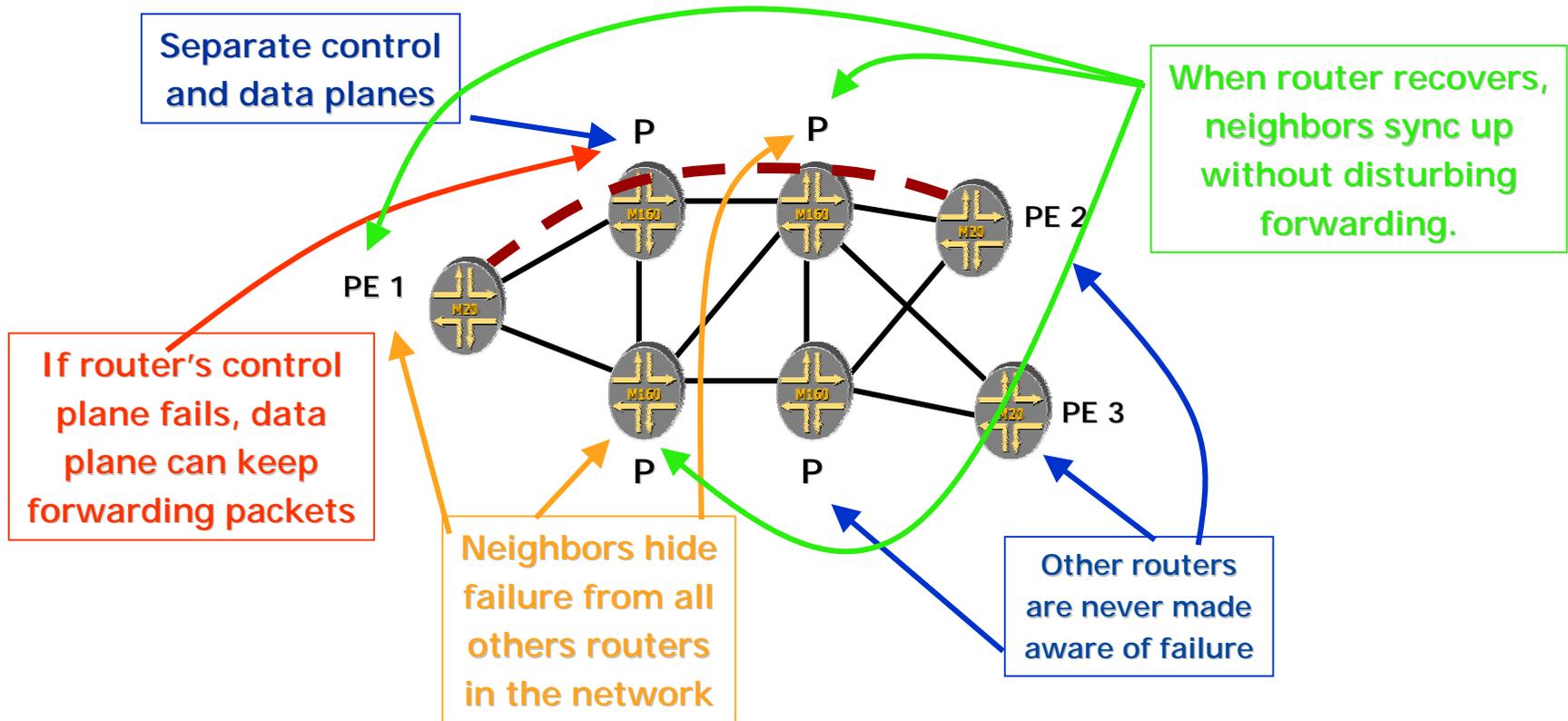
- **Faster convergence improves Network Reliability**

Features	Benefits
High Priority Flooding for Interested LSPs (ISIS / OSPF)	<ul style="list-style-type: none">■ Timer reduced from 100 to 20msec■ Faster propagation of major changes
Quick SPF Scheduling (ISIS / OSPF)	<ul style="list-style-type: none">■ Reduces time from 7 sec to 50 msec■ Speeds calculation of optimum path
Sub-second Hellos (ISIS)	<ul style="list-style-type: none">■ Lowest Hello Time possible for IS-IS, 333msec■ Faster Link Failure Detection
RIB and FIB Enhancements (BGP)	<ul style="list-style-type: none">■ Indirect Next Hop implies faster convergence

Routing Protocol Graceful Restart



Graceful Restart - How ?



Graceful Restart Protocol Details

Purpose - Continue forwarding (PFE) during a restart of routing (RE)

	Changes	IETF
BGP	<p>Protocol extensions</p> <p>Per-peer configuration</p> <p>Various timers with configurable defaults</p>	<p><i>Graceful Restart Mechanism for BGP</i></p> <p>draft-ietf-idr-restart-08.txt</p>
OSPF	<p>Protocol extensions</p> <p>New opaque-LSA type 9, "Grace-LSA"</p>	<p><i>Hitless OSPF Restart</i></p> <p>rfc3623</p>
IS-IS	<p>Protocol extensions</p> <p>3 new timers</p> <p>New "re-start" option (TLV) in IIH PDU</p>	<p><i>Restart Signaling for ISIS</i></p> <p>draft-shand-isis-restart-04.txt</p>
MPLS	<p>Protocol Extensions</p> <p>Uses signaling as described in "Graceful Restart Mechanism for BGP"</p>	<p>Graceful Restart Mechanism for BGP with MPLS</p> <p>draft-ietf-mpls-bgp-mpls-restart-03.txt</p>
RSVP	<p>Protocol Extensions</p> <p>Extend rfc 3473</p> <p>Recovery ERO</p>	<p>Graceful Restart Extensions</p> <p>draft-rahman-rsvp-restart-extensions-00.txt</p>

