



APRICOT 2013
Singapore

19 February - 1 March 2013



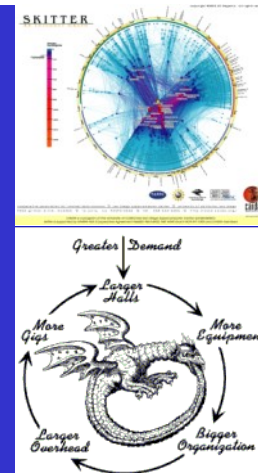
ISP and NSP Security Workshop

APRICOT 2013

Day 4

MPLS/L3VPN Security

MPLS / L3VPN Security





Before we start...

- Mainly of interest to providers/ISP/Carriers
 - Large enterprises, verticals also using MPLS internally
- To support MPLS in your network you **MUST** have:
 - Fully working IP network. If it's broken MPLS won't fix it
 - Hardware and Software support. Depends on vendors
 - Alcatel-Lucent
 - Service Routers (7705, 7210, 7450, 7750, 7950)
 - Juniper
 - M-series, T-series, J-series, E-series, MX-series
 - Cisco
 - Most platforms with CEF support (1800/2800/3600/6500/7200/7600/9000 etc)

Things I want you to know

- MPLS is a tool to solve problems
 - Not everyone has the same problems or pain
- In other words reason to deploy (choose 1+)
 - Traffic Engineering
 - Traffic Protection
 - Provider provisioned VPN's
 - Layer 3 and/or Layer 2
- Or in other words
 - Save money
 - Make money



What is MPLS?

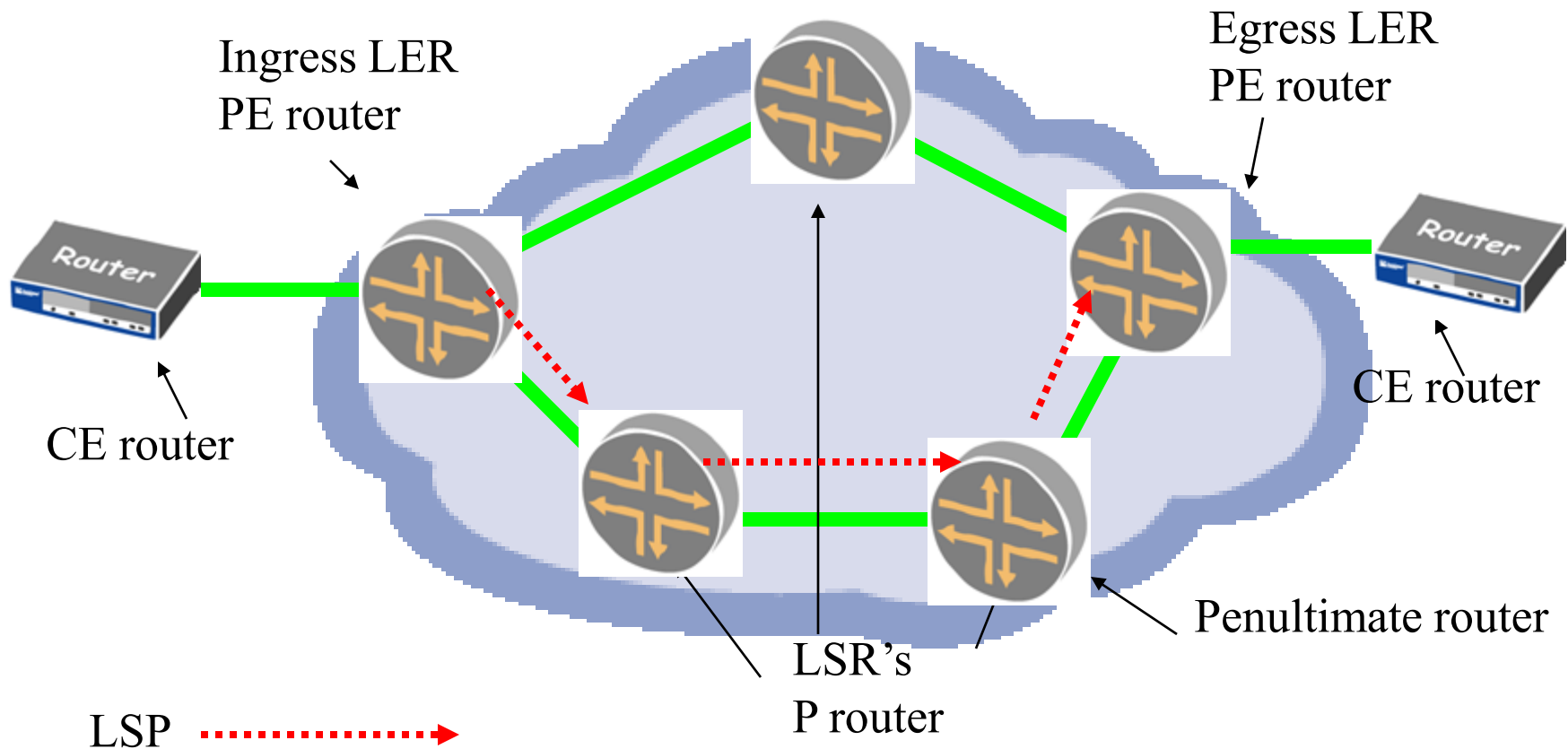
It's a tunnel!

- Multiprotocol Label Switching
- Connection Orientated Virtual Circuits over IP imple with label switching
- Grew out of
 - Cisco's Tag switching
 - Ipsilon (Nokia) IP switching
 - IBM ARIS
 - 3Com's FAST IP
- Expanding areas of application
 - Cost savings
 - New services
- Promise of Multiprotocol Unification (Core NOT edge)
- Defined by RFC 3031, RFC 3032



MPLS Terminology

- An LSP is a unidirectional flow of traffic



Push, Pop, Swap

- Push



- Pop

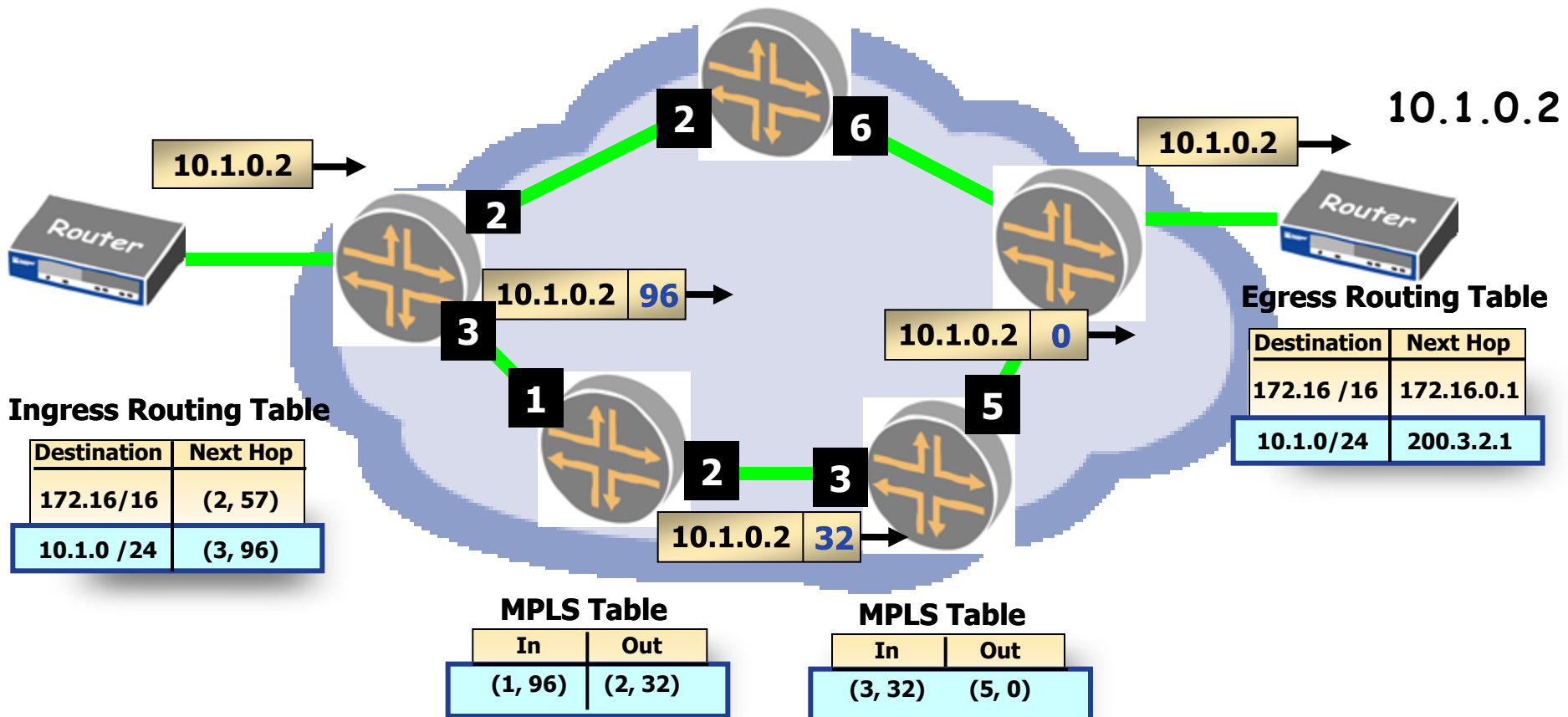


- Swap



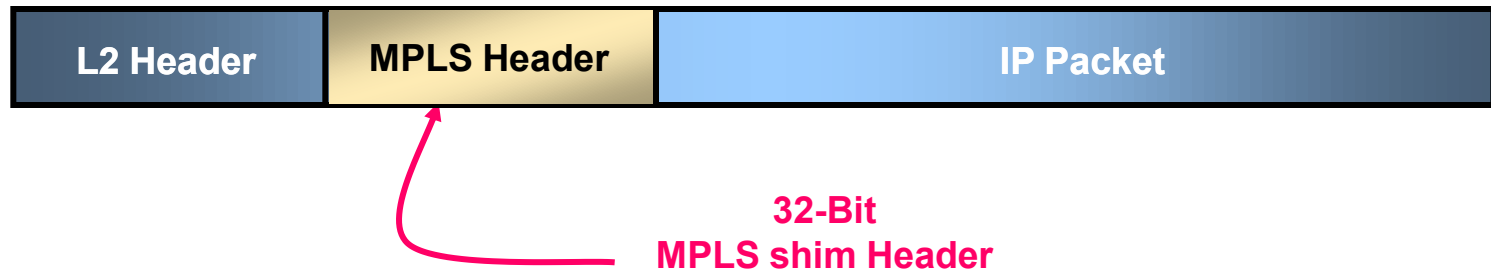
MPLS Forwarding Plane

MPLS Table	
In	Out
(2, 57)	(6, 0)



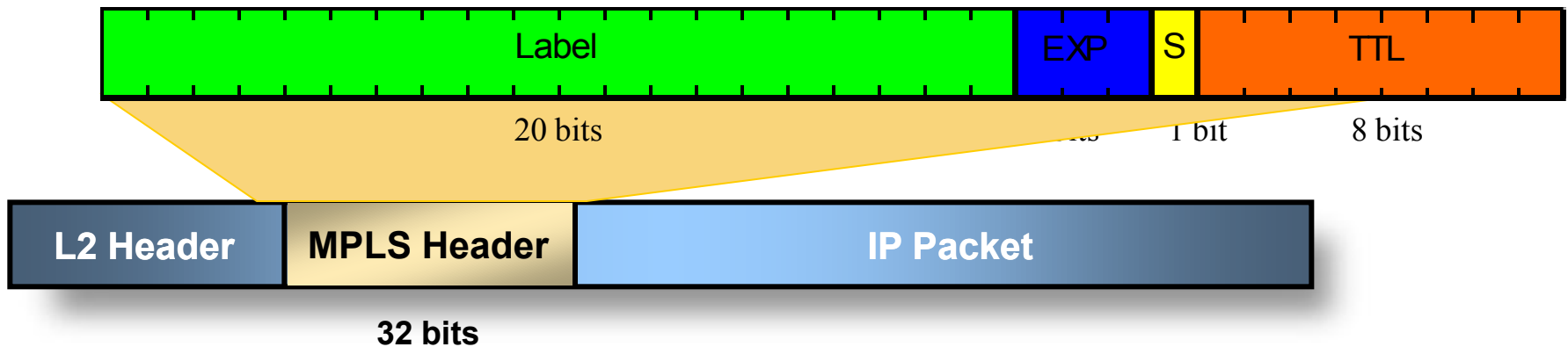
Labeled Packets

- MPLS header is prepended to packet with a *push* operation at ingress node
 - Label is added immediately after Layer 2 encapsulation header



- Packet is restored at the end of the LSP with a *pop* operation
 - Normally the label stack is popped at penultimate node

The Label



- Label
 - Used to identify virtual circuit
- EXP
 - Experimental. Currently this is used to identify class of service (CoS)
- S (Stack Bit)
 - Used to indicate if there is another label inside this packet or is it the original encapsulated data
- TTL
 - Time to live, functionally equivalent to IP TTL.