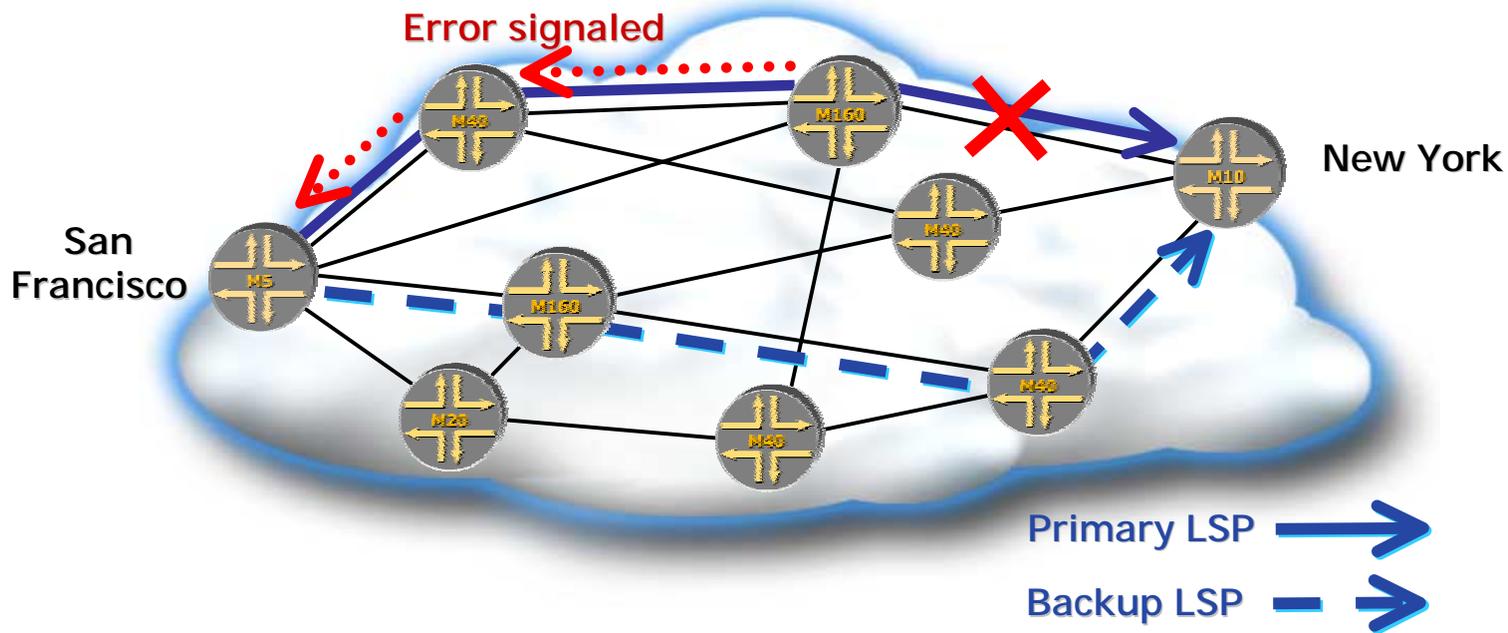
MPLS-based mechanisms

- Path protection (aka Secondary LSP)
- Local 1:1 (aka LSP/Detour Protection Fast Reroute)
 - Protects against both link failure and node failures
- Local 1:N (aka Facility-based Fast Reroute)
 - Link Protection Fast Reroute (Protects only against link failure)
 - Node Protection (Protects against both link failure and node forwarding plane failure)

Secondary LSPs

- An LSP may have multiple paths
- Primary path is the preferred path to set up and use
- Secondary paths are alternatives, to be used when the primary fails
 - Usually node/link disjointed from primary
 - The level of overlap between the primary and the secondary could be controlled
- Secondary path may result in wasting resources
 - Resources reserved for secondary are reserved all the time, yet used only when the primary fails

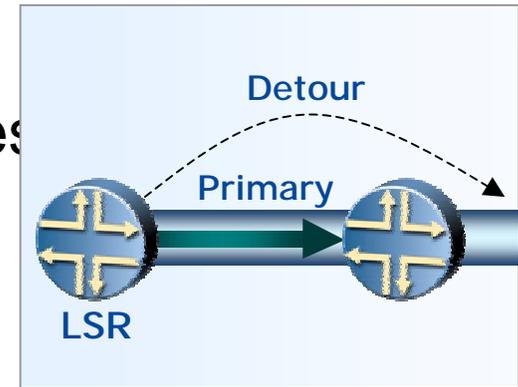
Secondary LSPs



- **Primary & secondary LSPs established a priori**
- **If primary fails**
 - Signal to ingress router to use secondary LSP
- **Faster response than routing protocol, requires wide area signaling**

MPLS Fast Reroute

- **Increasing demand for “APS/MSP-like” redundancy**
 - MPLS resilience to link/node failures
 - Control-plane protection required
 - Frequent code upgrades = instability
 - Cost of APS/MSP protection
- **Solution: MPLS Fast-reroute**
 - RSVP Extensions define Fast Reroute



Fast Reroute

- Head-end of LSP enables fast reroute
- When signaled, each intermediate node calculates its own path to the tail-end
 - Uses CSPF and reservation
 - Doesn't duplicate reservations on a single link (but does duplicate on the network as whole)
- If any node sees the interface over which the primary LSP is routed go down, that node can instantly switch to backup
- Head-end discovers later and can reroute in a way that is more globally optimal

Complexity Comparison

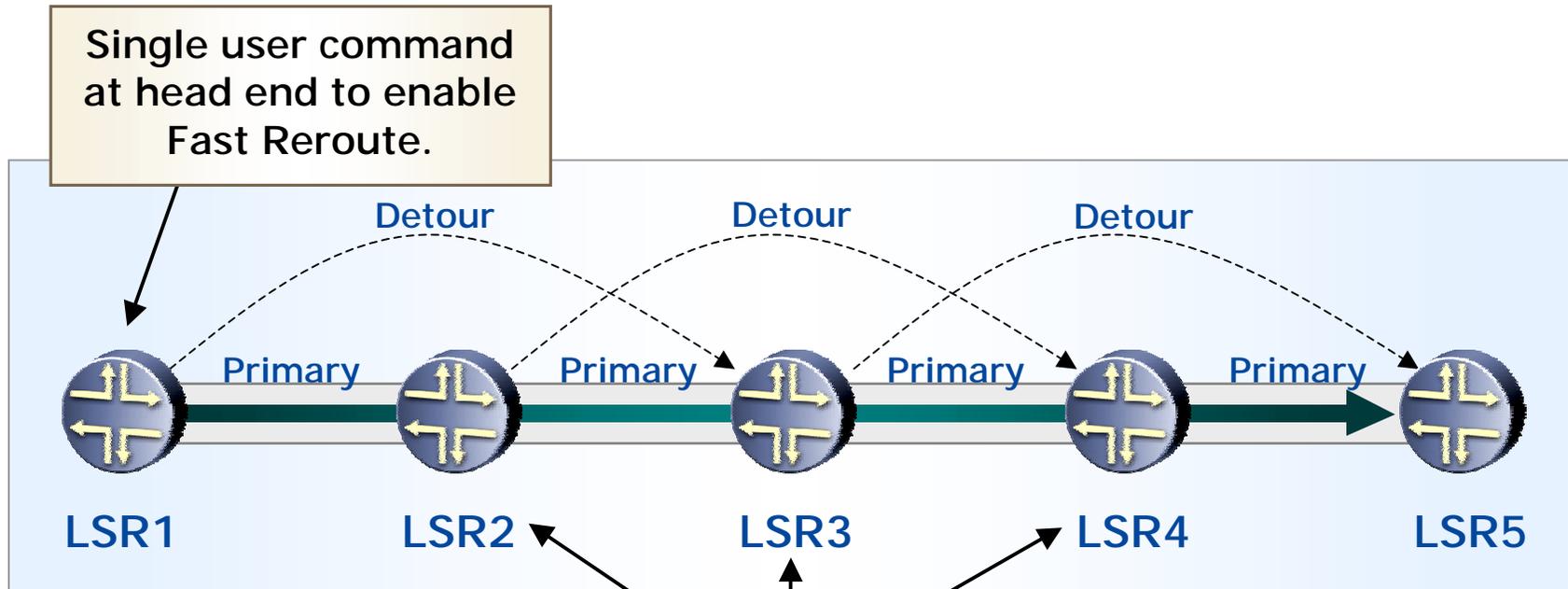
■ Secondary LSPs

- Signaled by ingress LSR only, protects path
- + additional constraints can be applied
- + tries to stay away from primary path nodes and links
- - additional management and planning
- - switch is done at the ingress router only
- + more scalable