Example - Ethernet



0 0 1 0 1 1 1 1 0 1 0 0 0 1 0 1 1

My Web Page

TCP | port = 80 (www)

IP Header | Protocol = TCP

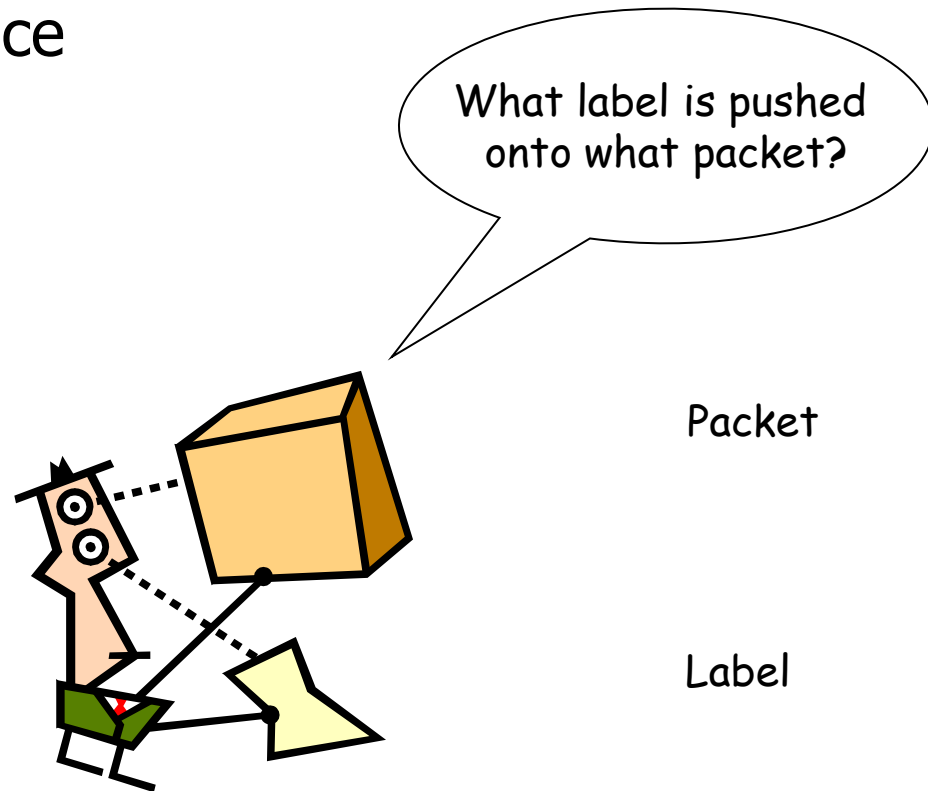
Label = 23 | EXP = BE | S = 0 | TTL = 254

Label = 47 | EXP = BE | S = 1 | TTL = 240

Dest. MAC Src. MAC Type = 8347

FEC – Forwarding Equivalency class

- All traffic with the same FEC will follow the same path and experience same level of service
- E.g. of FEC
 - Destination IP address
 - BGP next hop
 - VPN membership
 - Source address
 - Any combination of above



Signaling

- Protocols that are used to setup maintain and tear down LSP's.
- Can behave differently depending on function
- Let's describe a language / concepts to understand these differences in operation

Tell the routers what label to use on each hop!



Signalling Protocols

- LDP

- Label Distribution Protocol

- RSVP-TE

- Resource Reservation Protocol
with Traffic Engineering Extensions

- MBGP

- Multi-protocol BGP

Which you use depends
on why you are using MPLS!
Maybe you need all of them!



Which to choose...

RFC's mandate LDP
support for
L3 VPN's

- Traffic Engineering, Traffic Protection
 - RSVP
 - Link State protocol
- VPN's
 - LDP or RSVP (all LSR's)
 - MBGP (PE's only)
- Why use LDP at all?
 - Configuration scaling
 - LDP configuration is "per box"
 - RSVP configuration is "per LSP"

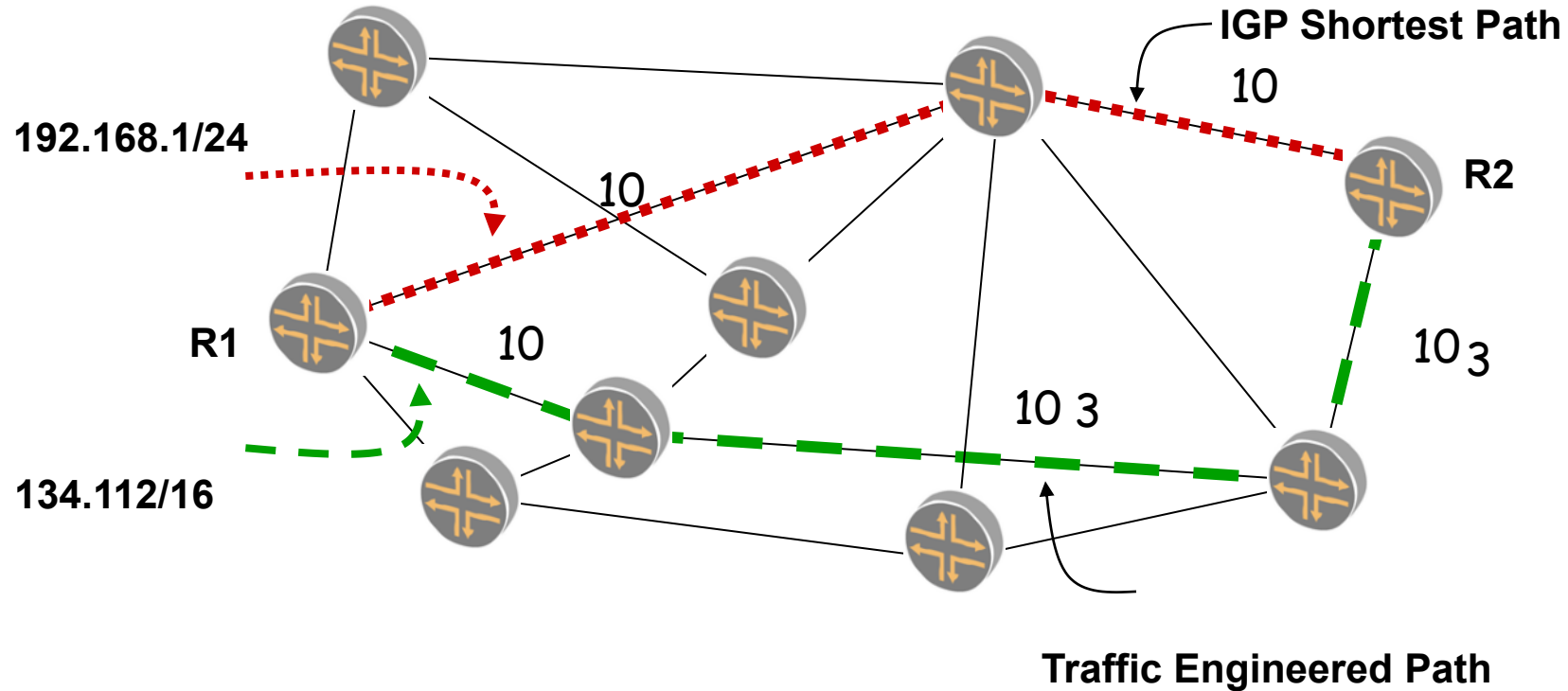




Traffic Engineering Defined

- Sub Optimal routing
- Network Engineering is putting bandwidth where the traffic is. Traffic Engineering is putting the traffic where the bandwidth is!
- To meet one of two requirements
 - To better utilize network capacity and resources.
 - To put traffic on a path that can support it's requirements
- Incorporate Traffic Protection to achieve SONET like failure recovery.

MPLS-Based Traffic Engineering





Traffic Engineering Options

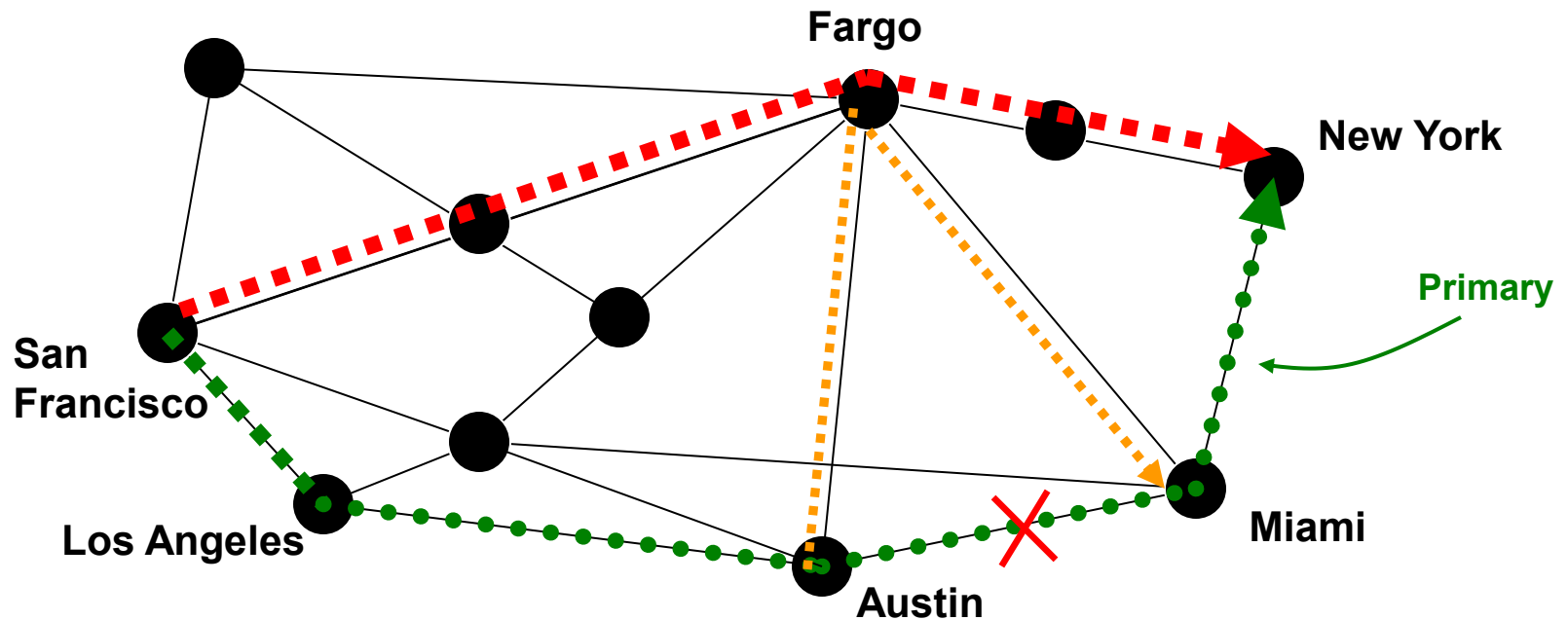
- Can we do this another way
 - IGP metrics ☹️
 - Flow = all traffic with same destination
- MPLS because
 - Granularity of flows
 - Flow = all traffic with same FEC
 - One network for all services
 - Less expensive



Traffic Protection

- Working definition
 - Reduce time of disruption
 - Reduce Packet Loss
 - “SONET like” sub millisecond recovery under failure conditions
- Can we do this another way
 - SONET/SDH
 - Lower IGP timers
- MPLS because
 - No extra capital – config change only
 - Pick which traffic needs it
 - One network for all services
 - Less expensive

Traffic Protection – example





Traffic Protection Variations

- Fast reroute
- Link Protection
- Link-Node Protection



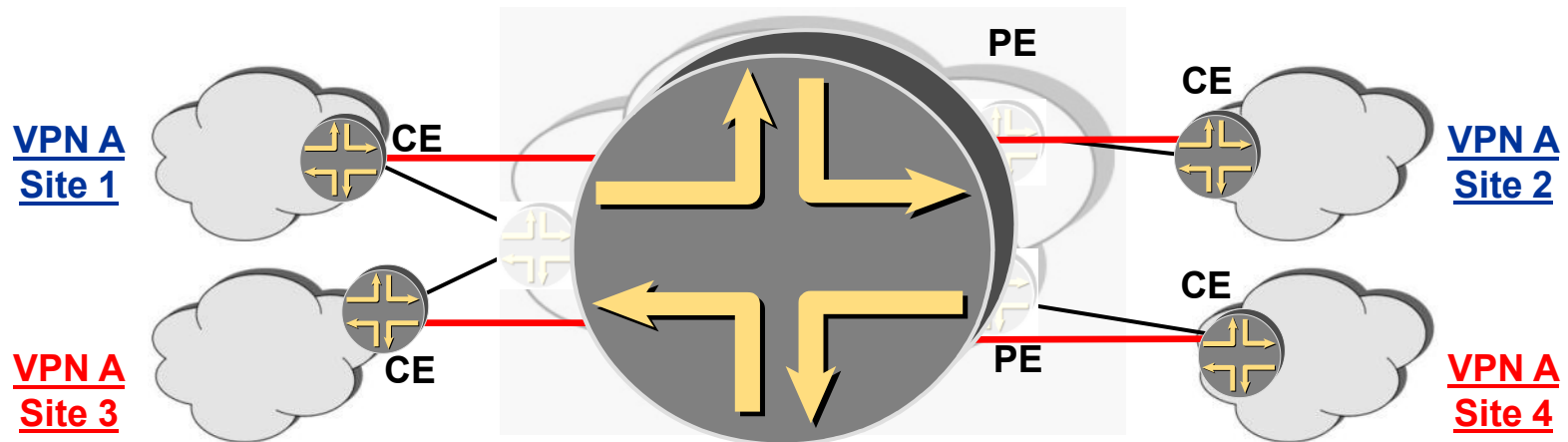
Layer 3 VPN (4364 BGP/MPLS VPN)

Provider provisioned VPN

- ISP runs backbone for customer
 - Customer can be another ISP!
- Attractive to
 - Customer who do not want to run their own backbone
- Not attractive to
 - Customer who doesn't trust carrier
 - Customers who's jobs are threatened

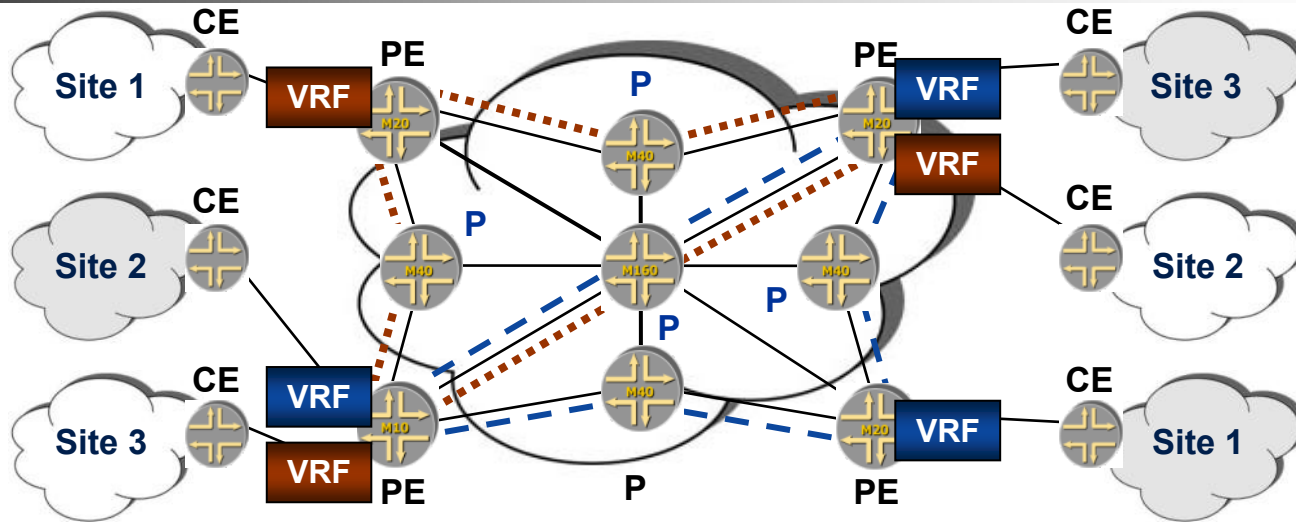
Customer View of L3VPN

- Make the cloud look like a router
- Single site provisioning



Layer 3 PP-VPNs: RFC 4364 (1 of 2)

Service Provider Network



■ Application: Outsource VPN

- PE router maintains VPN-specific forwarding tables for each of its directly connected VPNs
- Conventional IP routing between CE and PE routers
- VPN routes distributed using MP-BGP
 - Uses extended communities
- VPN traffic forwarded across provider backbone using MPLS



Layer 3 PP-VPNs: RFC 4364 (2 of 2)

- LDP or RSVP is used to set up PE-to-PE LSPs
- MP-BGP is used to distribute information about the VPN
 - Routing and reachability for the VPN
 - Labels for customer sites (tunneled in PE-PE LSP)
- Constrain connectivity by route filtering
 - Flexible, policy-based control mechanism



L3 VPN Options

- Can we do it another way
 - Separate Physical routers
 - Separate Logical Routers
- MPLS because
 - Scaling
 - Single site provisioning
 - Less expensive

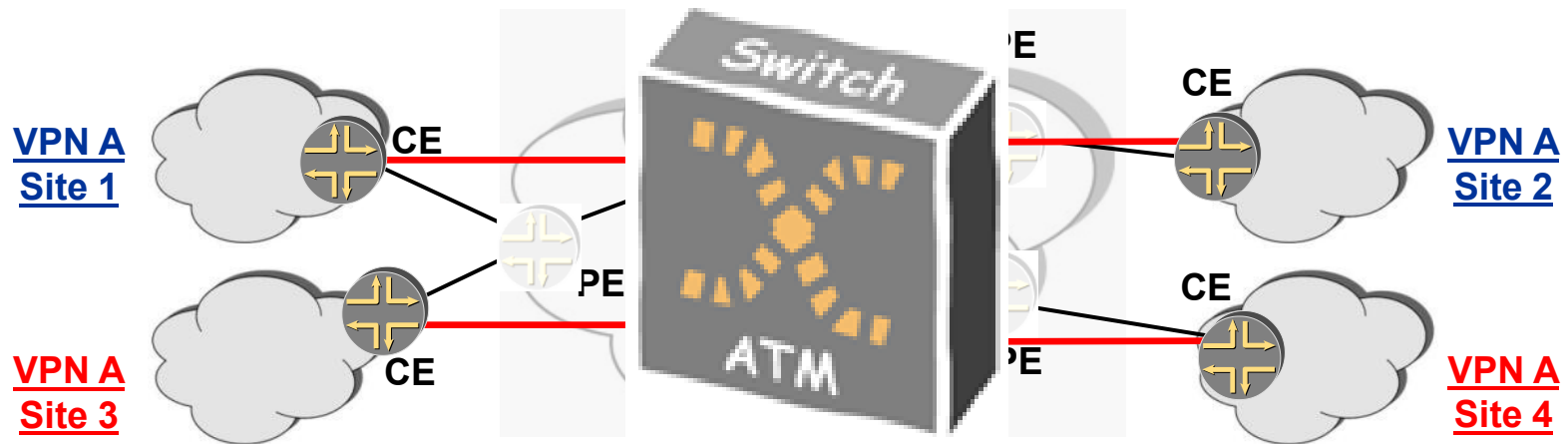


Layer 2 VPN's

- Provider provisioned VPN
 - ISP runs backbone for customer
 - Customer can be another ISP!
- Attractive to
 - Customers who want to preserve current CE technology
 - Customers who don't trust provider with L3
 - Carriers who want to offer another service
- Not Attractive to
 - Customers who do not want to run their own backbone

Customer View of L2VPN

- Make the cloud look like a ATM/FR network





L2 VPN Options

- Can we do it another way?
 - Traditional ATM/FR/leased line infrastructure
- MPLS because
 - One network for all services
 - Less expensive
 - Scaling
 - Single site provisioning *

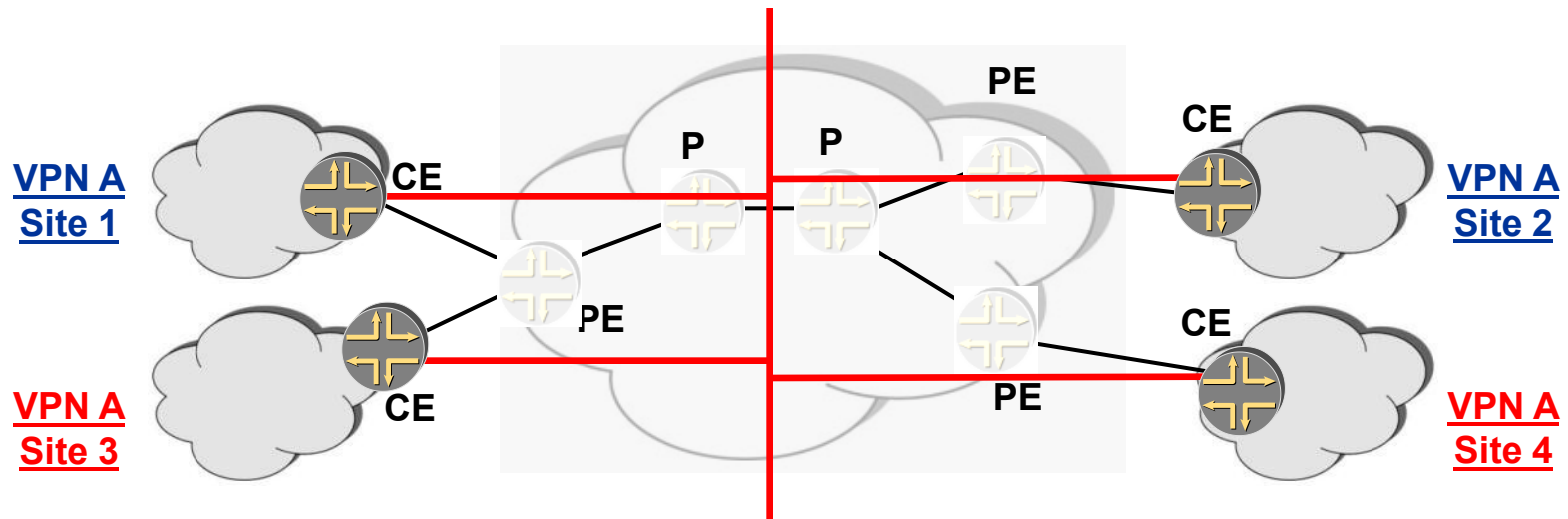


VPLS

- Virtual Private LAN Service
- Attractive to
 - Customers who like ethernet as CE
 - Lots of locations close together with 'high' WAN bandwidth requirements (kiosks)
 - No routing required
- Not attractive to
 - Customers who like control and visibility of core. "what can I ping to identify fault-domain?"
 - Controlling broadcasts

VPLS

- Make the cloud look like an ethernet switch



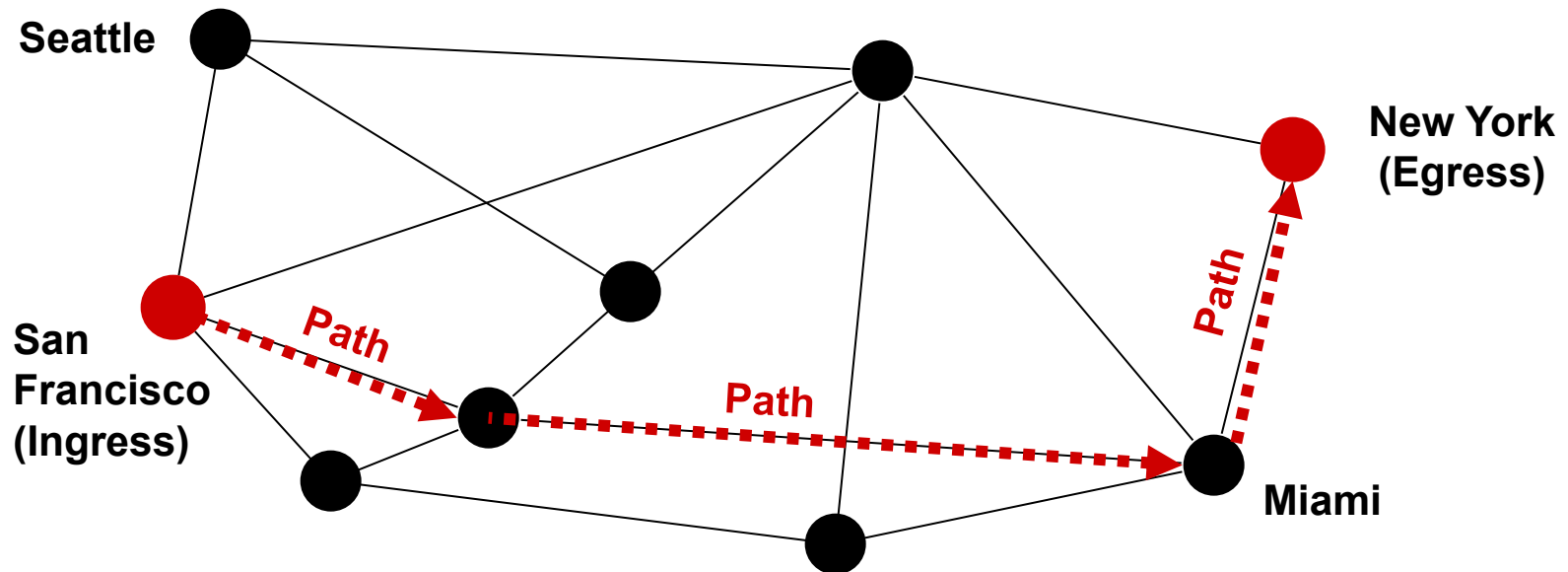


VPLS Options

- Can we do it another way?
 - Separate physical switches tying all customer sites
 - VLAN's over layer 2 backbone
- MPLS because
 - Scaling
 - One network for all services
 - Less expensive