■ FRR

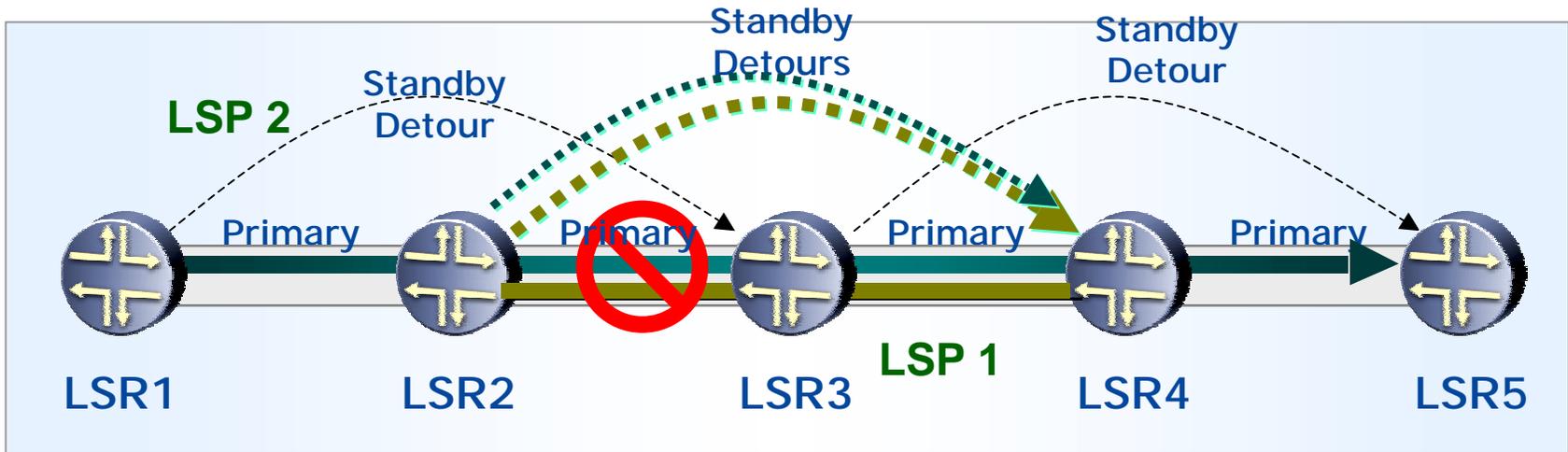
- Each LSR along the path protects configured links
- - limited path constraints
- + no additional path definitions configuration

Local 1:1 Protection Operation



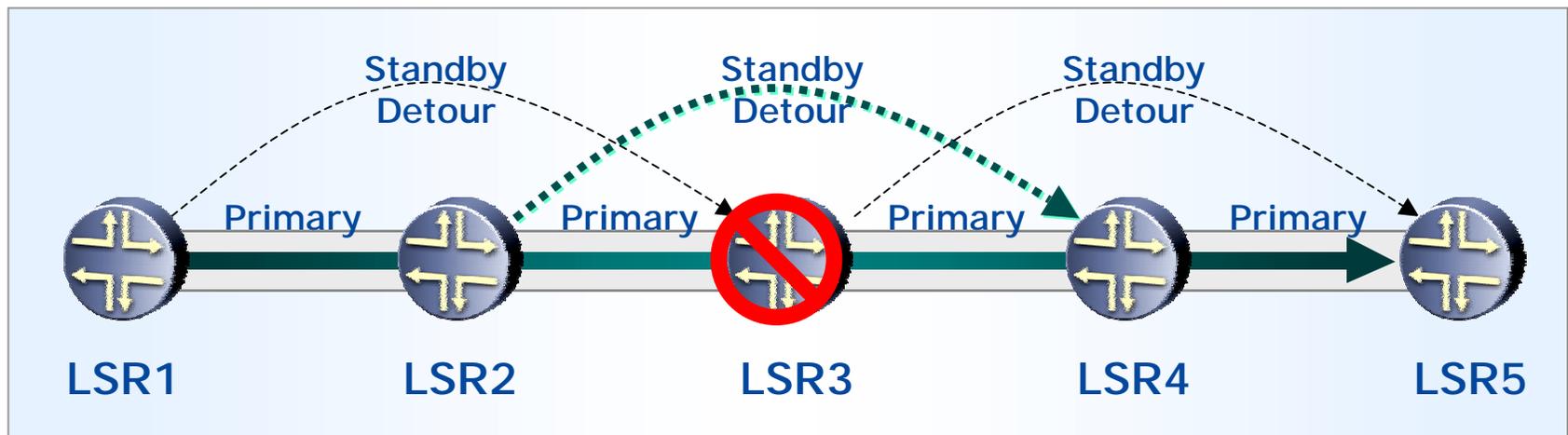
- Fast reroute is signaled to each LSR in the path
- Each LSR computes and sets up a detour path that avoids the next link and next LSR
- Each LSR along the path uses the same route constraints used by head-end LSR