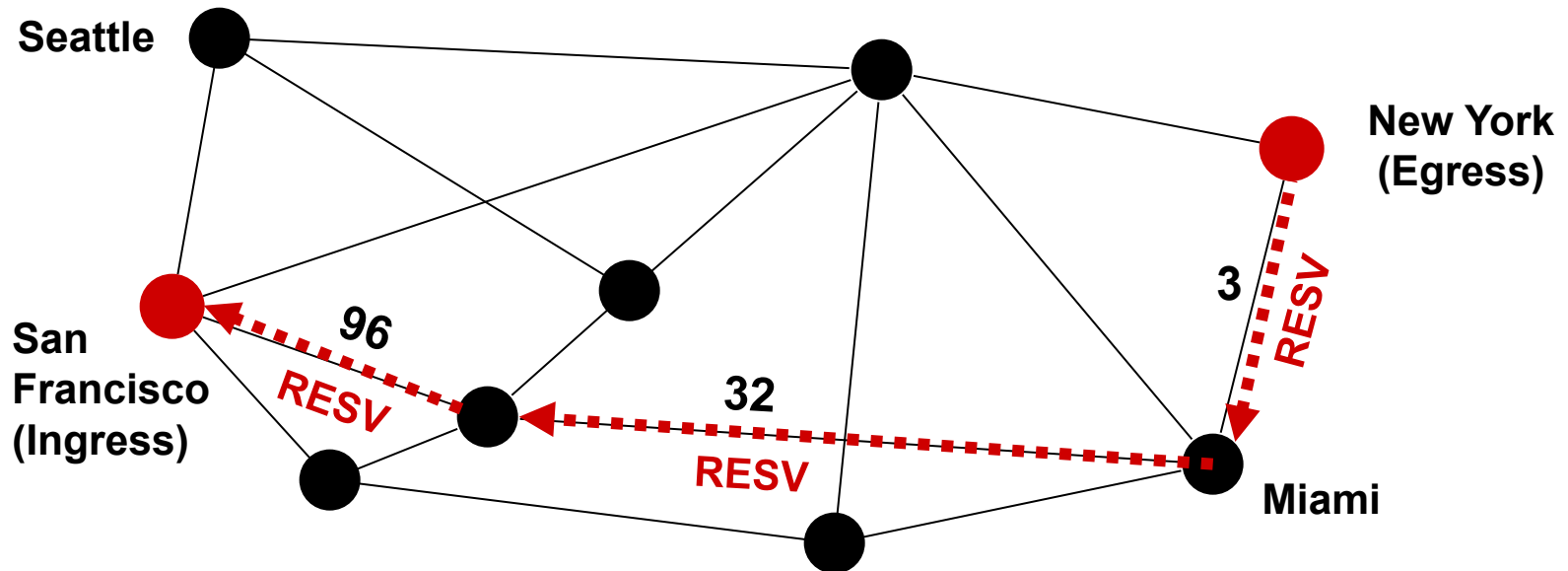
RSVP Signaling Example: Path

RSVP sets up path from San Francisco to New York



RSVP Signaling Example: Reservation

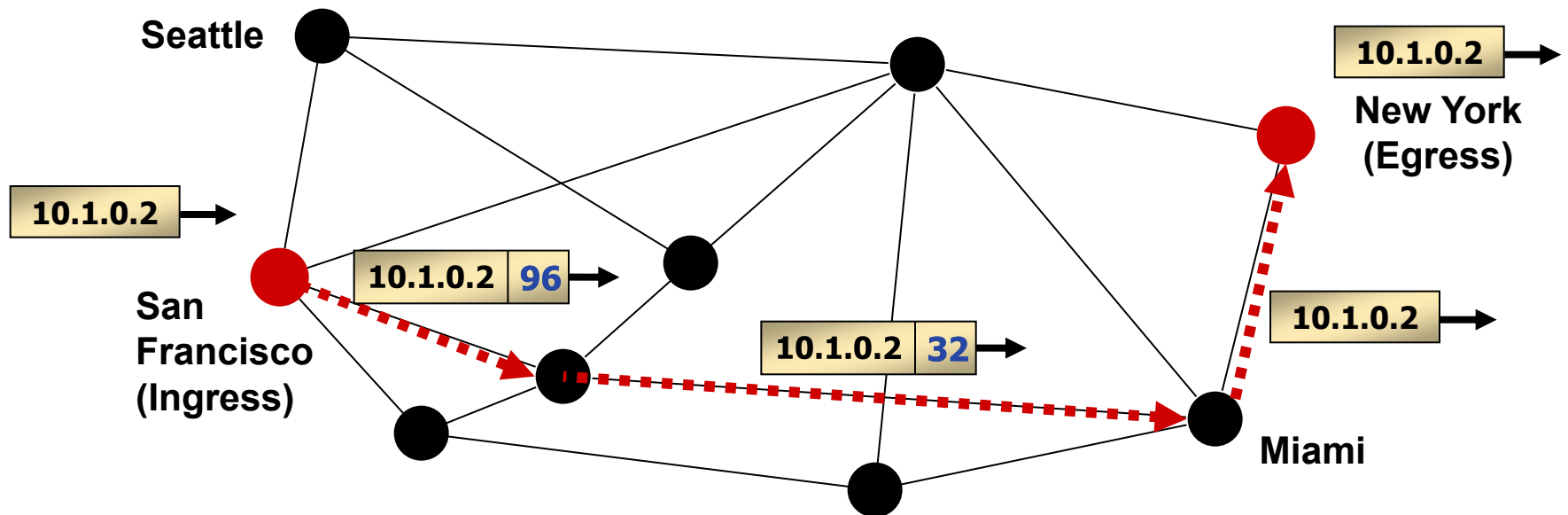
- The resv message visits each router on the path in reverse order
 - Labels assigned hop to hop in the upstream direction



LSP Established!

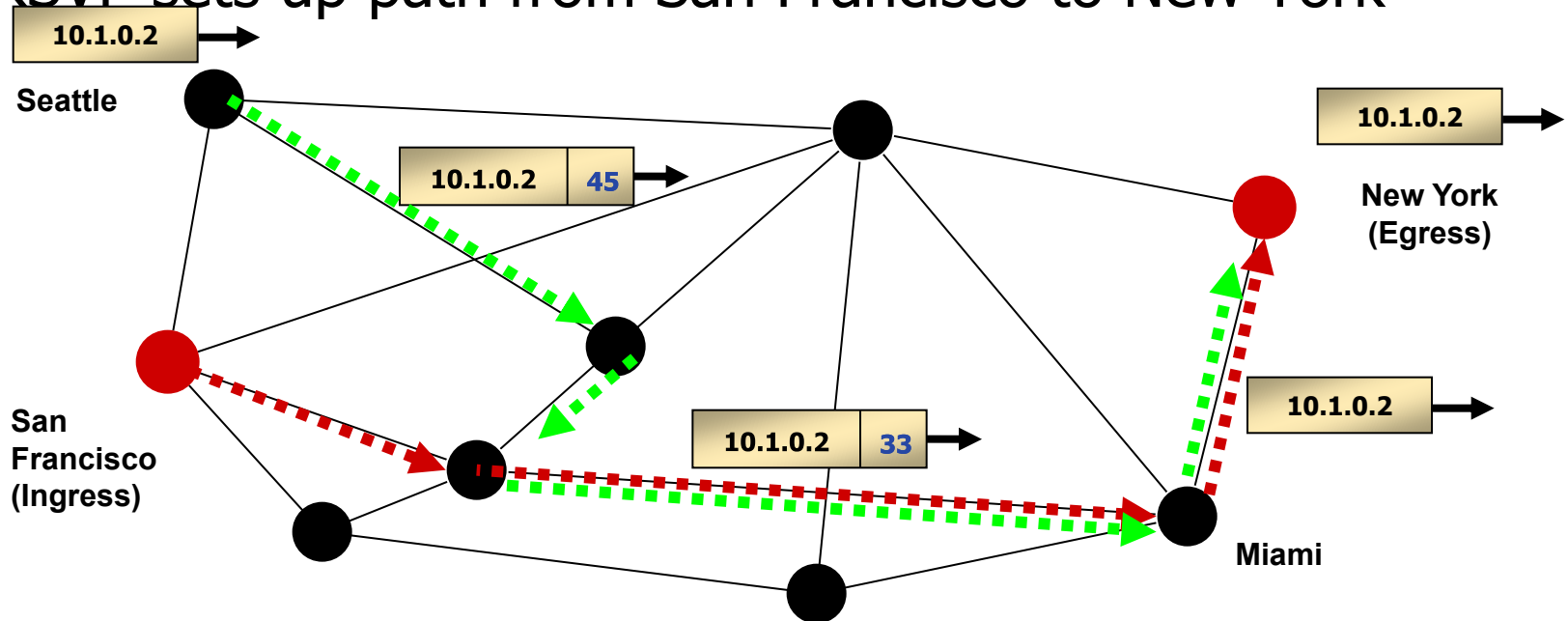
RSVP Signaling Example: Forwarding

RSVP sets up path from San Francisco to New York



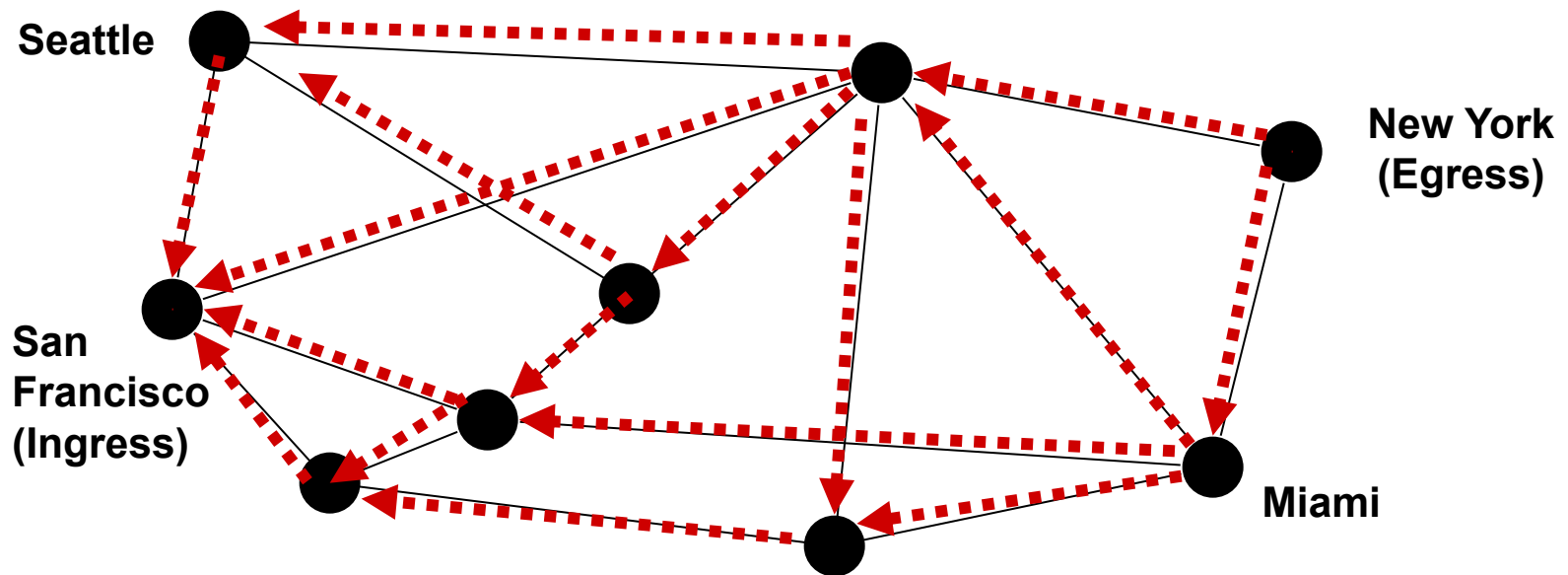
RSVP Signaling Example: Forwarding 2

RSVP sets up path from San Francisco to New York



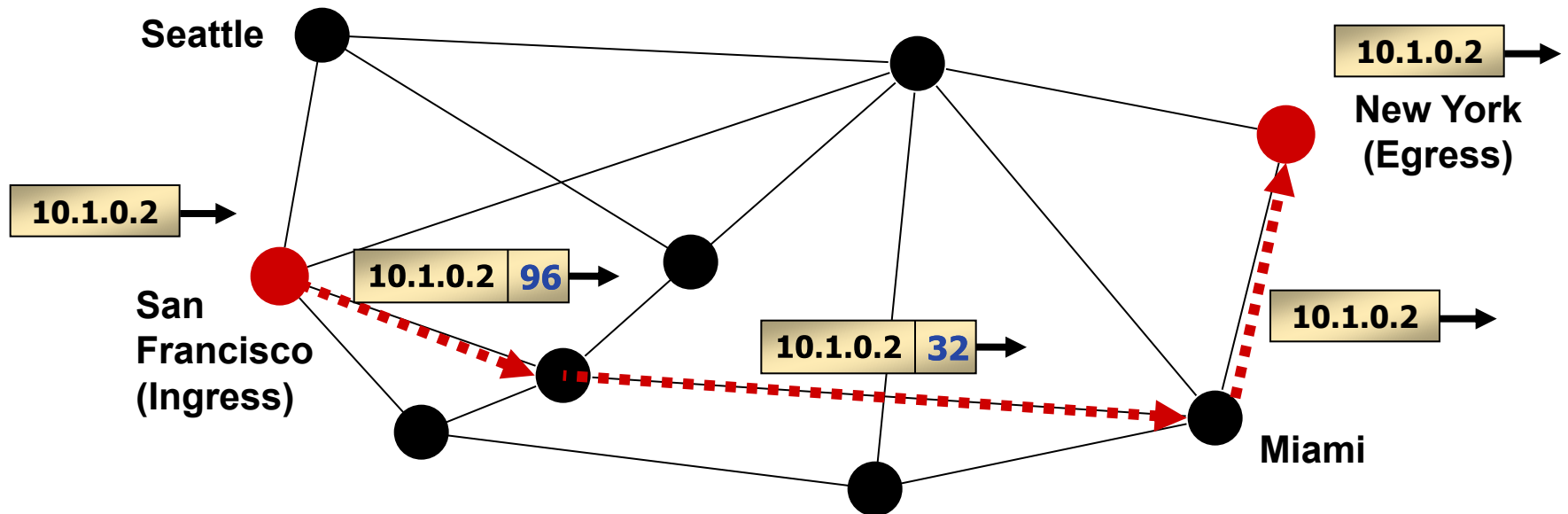
LDP Signaling Example: Label Binding

- Label Mappings are made for entries in the routing table
 - Labels assigned hop to hop in the upstream direction



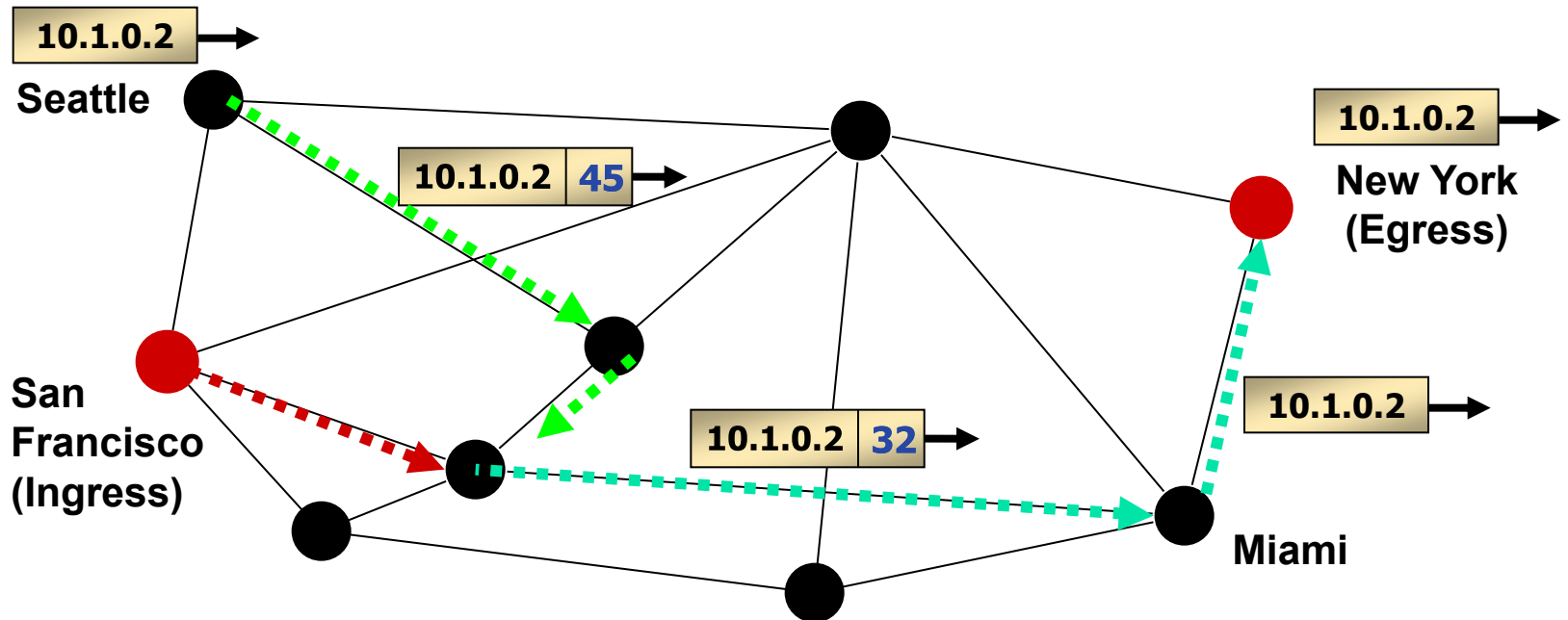
LDP Signaling Example: Forwarding

LDP path available to egress



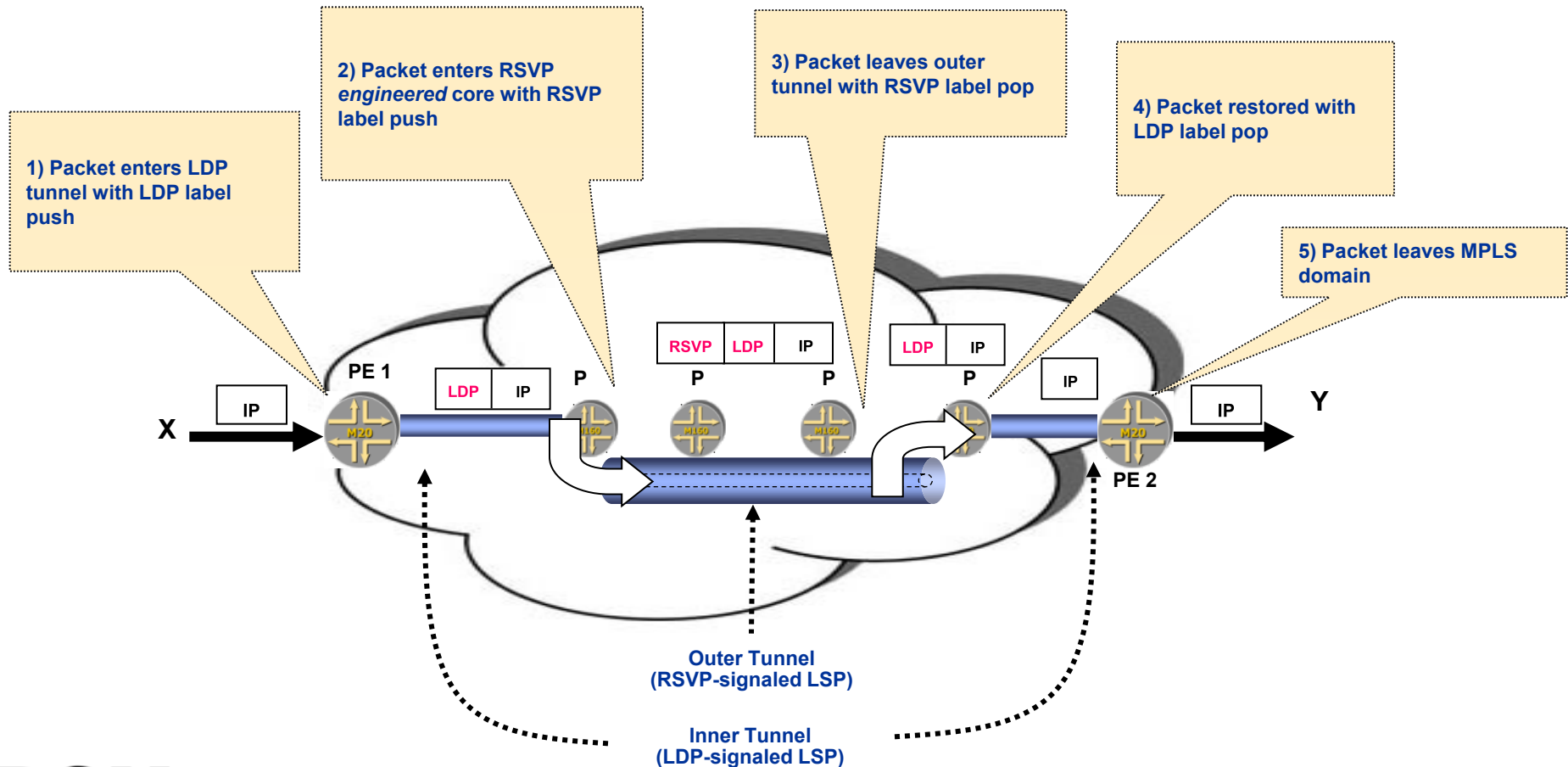
LDP Signaling Example: Forwarding 2

LSP Merging occurs



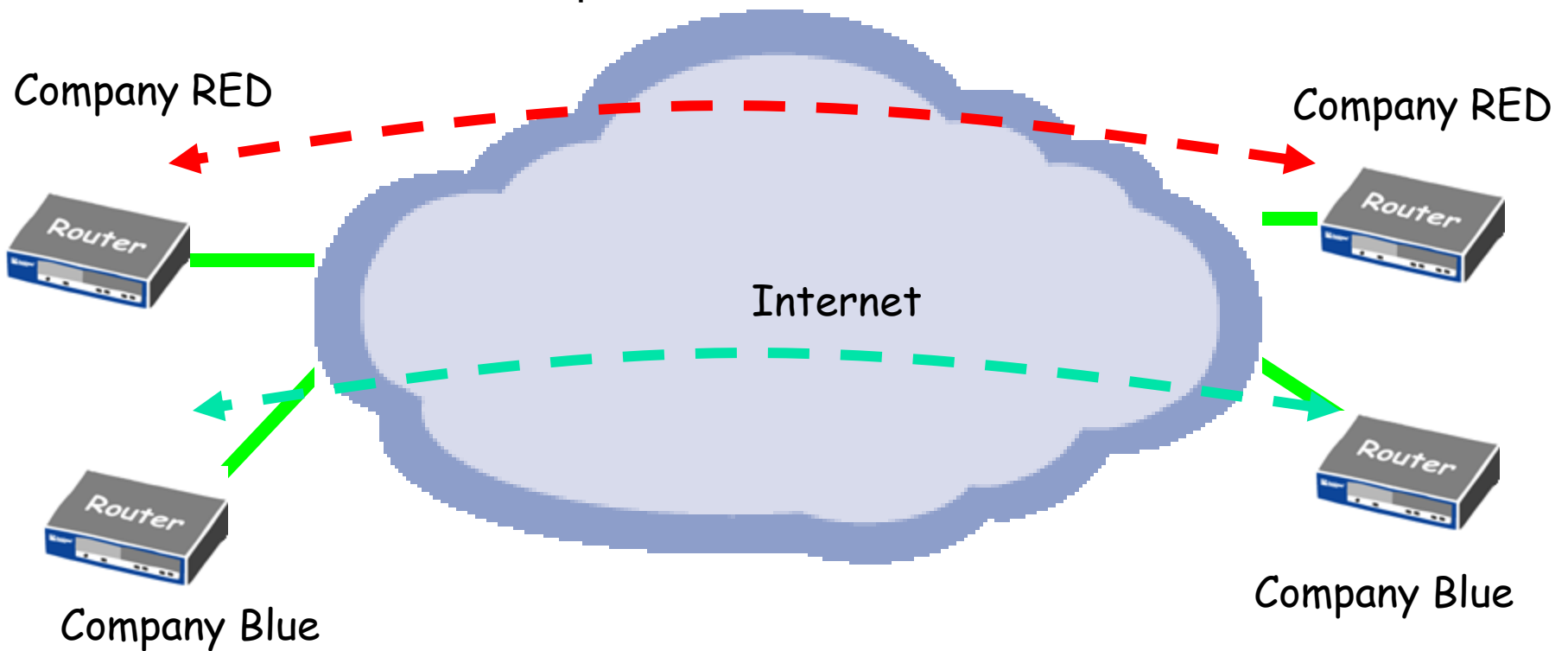
Label Stacking

- Label stacking improves scalability
 - Similar to ATM's VP and VC hierarchy



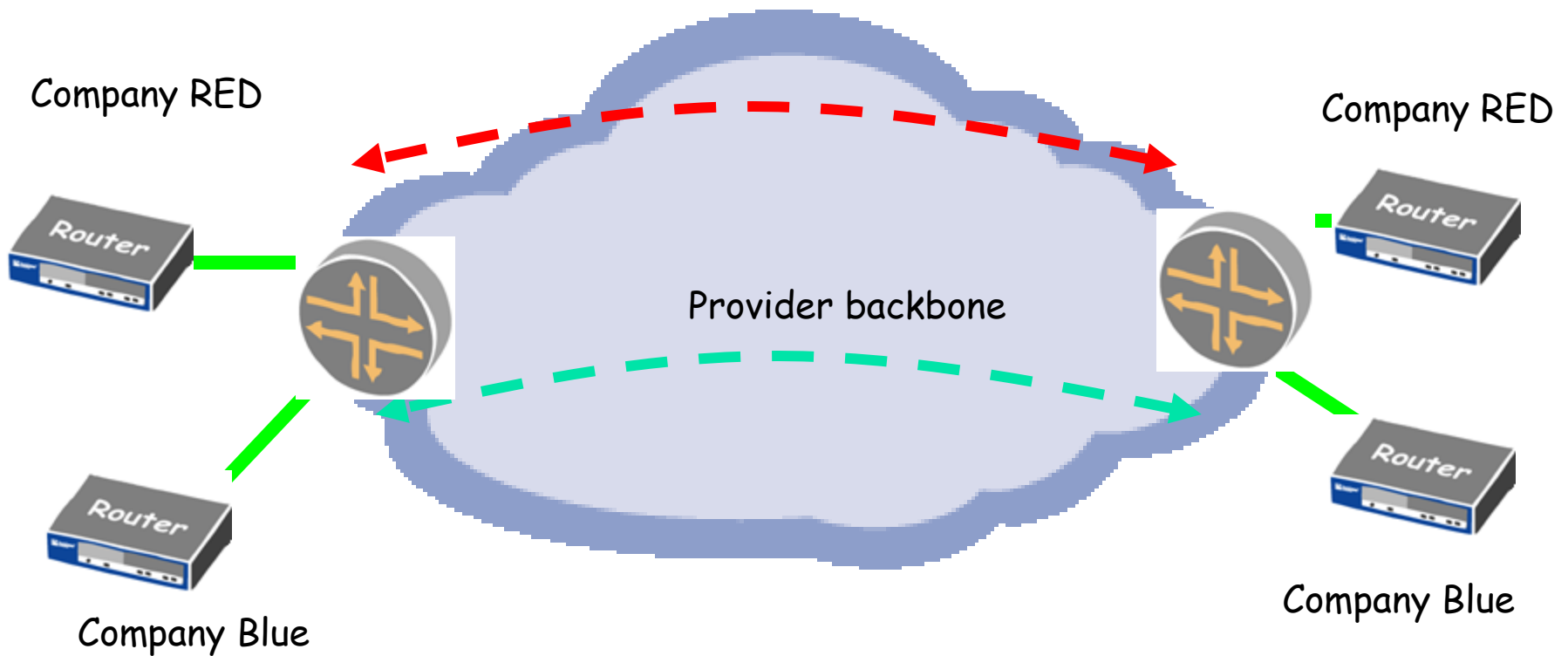
Traditional VPN's

- CPE based
- Customer controlled
- No value add for provider



Provider provisioned VPN's - PPVPN

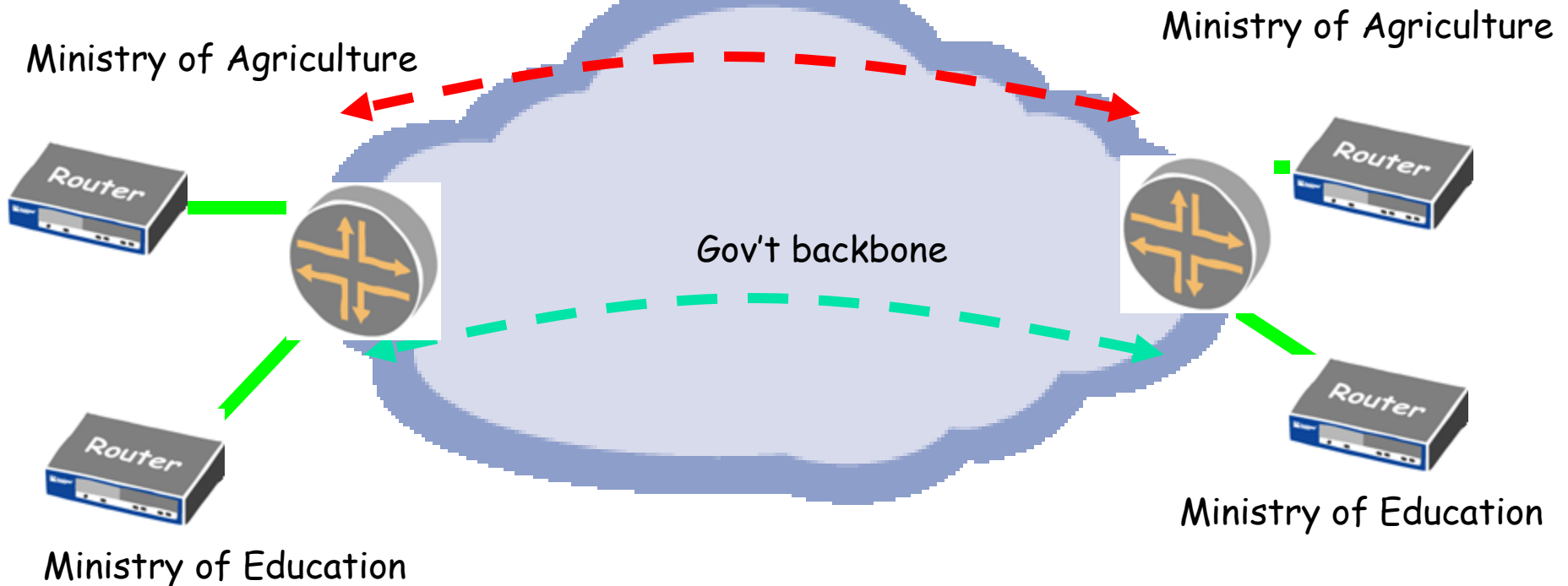
- PE based
- Customer outsource backbone
- Value add for provider
- Single Site Provisioning (BGP, + Route refresh + Route Target Filtering)