Local 1:1 Protection Operation: Link Failure



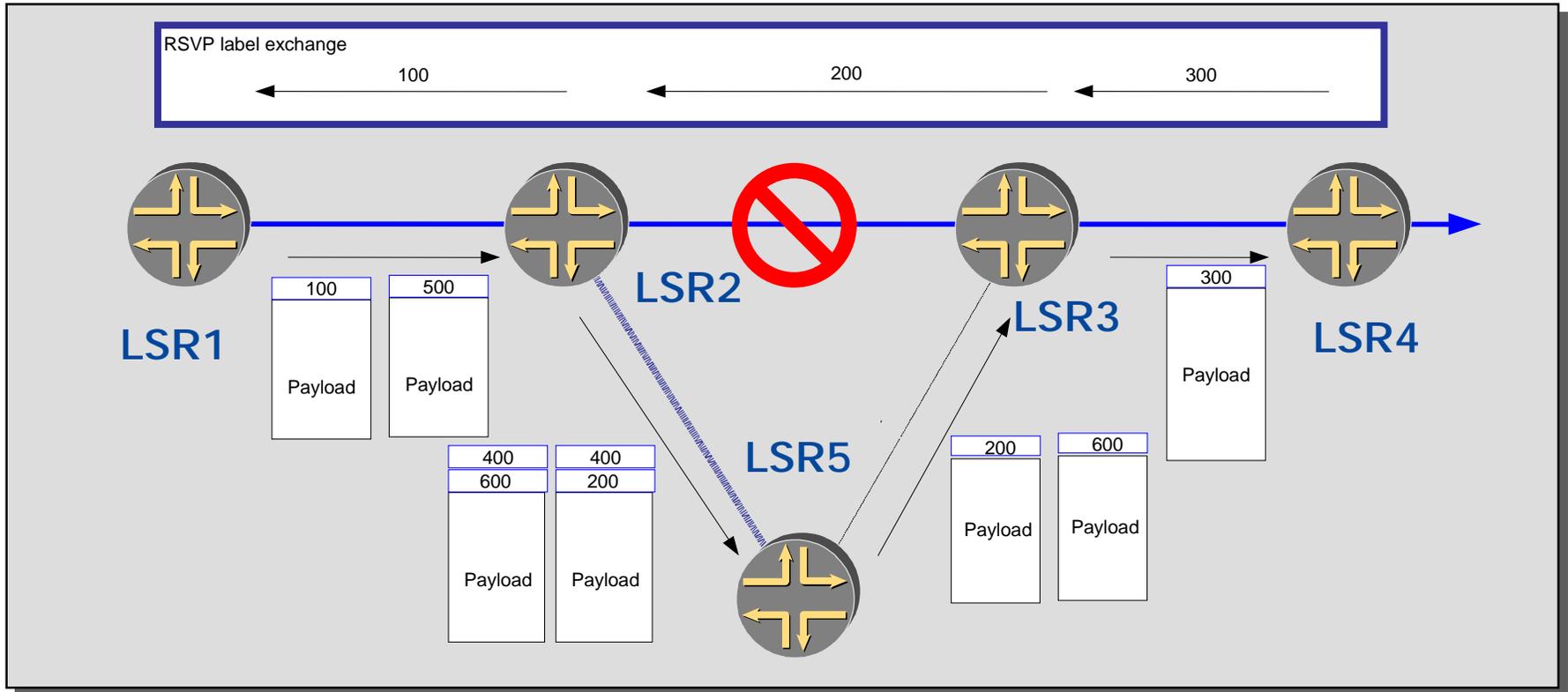
- **LSR2 detects that an interface in an LSP has gone down and reroutes via standby detour**
 - Recovery time is limited by the time to detect the failure
 - Comparable to SONET APS
- **Packet loss is minimized to the unlucky few that were transiting at the time of failure**

Local 1:1 Protection Operation: Node Failure



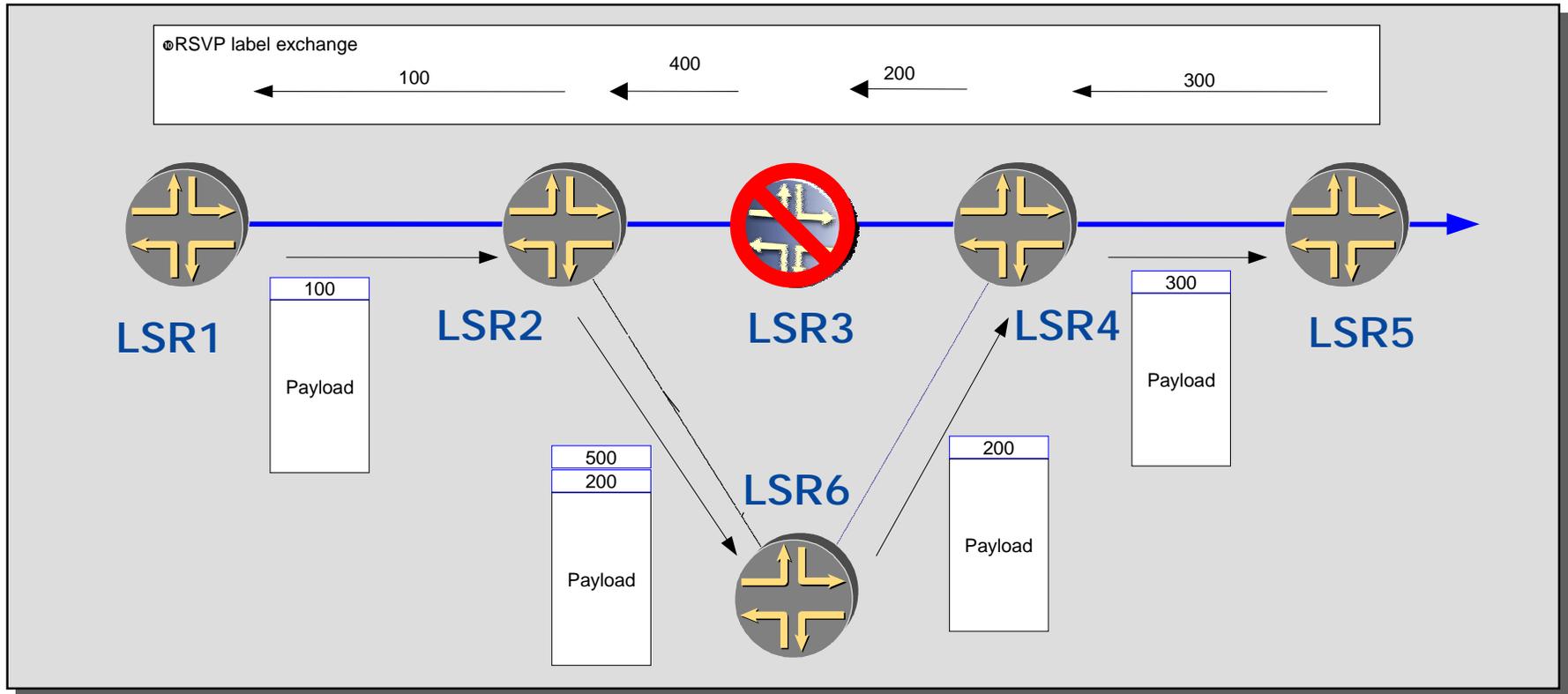
- **LSR2 detects that neighbor's (LSR3) forwarding plane has gone down and reroutes via standby detour**
 - Recovery time is limited by the time to detect the failure
- **Packet loss is minimized to the unlucky few that were transiting at the time of failure**