Sharing Network backbones

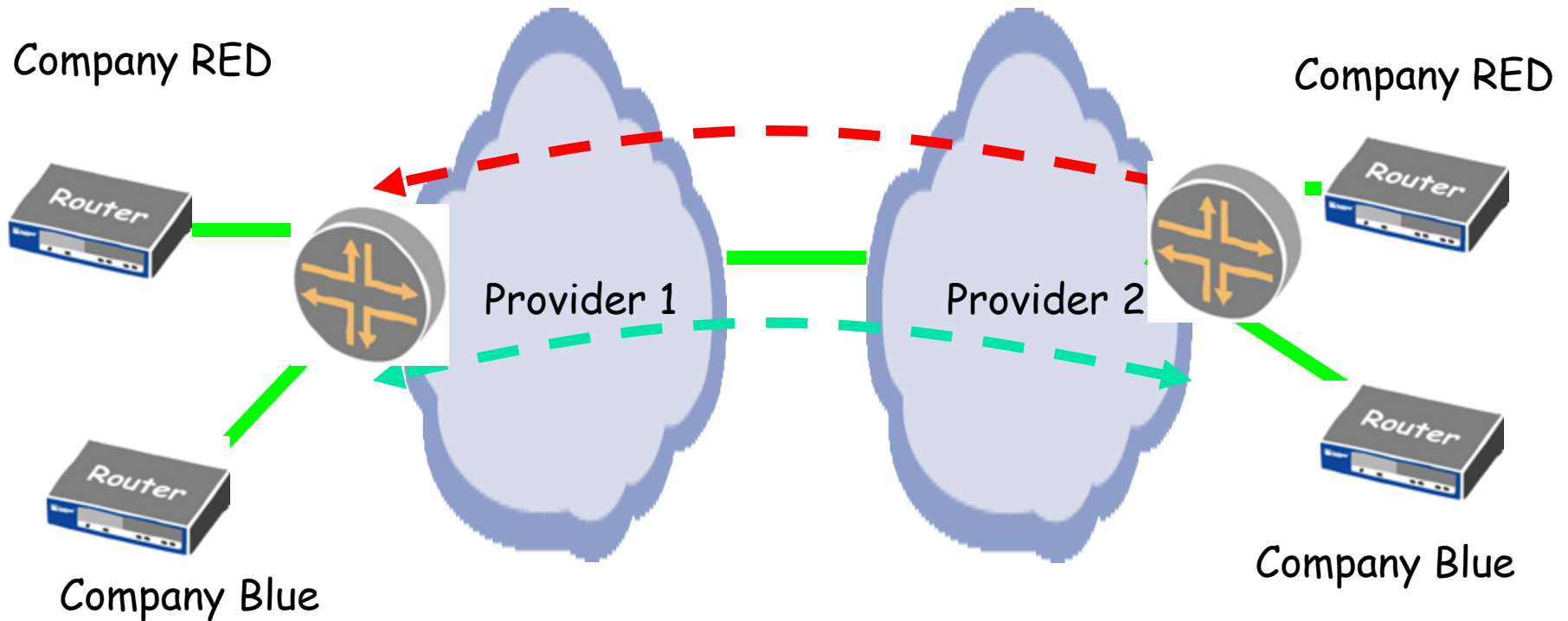
- Infrastructure built by one department
- Shared by other departments
- Cost effective government spending

- Examples
 - Gov't backbones
 - Industry Aligned



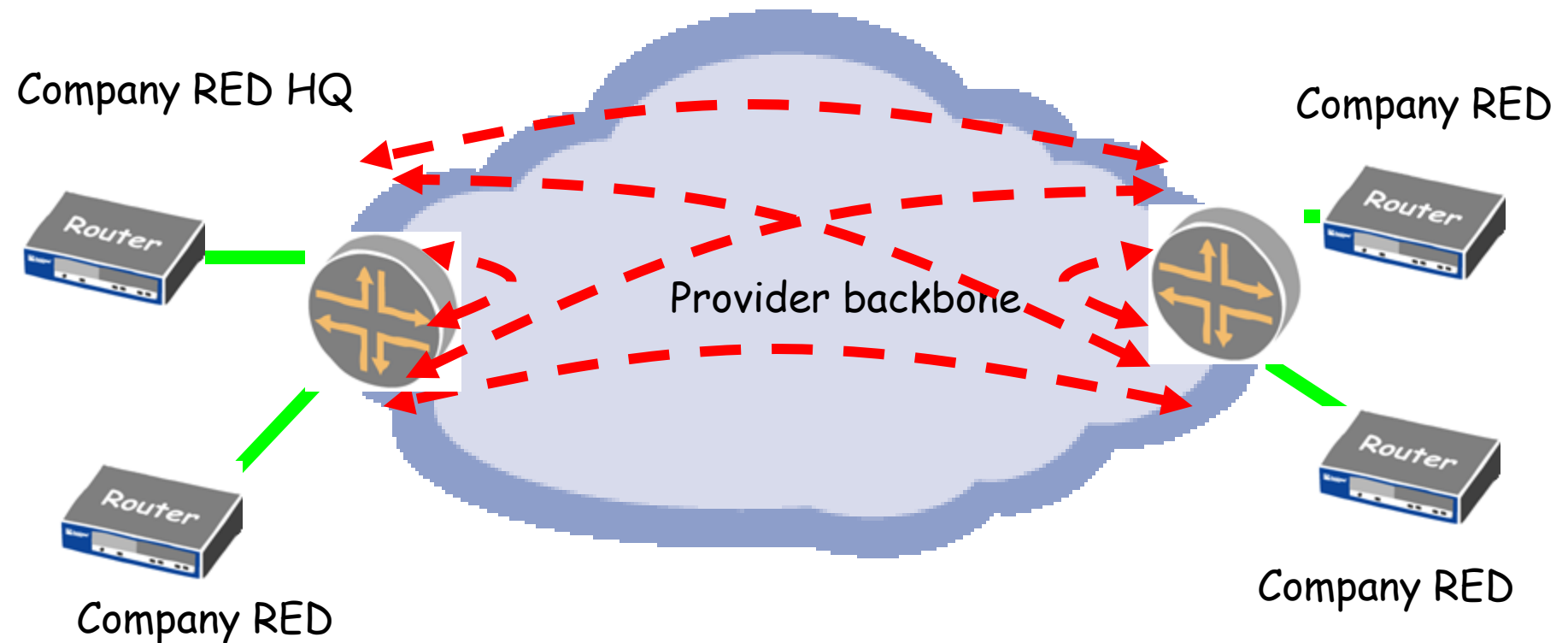
InterAS VPN's

- Requires Co-operation
- Opportunity for global coverage



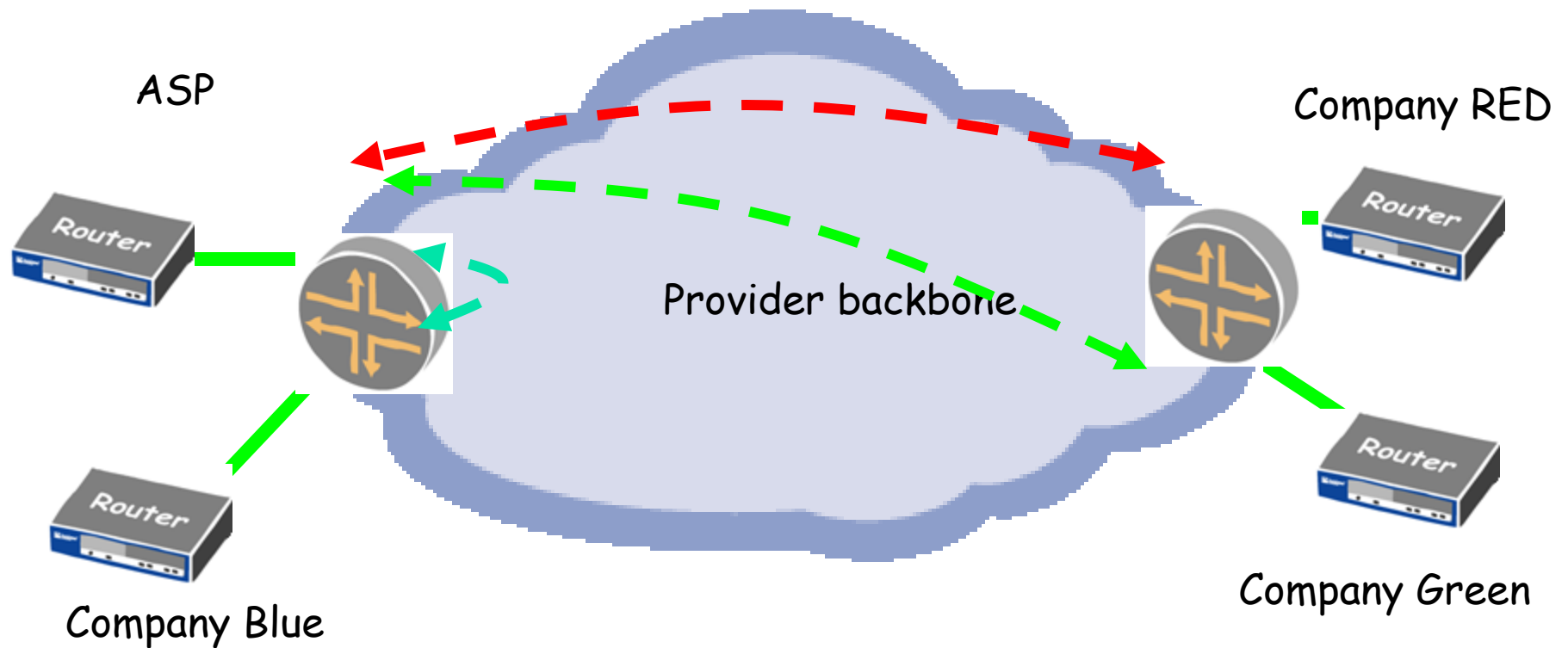
Site Connectivity

- Partial or Full Mesh is supported
- Full Mesh is more cost effective and competitive with traditional solutions