1:N Link Protection



- ◆ Each LSR detects that an interface has gone down and reroutes all the Protected LSPs traversing the interface via the Bypass LSP
 - ❖ Recovery time is limited by the time to detect the failure
- ◆ Packet loss is minimized to the unlucky few that were transiting at the time of failure

1:N Node Protection

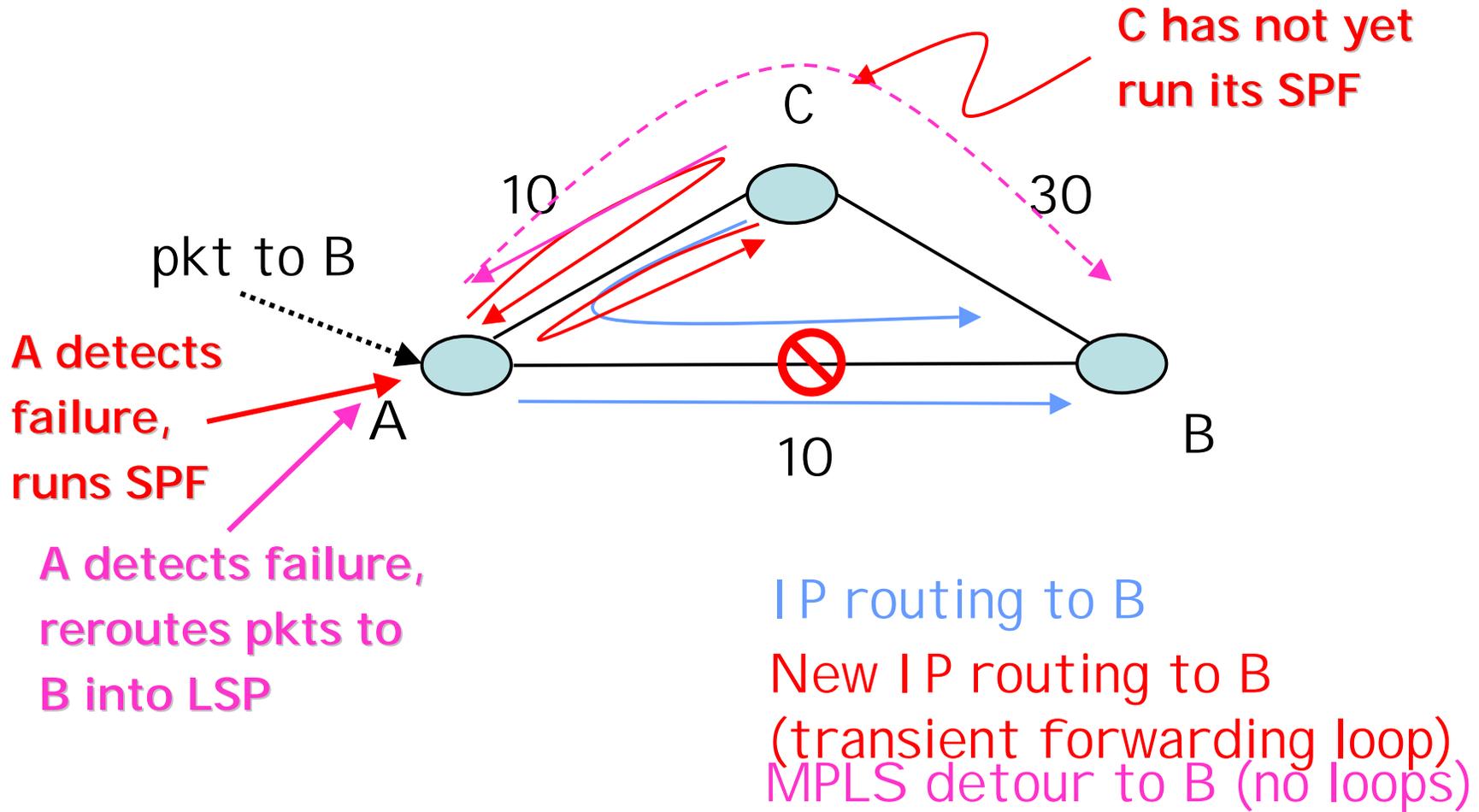


- ◆ Each LSR detects that an interface has gone down and reroutes all the Protected LSPs traversing the interface via the Bypass LSP
 - ❖ Recovery time is limited by the time to detect the failure
- ◆ Packet loss is minimized to the unlucky few that were transiting at the time of failure

Which one to use?

- 1:1 Detour Backup
 - The number of LSPs to be protected is small
 - Finer control (at the granularity of individual LSPs) with respect to LSP priority, bandwidth, link coloring for detour/bypass LSPs is important
 - Simpler configuration is desired
 - Suitable if LSP's have divergent paths
- 1:n Facility Backup
 - Ability to protect all the LSP's on a link with a single LSP with stacking

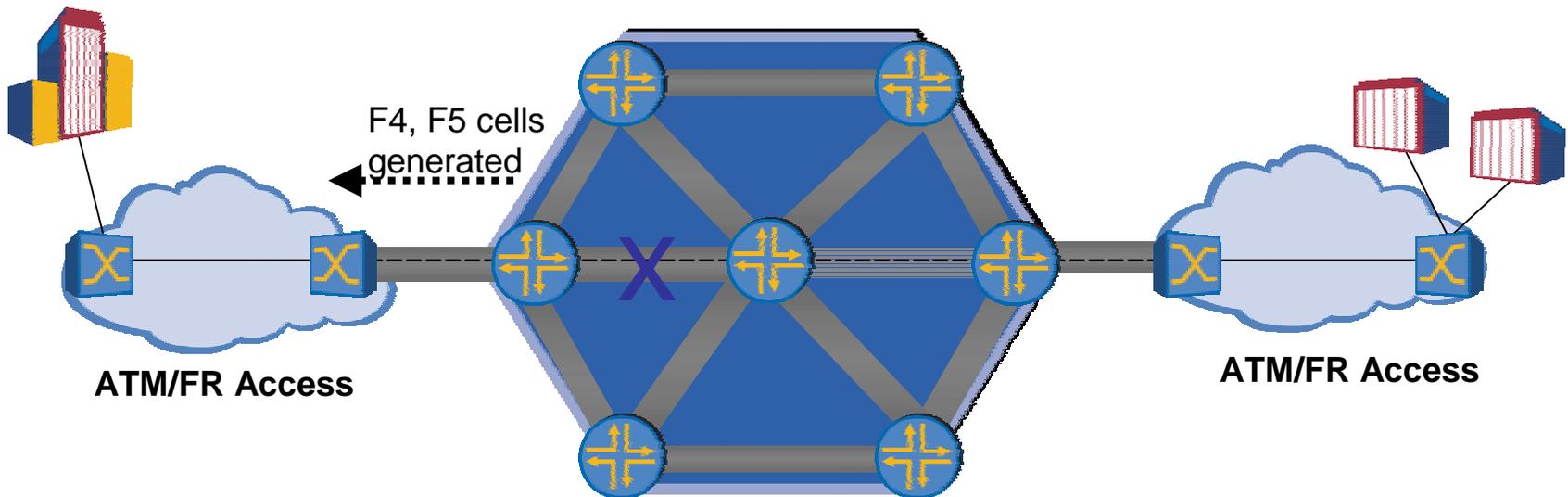
MPLS Fast Reroute vs IP



Extending to Legacy Networks

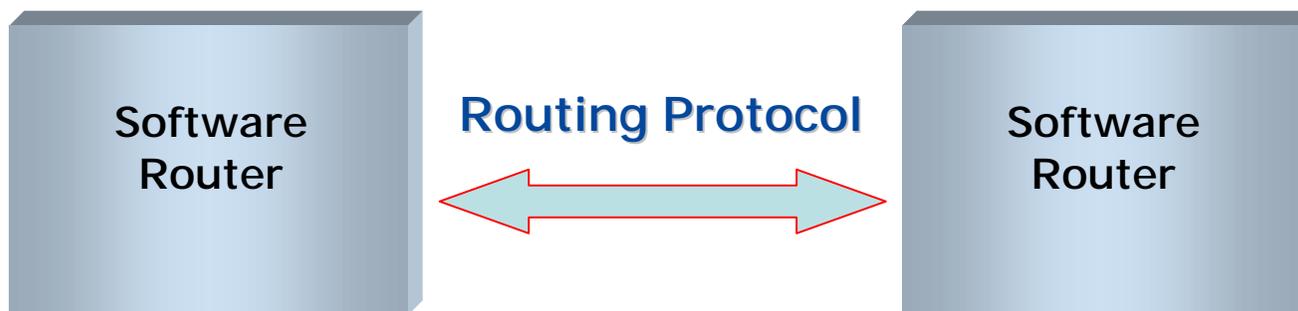
- **MPLS OAM features**

- Use BFD and FRR, along with other mechanisms
- Provides notification to external networks if LSP fails



BFD: Forwarding Liveliness (Bidirectional Forwarding Detection)

- **In IP, historically a function of the routing protocol**
 - Because formerly, routing = forwarding
 - Fault resolution in perhaps tens of seconds
 - This is too slow for anything but best-effort IP
 - Sometimes there is no routing protocol!