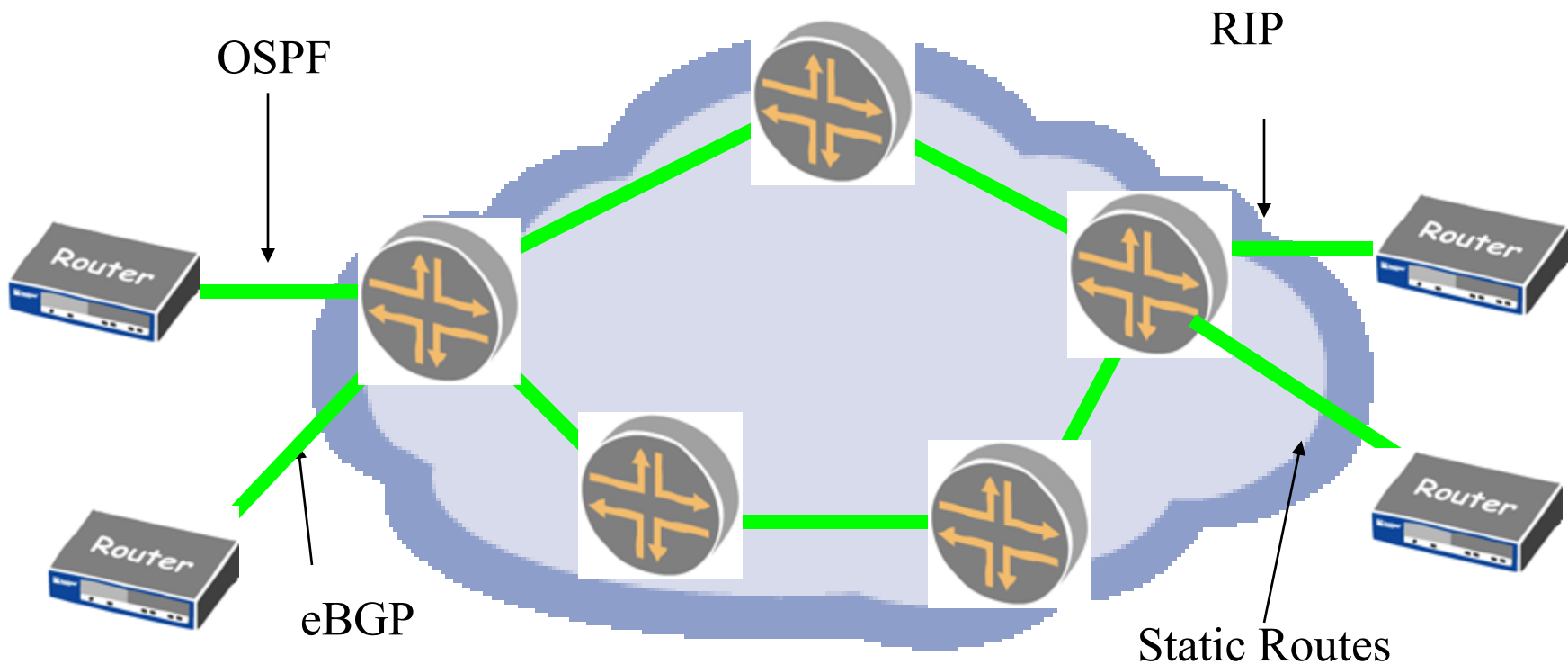
Overlapping VPN's

- Suites application / service providers



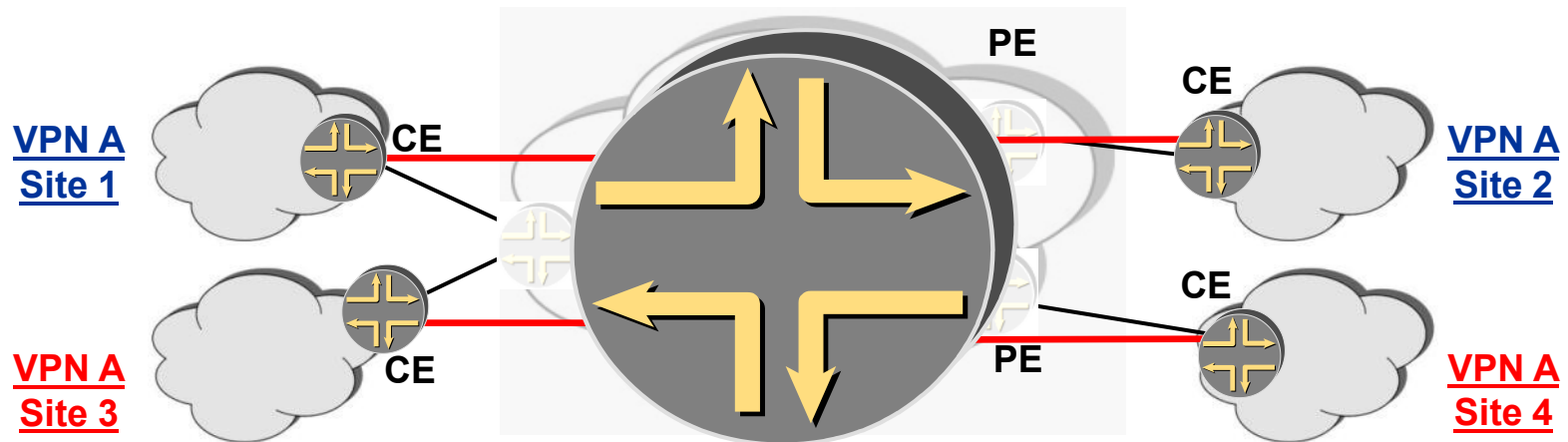
CE-PE interaction

- Any L2 connection, Any routing protocol
- CE peers at layer 3 with PE



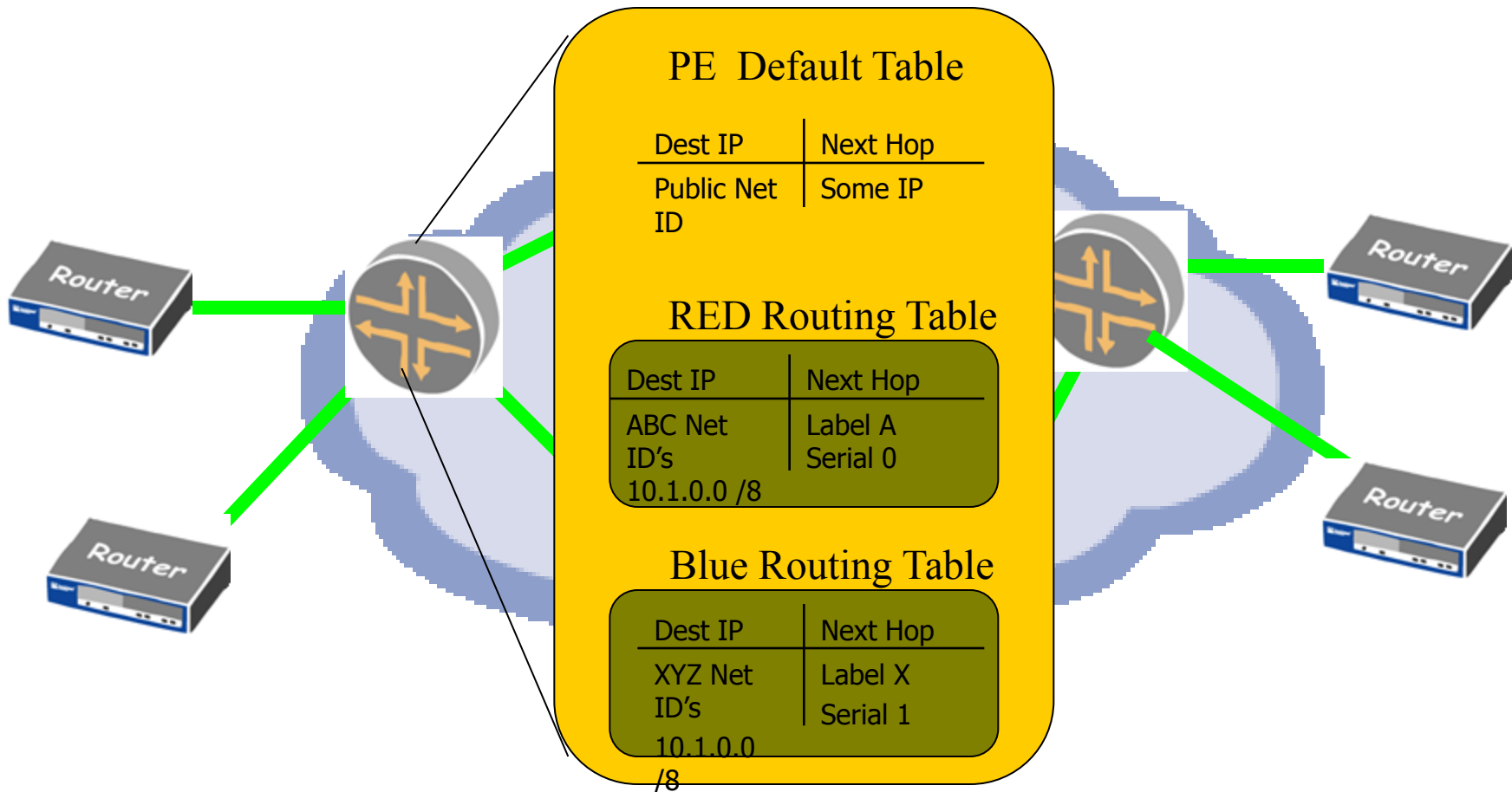
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- Single site provisioning



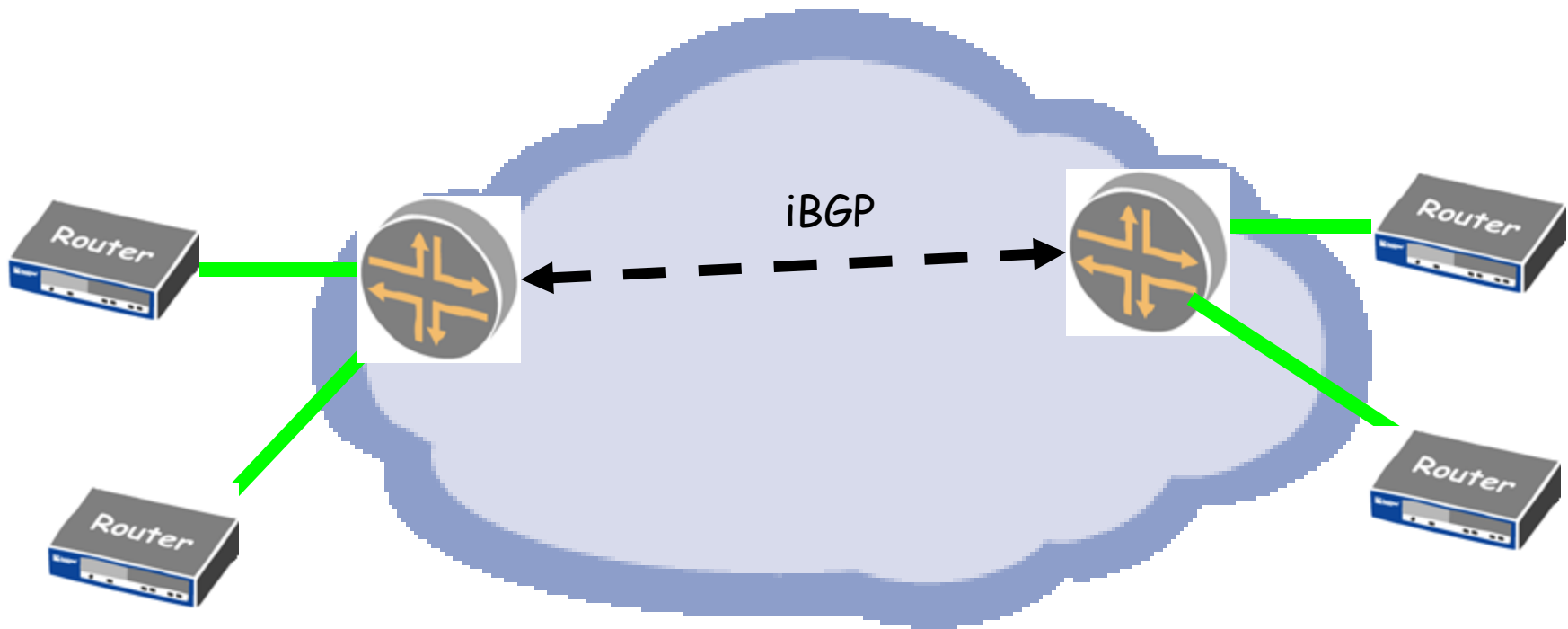
VRF – Virtual Routing and Forwarding instance

- VRF per VPN on PE
- Logical Interface packet arrives on defines the VRF used



PE-PE interaction

- iBGP between PE's carries routing information
- Assigns label per VPN



Route Distinguishers

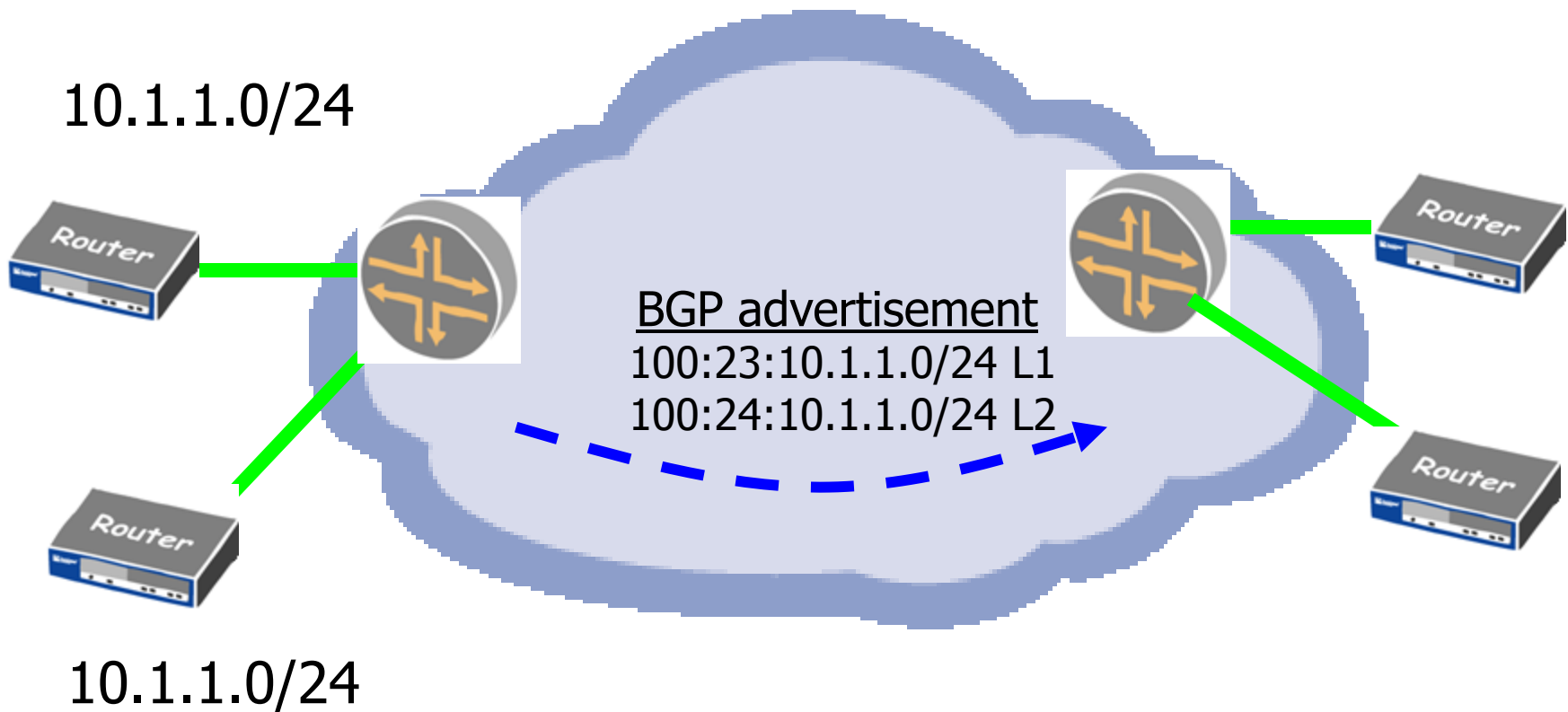
RD's have nothing to do with defining VPN membership

- Used to disambiguate possibly duplicate routes from VRF's
 - i.e. guarantee unique addressing space
 - AS:nn e.g. 100:23
 - IPv4:nn e.g. 192.168.1.1:23
- Creates a guaranteed unique address th BGP can advertise in a single database
- VPNIPv4 addresses



RD's in action

- Per VPN via BGP label assignment
- PE – PE set up via LDP or RSVP (saves state)



Route Targets

RT's tell you
which routes go into
which VPN's

- PE receives VPN IPv4 NLRI's
- Routes then placed into VRF based upon
 - Extended BGP community,
 - AS:nn 100:45
 - IPv4:nn e.g. 192.168.1.1:45
- A route may have one or more RT





Route Targets in action

- When routes are advertised, they are exported with one or more RT's
- A VRF can import routes with matching RT's
- Security of this architecture depends on YOUR provisioning integrity

Why RD's and RT's?

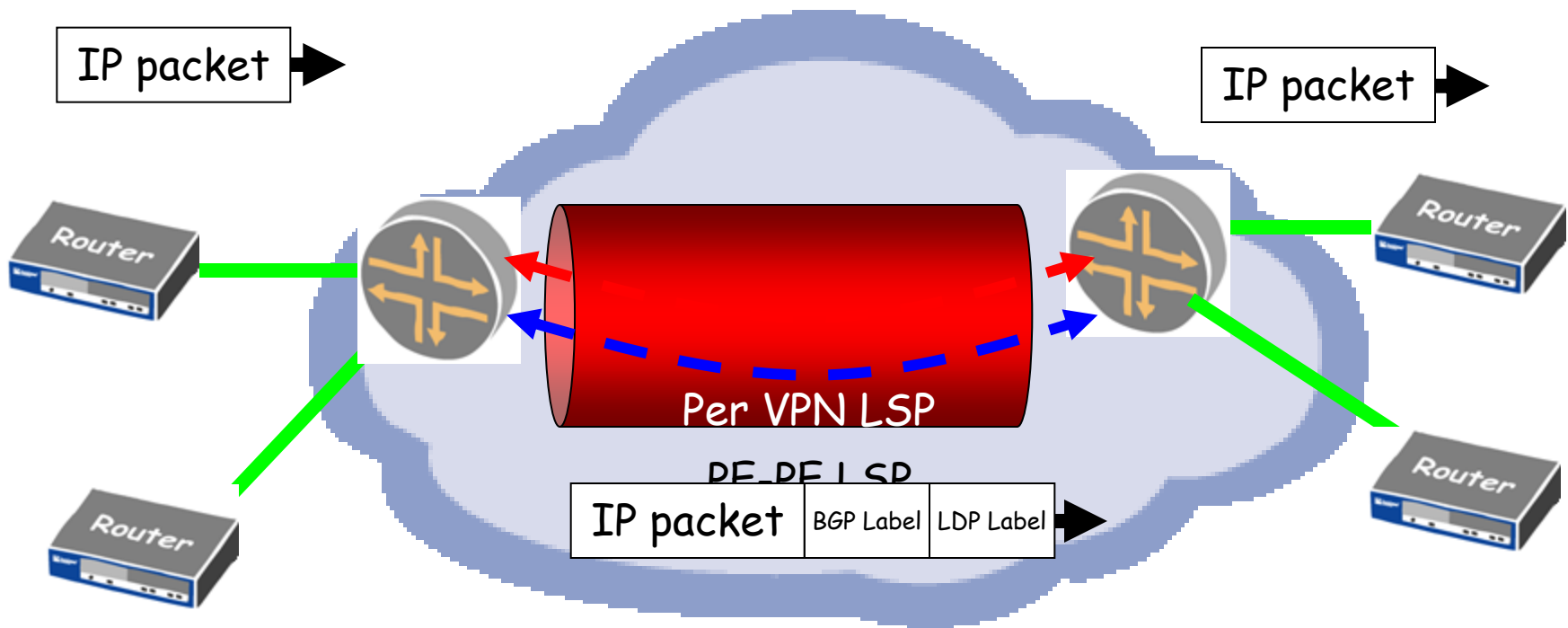
RT's tell you
which routes go into
which VPN's

- Overhead is better when
 - Advertisements get bigger, as opposed to
 - More advertisements
- Allows for overlapping VPN's
- Can be the same
 - But don't lock yourself in



LSP establishment

- Per VPN via BGP label assignment
- PE – PE set up via LDP or RSVP (saves state)



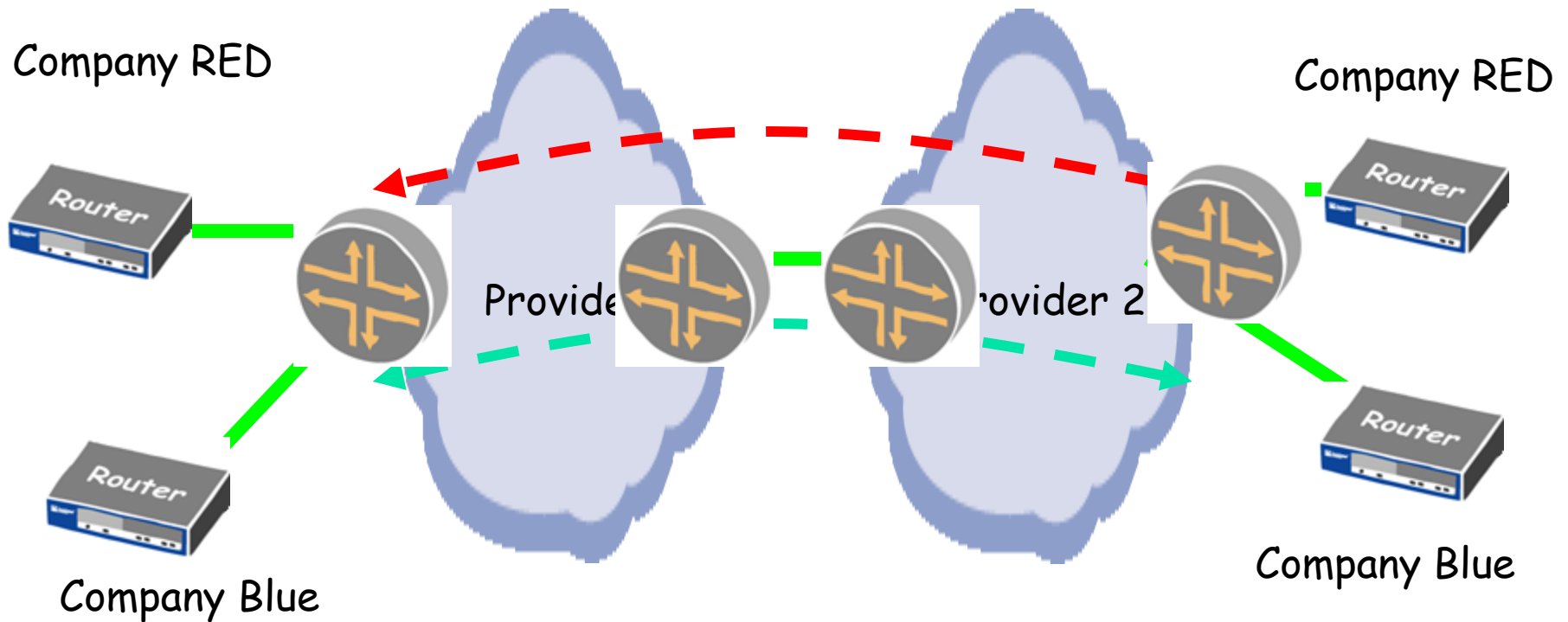


Connectivity

- Hub and spoke
 - Outsourcing internet access and Applications
- Full Mesh

InterAS VPN's

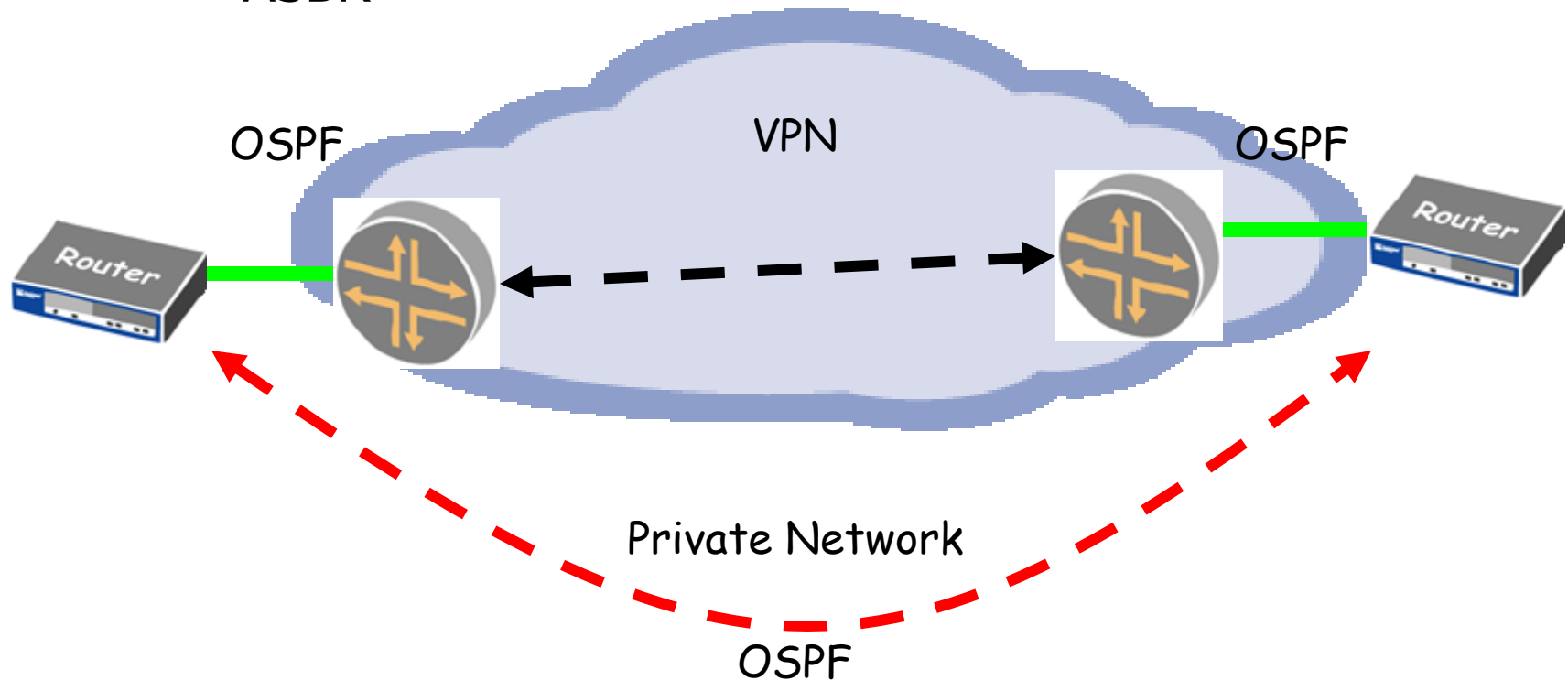
- VRF-to-VRF
- MBGP between ASBR (not OSPF)
- MBGP between PE's



VPN as backup

Do you want PE to appear as

- Intra Area Router (Sham Links)
- ABR
- ASBR





Issues

- BGP scaling
 - RR, often separate from IP RR
- Inter-AS scaling
 - MBGP between PE's is desirable
- Management
 - Usual MPLS, OAM, root cause automation.
 - Overlap NOC with VPN? Addressing?
- QoS
 - Carriers mapping 4+ queues



Security

- Routing protocol security is just as critical in an MPLS VPN environment as it is in the Internet environments
- PE-CE security is important
 - RIP
 - OSPF
 - BGP
- When using BGP, we need to ensure customers are not injecting malicious communities to hijack VPN traffic
- Labels
 - Label injection exploits are possible
 - Malicious PE injects traffic with a label for a different service
 - Misconfigured PE may accept labeled traffic from a CE



Configuring L3VPN's



Enable MPLS and LDP

```
ip cef
mpls ip
mpls label protocol ldp
!
interface fast 0/1
mpls ip
mpls label protocol ldp
!
```



PE-PE MP-IBGP Peering

- PE-to-PE MP-IBGP sessions require VPN-IPv4 NLRI

!

```
router bgp 150  
neighbor 192.168.16.1 activate
```

!

```
address-family vpnv4  
neighbor 192.168.16.1 activate  
neighbor 192.168.16.1 send-community extended
```

!

MP-IBGP Peering: PE-PE

```
lab@Amsterdam> show bgp neighbor
Peer: 192.168.16.1+179 AS 65412 Local: 192.168.24.1+1048 AS 65412
  Type: Internal      State: Established      Flags: <>
  Last State: OpenConfirm  Last Event: RecvKeepAlive
  Last Error: None
  Options: <Preference LocalAddress HoldTime AddressFamily Rib-group Refresh>
  Address families configured: inet-unicast inet-vpn-unicast
  Local Address: 192.168.24.1 Holdtime: 90 Preference: 170
  Number of flaps: 0
  Peer ID: 192.168.16.1      Local ID: 192.168.24.1      Active Holdtime: 90
  Keepalive Interval: 30
  NLRI advertised by peer: inet-unicast inet-vpn-unicast
  NLRI for this session: inet-unicast inet-vpn-unicast
  Peer supports Refresh capability (2)
  Table inet.0 Bit: 10000
    Send state: in sync
    Active prefixes: 0
    Received prefixes: 0
    Suppressed due to damping: 0
  Table bgp.l3vpn.0 Bit: 30000
    Send state: in sync
    Active prefixes: 8
    Received prefixes: 8
    Suppressed due to damping: 0
  Table vpn-a.inet.0 Bit: 40000
    Send state: in sync
    Active prefixes: 7
    Received prefixes: 8
```



Assigning the Route Distinguisher

- Manually assign the RD per VRF table

```
IOS
---
ip vrf ODD_Customer
rd 150:101
...
```

- Automatic RD assignment is possible on some platforms



A Sample VRF Table Configuration

Create a VRF table called *vpn-a* with BGP running between the PE and CE routers using the *vrf-target* statement:

```
ip vrf vpn-a  
rd 3:101
```

```
interface fastethernet 0/0  
ip vrf forwarding vpn-a  
ip address 200.1.9.1 255.255.255.0
```

```
ip vrf vpn-a  
route-target export 3:111  
route-target import 3:111
```