Goals of BFD

- **Faster convergence of routing protocols, particularly on shared media (Ethernet)**
- **Semantic separation of forwarding plane connectivity and control plane connectivity**
- **Detection of forwarding plane-to-forwarding plane connectivity (including links, interfaces, tunnels etc.)**
- **A single mechanism that is independent of media, routing protocol, and data protocol**
- **Requiring no changes to existing protocols**

BFD Protocol Overview

- **At its heart, Yet Another Hello Protocol**
- **Packets sent at intervals; neighbor failure detected when packets stop arriving**
- **Intended to be implemented in the forwarding plane where possible**
- **Context defined by encapsulating protocol**
- **Always unicast, even on shared media**

BFD Applications

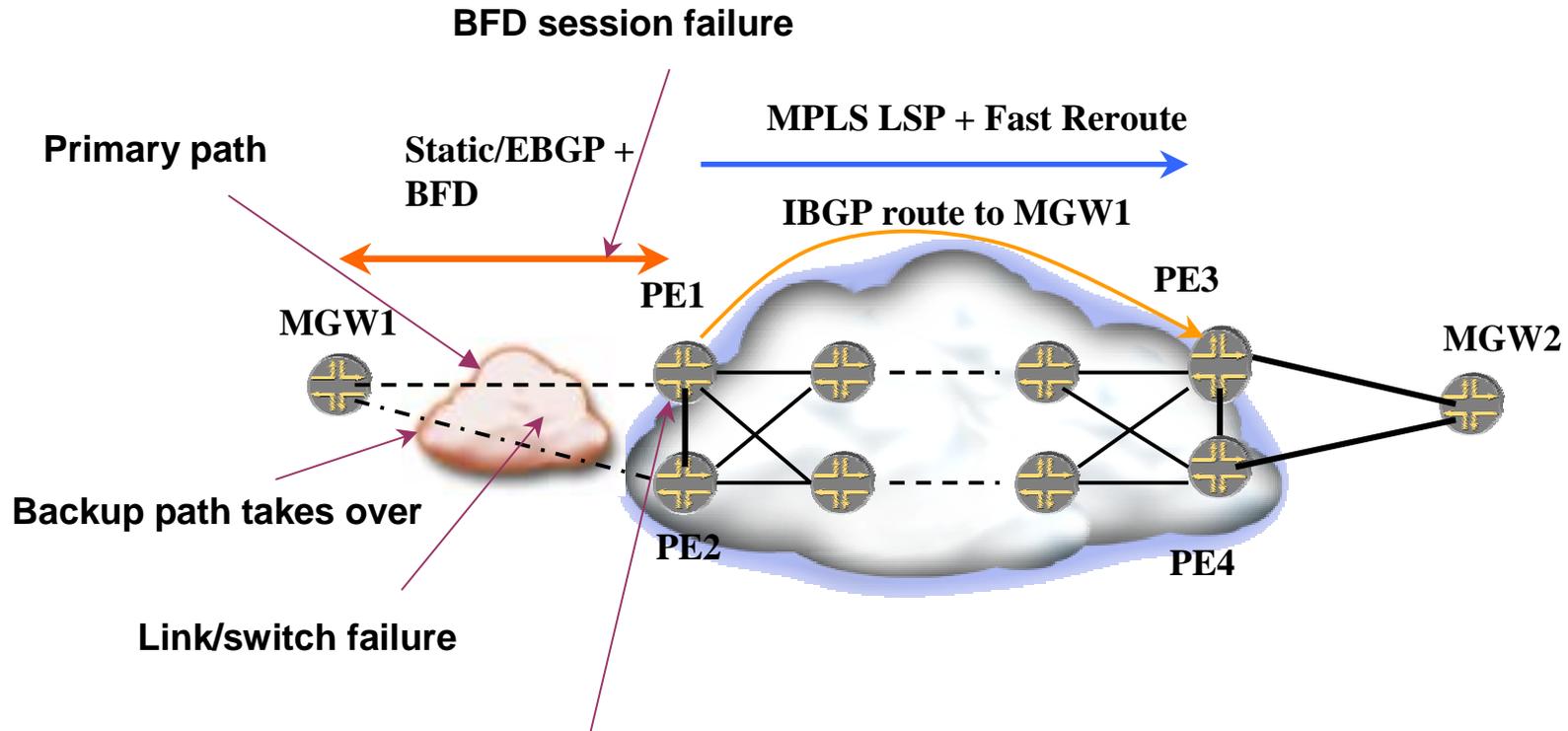
- **IGP liveness detection**
- **Tunnel liveness detection**
 - MPLS LSPs
 - IP-in-IP/GRE tunnels
- **Edge network availability**
- **Liveness of static routes**
- **Host reachability (e.g media gateways)**
- **Switched Ethernet integrity**

BFD for IGP Liveliness Detection

- One of the first motivations for BFD
- Faster convergence particularly on shared media
 - Sub-second IGP adjacency failure detection
- IGP hellos can be set to higher intervals
 - Can improve IGP adjacency scaling



BFD for Edge Availability Voice over IP



PE1 switches to a backup route through PE2 to reach MGW1

Summary

Dependability:

- ❖ Is a culture
- ❖ Has many layers
- ❖ Is business critical
- ❖ Must be designed into networks from the start

❖ Luckily:

- ❖ Vendors are providing tools for reliability
- ❖ Many architectural options from which to choose
- ❖ Also many protocols and mechanisms

Thank you!

